

# PS3 Programming

Week 3. More DMA techniques

Chap 7 and chap 12

# Outline

- Multi-buffered DMA
- DMA request lists
- DMA synchronization
- Homework

# **MULTI-BUFFERED DMA**

# A standard SPU's code

```
/* Read unprocessed data from main memory */
mfc_get(buff, argp, sizeof(buff), TAG, 0, 0);
mfc_write_tag_mask(1<<TAG);
mfc_read_tag_status_all();

/* Process the data */
....
```

```
/* Write the processed data to main memory */
mfc_put(buff, argp, sizeof(buff), TAG, 0, 0);
mfc_write_tag_mask(1<<TAG);
mfc_read_tag_status_all();
```

# Synchronized data access

- A standard process
  - SPUs read data from main memory into LS
  - SPUs process data (need to wait read data)
  - SPU write data from LS to main memory
- There is no multi-threading in SPU to overlap the communication/computation.
  - SPUs are idle when waiting data
  - Use multi-buffer to hide the communication time

# Double buffer

- Create a buffer as large as the incoming data

```
/* The buffer is twice the size of the data */
```

```
vector<unsigned int> buff[SIZE*2]
    __attribute__((aligned(128)));
unsigned short block_size = sizeof(buff)/2;
```

- Let SPUs process the data in one half of the buffer and let MFC get the data in another half of the buffer simultaneously.
- Change the buffer alternatively

```
/* Fill low half with unprocessed data */
mfc_get(buff, argp, block_size, 0, 0, 0);

for(i=1; i<8; i++) {
    /* Fill new buffer with unprocessed data */
    mfc_get(buff +(i&1)*SIZE, argp+i*block_size,
            block_size, i&1, 0, 0);

    /* Wait for old buffer to fill/empty */
    mfc_write_tag_mask(1<<(1-(i&1)));
    mfc_read_tag_status_all();
    /* Process data in old buffer */
    ....

    /* Write data in old buffer to memory */
    mfc_put(buff+(1-(i&1))*SIZE, argp+(i-1)*block_size,
            block_size, 1-(i&1), 0, 0);
}
```

# **DMA REQUEST LISTS**

# Communication cost

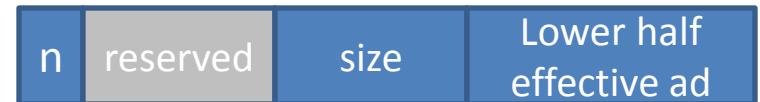
- Communication cost includes
  - # words moved / bandwidth
  - # messages \* latency
- Exponentially growing gaps between
  - Flop/s << 1/Network BW << Network Latency
    - Improving 59%/year vs 26%/year vs 15%/year
  - Flop/s << 1/Memory BW << Memory Latency
    - Improving 59%/year vs 23%/year vs 5.5%/year

# Vectored I/O

- Put multiple memory access calls into one
  - Reduce the latency of I/O
- DMA request list (for **scatter/gather**)
  - Can hold up to 2,048 DMA transfers
  - Maximum data movement is  $2048 \times 16K = 32\text{MB}$
  - Cannot call `mcf_get` and `mcf_put` in the same list
- Two-step process
  - Create a **list element** data structure
  - Call a DMA list function

# DMA list elements

```
typedef struct mfc_list_element{  
    unsigned int notify      : 1;  
    unsigned int reserved   : 16;  
    unsigned int size       : 15;  
    unsigned int eal        : 32;  
} mfc_list_element_t;
```



- Defined in `spu_mfcio.h`
  - Lower half effective address (eal): more latter
  - The size of data to transfer (<16KB)
  - If `notify==1`, the list call stops.

# DMA list functions

- Six functions: `mfc_getl`, `mfc_putl`, `mfc_getlf`,  
`mfc_putlf`, `mfc_getlb`, `mfc_putlb`
- Arguments
  - **volatile** void \*`ls`: LS address of the data
  - unsigned long long `ea`: the effective address (EA)
  - **volatile** `mfc_list_element_t` \*`list`: array of list elm
  - unsigned int `size`: size of list
  - unsigned int `tag`, `tid`, `rid`: same as they are in  
`mfc_get` / `mfc_put`

# Some pitfalls

- Only the most significant 32 bits of **ea** are used. The least significant 32 bits are provided by the **lea** in list elements.
  - Use **mfc\_ea2l**, **mfc\_ea2h**, **mfc\_hl2ea**, and **mfc\_ceil128** to manipulate effective address
- The **argument size** is the size of **\*list**, not the size of data. The size of data is specified by the **size** in list elements.

