

# **Index Structures of User Profiles for Efficient Web Page Filtering Services**

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## **Outline**

- Introduction**
- Related Approaches**
- Our Approaches**
- Comparisons**
- Conclusion**

## Introduction

### □ Motivation

- ▶ Searching problem on the WWW
  - ☞ search engine
  - ☞ meta-search engine
- ▶ The performance may get worse if the number of web pages grows rapidly

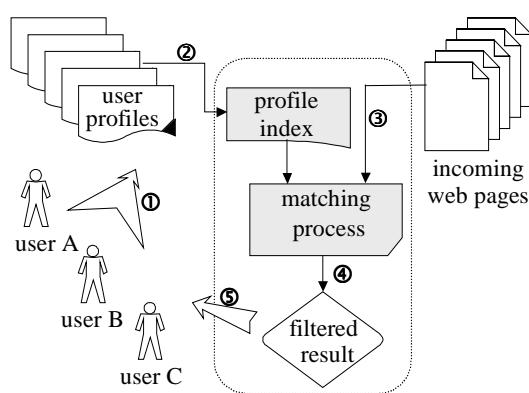
### □ Goal

- ▶ Filtering approach
  - ☞ find the matched profiles for each web page

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## Introduction

### □ A Web Page Filtering Service



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## Introduction

### □ An Example

- ▶ Conjunction of keywords
- ▶ Boolean model
- ▶ Matches
- ▶  $\neg\{P_1, P_4\}$

Profile	Keyword
P <sub>1</sub>	a b
P <sub>2</sub>	a d
P <sub>3</sub>	a d e
P <sub>4</sub>	b f
P <sub>5</sub>	c d e f

Example page
a b c f

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## Related Approaches

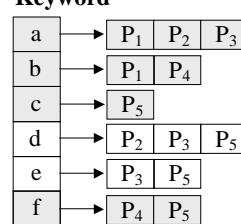
### □ The Counting Method

- ▶ Keyword array: inverted lists
- ▶ Profile arrays: TOTAL, COUNT
- ▶ Matching criteria: COUNT=TOTAL

Profile

	TOTAL	COUNT
P <sub>1</sub>	2	2
P <sub>2</sub>	2	1
P <sub>3</sub>	3	1
P <sub>4</sub>	2	2
P <sub>5</sub>	4	2

Keyword



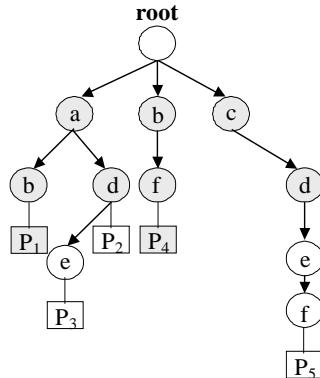
Example page
a b c f

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## Related Approaches

### □ The Tree Method

- ▶ K-nodes: internal nodes
- ▶ P-nodes: leaf nodes
- ▶ External path
  - ☞ root → p-node
  - ☞ a profile
- ▶ Matches
  - ☞ root → a → b → P<sub>1</sub>
  - ☞ root → b → f → P<sub>4</sub>

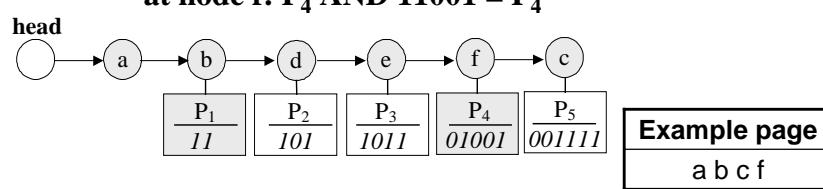


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## Our Approaches

### □ Method 1

- ▶ Index path with path signatures
- ▶ Path signature of the example page: 110011
- ▶ Matches
  - ☞ at node b: P<sub>1</sub> AND 11 = P<sub>1</sub>
  - ☞ at node f: P<sub>4</sub> AND 11001 = P<sub>4</sub>