# Scatter / gather I/O

- There is only one LS address

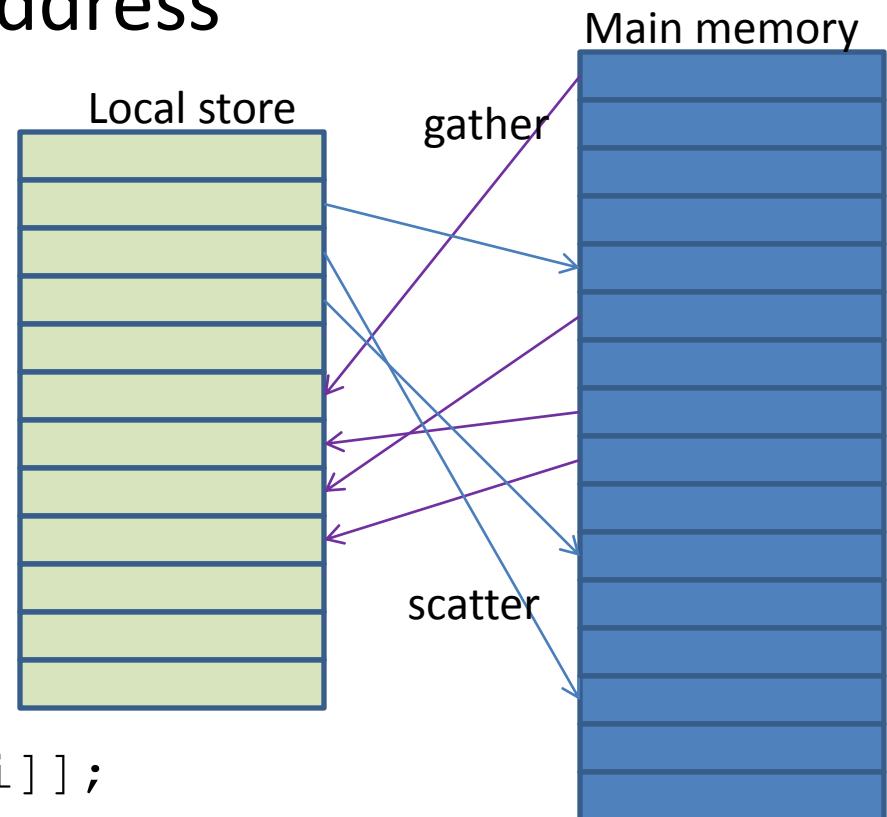
- The initial address of data transfer

- Gather:

```
for(i=0; i<N; ++i)  
    A[i]=B[i]+C[D[i]];
```

- Scatter:

```
for(i=0; i<N; ++i)  ++A[B[i]];
```



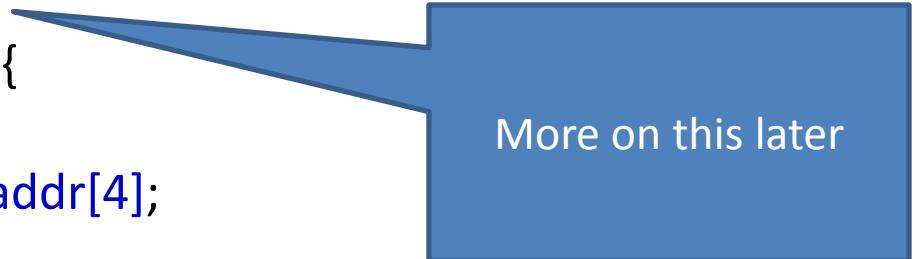
# Example: spu's code

```
#include <spu_mfcio.h>
#include <spu_intrinsics.h>

#define SIZE 4096
#define TAG 3

unsigned int hold_array[SIZE*4] __attribute__ ((aligned (128)));

int main(vector<unsigned long long> arg1,
          vector<unsigned long long> arg2,
          vector<unsigned long long> arg3) {
    unsigned long long get_addr, put_addr[4];
    int i;
```



More on this later

# Get list data

```
/* Retrieve the five addresses from the input parameters */
get_addr = spu_extract(arg1, 0);

/* Create list elements for mfc_getl */
mfc_list_element_t get_element[4];
for (i=0; i<4; i++) {
    get_element[i].size = SIZE*sizeof(unsigned int);
    get_element[i].eal = mfc_ea2l(get_addr) + i*SIZE*sizeof(unsigned int);
}

/* Transfer data into LS */
mfc_getlb(hold_array, get_addr, get_element,
           sizeof(get_element), TAG, 0, 0);
mfc_write_tag_mask(1<<TAG);
mfc_read_tag_status_all();
```

# Put list data

```
/* Retrieve the addresses from the input parameters */
put_addr[0] = spu_extract(arg1, 1);
put_addr[1] = spu_extract(arg2, 0);
put_addr[2] = spu_extract(arg2, 1);
put_addr[3] = spu_extract(arg3, 0);

/* Create list elements for mfc_putl */
mfc_list_element_t put_element[4];
for (i=0; i<4; i++) {
    put_element[i].size = SIZE*sizeof(unsigned int);
    put_element[i].eal = mfc_ea2l(put_addr[i]);
}
/* Transfer data out of LS */
mfc_putl(hold_array, put_addr[0], put_element,
          sizeof(put_element), TAG, 0, 0);
mfc_write_tag_mask(1<<TAG);
mfc_read_tag_status_all();
}
```

# PPU's code

```
#include <stdio.h>
#include <stdlib.h>
#include <libspe2.h>
#include <ppu_intrinsics.h>

#define SIZE 4096
/* Program handle representing the SPU object */
extern spe_program_handle_t spu_dmalist;

int main() {
    int i, retval;  spe_context_ptr_t spe;
    unsigned entry = SPE_DEFAULT_ENTRY;  spe_stop_info_t stop_info;

    /* Array to be transferred into the SPU's LS */
    unsigned int large_array[SIZE*4] __attribute__ ((aligned (128)));
```

```
/* Arrays to hold data transferred out of the LS */
unsigned int small_1[SIZE] __attribute__ ((aligned (128)));
unsigned int small_2[SIZE] __attribute__ ((aligned (128)));
unsigned int small_3[SIZE] __attribute__ ((aligned (128)));
unsigned int small_4[SIZE] __attribute__ ((aligned (128)));

/* Fill the array with whole numbers */
for(i=0; i<SIZE*4; i++)  large_array[i] = i;

/* Initialize the array of array addresses */
unsigned long long control_block[6];
control_block[0] = (unsigned long long)large_array;
control_block[1] = (unsigned long long)small_1;
control_block[2] = (unsigned long long)small_2;
control_block[3] = (unsigned long long)small_3;
control_block[4] = (unsigned long long)small_4;
control_block[5] = (unsigned long long)NULL;
```

```
/* Create the SPE Context */
spe = spe_context_create(0, NULL);
if (!spe) {
    perror("spe_context_create"); exit(1);
}

/* Load the program into the context */
retval = spe_program_load(spe, &spu_dmalist);
if (retval) {
    perror("spe_program_load"); exit(1);
}

/* Run the program inside the context */
retval = spe_context_run(spe, &entry, SPE_RUN_USER_REGS,
                        control_block, NULL, &stop_info);
if (retval < 0) {
    perror("spe_context_run"); exit(1);
}
```

```
/* Deallocate the context */
retval = spe_context_destroy(spe);
if (retval) {
    perror("spe_context_destroy");
    exit(1);
}

/* Check that the received data is valid */
int test = 1;
for(i=0; i<SIZE; i++) {
    if(small_1[i] != i)          test = 0;
    if(small_2[i] != SIZE+i)    test = 0;
    if(small_3[i] != 2*SIZE+i)  test = 0;
    if(small_4[i] != 3*SIZE+i)  test = 0;
}
if(test)  printf("DMA Passed.\n");
else     printf("DMA Failed.\n");
return 0;
}
```

# **SYNCHRONIZATION**

# Synchronization library (libsync)

- Functions are divided into five categories
  1. Atomic operations: uninterruptable operations
    - atomic\_read, atomic\_set, atomic\_inc, ...
  2. Mutexes
  3. Reader/writer locks
  4. Condition variable operations
  5. Completion operations
    - A special condition variable that signals others when the thread is finished.

# Mutex (mutual exclusive)

- Protect critical regions

```
mutex_lock(a)  
// critical region (CR)  
// access mutual exclusive resources.  
mutex_unlock(a)
```

Spinning lock  
Wait until a==0  
And then set a=1

- Functions

```
mutex_init(mutex_ea_t lock)  
mutex_lock(mutex_ea_t lock)  
mutex_trylock(mutex_ea_t lock)  
mutex_unlock(mutex_ea_t lock)
```

# Reader/writer locks

- A set of more sophisticated mutex functions
  - Mutex only allows one thread to enter the CR
  - Sometimes, it is ok for more than 1 threads to access data, like `read`. Mutual exclusion only happens when there is a `write` operation.
- Rules: suppose the lock is “a”. (typed `mutex_ea_t`)
  - If  $a \neq 0$ , the `write_lock` causes a spinning lock
    - After it sees  $a == 0$ , it sets  $a = -1$
  - If  $a == -1$ , `read_lock` causes spinning lock
    - After it sees  $a \neq -1$ , it sets  $a = a + 1$