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## Our Approaches

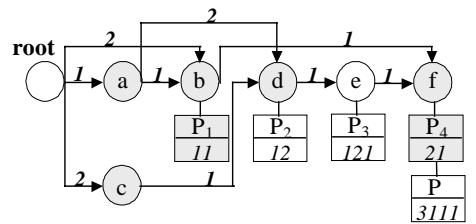
### □ Method 2

► Index graph with path signatures

► Matches

☞ root → a → b → P<sub>1</sub>: 11

☞ root → b → f → P<sub>4</sub>: 21



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## Our Approaches

### □ Method 3

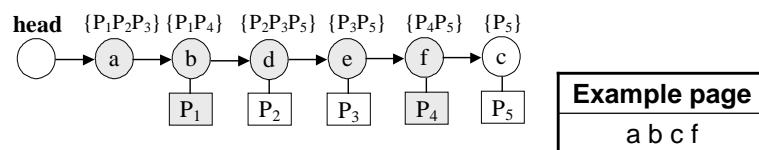
► Index path with profile sets

► Candidate set

► Target set  $\Leftarrow$  candidate set  $\cap$  profile set

☞ at node b: T = {P<sub>1</sub>P<sub>2</sub>P<sub>3</sub>P<sub>4</sub>P<sub>5</sub>}  $\cap$  {P<sub>1</sub>P<sub>4</sub>} = {P<sub>1</sub>P<sub>4</sub>}

☞ at node d: T = {P<sub>2</sub>P<sub>3</sub>P<sub>4</sub>P<sub>5</sub>}  $\cap$  {P<sub>2</sub>P<sub>3</sub>P<sub>5</sub>} = {P<sub>2</sub>P<sub>3</sub>P<sub>5</sub>}

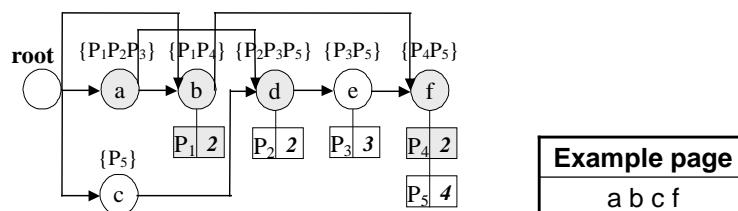


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## Our Approaches

### □ Method 4

- ▶ Index graph with profile sets
- ▶ Path length
- ▶ Matching criteria: matching of keywords  $\wedge$   
profile id  $\subset$  target set  $\wedge$  equal path length



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## Comparisons

### □ Notation

Symbol	Description
P	The set of all profiles
K	The set of all distinct keywords
n	Average number of keywords in a profile
f	Average number of profiles in which a specific keyword is specified
m	Average number of keywords to represent a web page

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## Comparisons

### □ Summary

Approaches Criteria	Counting Method	Tree Method	Method 1	Method 2	Method 3	Method 4
Duplication of information	profile	keyword	no	no	no	no
Sorting of keywords	no	yes	no	yes	no	yes
Storage space	$O( P +f K )$	$O(n P )$	$O( P + K )$	$O( P + K )$	$O( P + K )$	$O( P + K )$
Insertion/Deletion time	$O(nf)$	$O(nf)$	$O(n+f)$	$O(nf)$	$O(n+f)$	$O(n+f)$
Matching time	$O(mf+ P )$	$O(mf)$	$O(mf)$	$O(mf)$	$O(mf)$	$O(mf)$
Modification time	$O(nf)$	$O(nf)$	$O(n+f)$	$O(nf)$	$O(n+f)$	$O(n+f)$

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## Conclusion

### □ Contribution

- ▶ Four new methods for profile indexing
- ▶ Comparisons by complexity analyses
- ▶ Efficient web page filtering service

### □ Future Work

- ▶ Prototype system for real data
- ▶ Dissemination and display of the filtered results
- ▶ More predicates for specifying the user profiles

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