# Condition variables

- The goal of a condition variable is to have all threads wait until a condition happens.
  - The condition is triggered by other threads.
- Condition variables are always associated with mutexes to avoid race conditions.

```
mutex_lock(a)  
cond_wait(cond, a)  
mutex_unlock(a)
```

It will release “a”. Until  
the condition is satisfied,  
it re-locks “a”.

```
mutex_lock(a)  
cond_signal(cond)  
mutex_unlock(a)
```

Or use  
cond\_croadcast

# SPU's cashier code

```
#include <spu_mfcio.h>
#include <spu_intrinsics.h>
#include <libsync.h>
#define TAG 3

/* SPU initialization data */
typedef struct _control_block {
    cond_ea_t cashier;
    mutex_ea_t cashier_mutex, served_mutex;
    unsigned long long served_addr;
} control_block;

control_block cb __attribute__ ((aligned (128)));

int main(unsigned long long speid, unsigned long long argp,
        unsigned long long envp) {
```

```
/* Get the control block from main memory */
mfc_get(&cb, argp, sizeof(cb), TAG, 0, 0);
mfc_write_tag_mask(1<<TAG); mfc_read_tag_status_all();

/* Enter the store: get lock to wait for cashier */
mutex_lock(cb.cashier_mutex);
/* Wait for cashier */
cond_wait(cb.cashier, cb.cashier_mutex);
/* Allow others to wait for the cashier */
mutex_unlock(cb.cashier_mutex);

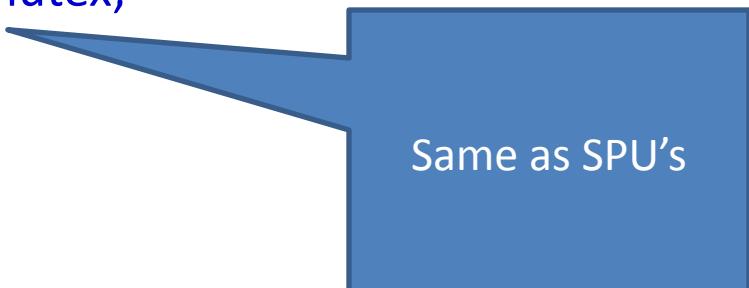
/* Leave the store: get lock to increment num_served */
mutex_lock(cb.served_mutex);
/* Increment the number of customers served */
atomic_inc((atomic_ea_t)cb.served_addr);
printf("Thread %llu incremented num_served to %u\n",
    speid, atomic_read((atomic_ea_t)cb.served_addr));
/* Allow others to access num_served */
mutex_unlock(cb.served_mutex);
return 0;
}
```

# PPU's cashier code

```
/* SPU initialization data */
typedef struct _control_block {
    cond_ea_t cashier;
    mutex_ea_t cashier_mutex, served_mutex;
    unsigned long long served_addr;
} control_block;

/* SPU program handle */
extern spe_program_handle_t spu_cashier;
ppu(pthread_data_t data[6];
control_block cb __attribute__ ((aligned (128)));

/* Declare variables */
volatile int cashier_var __attribute__ ((aligned (128)));
volatile int cashier_mutex_var __attribute__ ((aligned (128)));
volatile int served_mutex_var __attribute__ ((aligned (128)));
volatile int num_served __attribute__ ((aligned (128)));

Same as SPU's


```

```
/* The data sent to the pthread */
typedef struct ppu_pthread_data {
    spe_context_ptr_t speid;
    pthread_t pthread;
    void *argp;
} ppu_pthread_data_t;

/* The function executed in the pthread */
void *ppu_pthread_function(void *arg) {
    ppu_pthread_data_t *data = (ppu_pthread_data_t *)arg;
    int retval;
    unsigned int entry = SPE_DEFAULT_ENTRY;
    if ((retval = spe_context_run(data->speid, &entry, 0, data->argp, NULL,
NULL)) < 0) {
        perror("spe_context_run");
        exit(1);
    }
    pthread_exit(NULL);
}
```

```
int main(int argc, char **argv) {
    int i, retval, spus;

    /* Create condition variable, mutexes, pointer */
    cb.cashier = (cond_ea_t)&cashier_var;
    cb.cashier_mutex = (mutex_ea_t)&cashier_mutex_var;
    cb.served_mutex = (mutex_ea_t)&served_mutex_var;
    cb.served_addr = (unsigned long long)&num_served;

    /* Initialize condition variable, mutexes */
    cond_init(cb.cashier);
    mutex_init(cb.cashier_mutex);
    cond_init(cb.served_mutex);
    num_served = 0;

    /* Determine number of available SPUs */
    spus = spe_cpu_info_get(SPE_COUNT_USABLE_SPES, 0);
```

```
/* Create contexts and threads */
for (i=0; i<spus; i++) { /* Create context */
    if ((data[i].speid = spe_context_create(0, NULL)) == NULL) {
        perror("spe_context_create");
        exit(1);
    }

    /* Load program into the context */
    if ((retval = spe_program_load(data[i].speid, &spu_cashier)) != 0) {
        perror("spe_program_load");
        exit (1);
    }
    data[i].argp = &cb;

    /* Create thread */
    if ((retval = pthread_create(&data[i].pthread,
                                NULL, &ppu_thread_function, &data[i])) != 0) {
        perror("pthread_create");
        exit (1);
    }
}
```

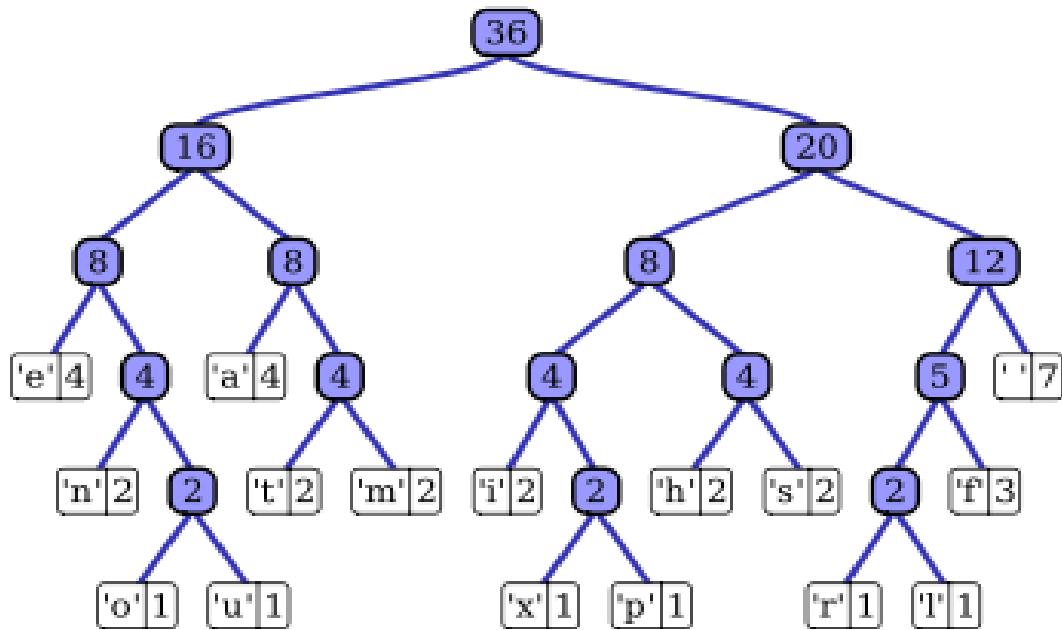
```
int count = 0;
while (count < spus) {
    mutex_lock(cb.cashier_mutex);
    cond_signal(cb.cashier);
    count = atomic_read((atomic_ea_t)cb.served_addr);
    mutex_unlock(cb.cashier_mutex);
}

for (i = 0; i < spus; i++) { /* Wait for the threads */
    if ((retval = pthread_join (data[i].pthread, NULL)) != 0) {
        perror("pthread_join");
        exit (1);
    }

    /* Deallocate the contexts */
    if ((retval = spe_context_destroy (data[i].speid)) != 0) {
        perror("spe_context_destroy");
        exit (1);
    }
}
return 0;
```

# **HOMEWORK**

# Huffman coding



- Given a set of symbols and their counts, find a **prefix-free binary code** with minimum expected codeword length

Char	Freq	Code
space	7	111
a	4	010
e	4	000
f	3	1101
h	2	1010
i	2	1000
m	2	0111
n	2	0010
s	2	1011
t	2	0110
l	1	11001
o	1	00110
p	1	10011
r	1	11000
u	1	00111
x	1	10010

# Huffman coding for compression

1. Read in a file (binary or text)
2. Count its char frequency
3. Build the Huffman tree according to the char frequency
  - Use priority queue (heap sort)
4. Write out the Huffman coded file
  - The code table and the compressed data
5. Design the decompression function (optional)

# Homework

- Implement a sequential Huffman compression and decompression codes using PPU
  - Use as many vectored statements as possible
  - You don't need to implement a priority queue, but it would be a good practice
  - You can designed your own compression file format
- Use SPUs to help