

CS5371

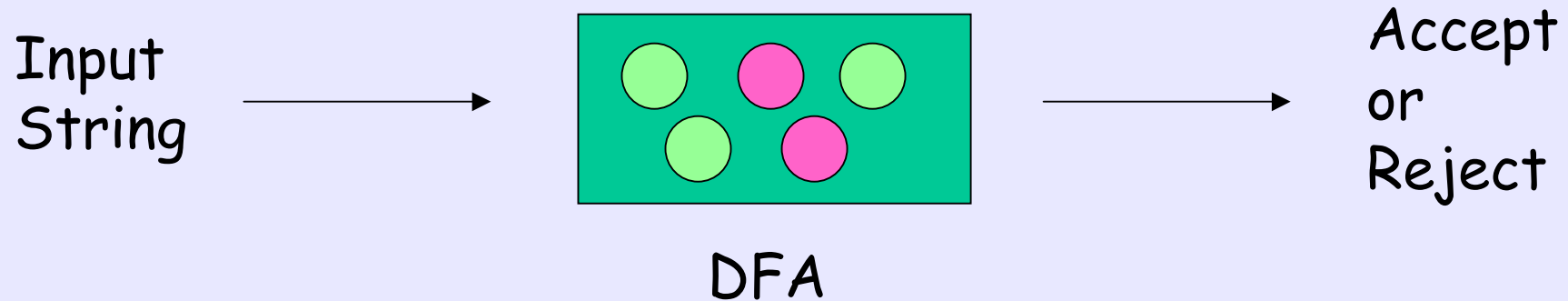
Theory of Computation

Lecture 3: Automata Theory I
(DFA and NFA)

Objectives

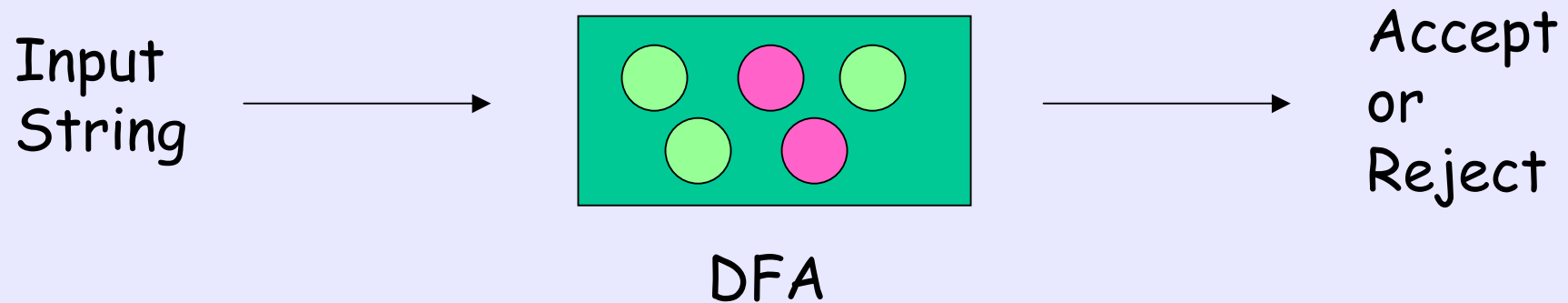
- This time, we will look at how to define a very simple "computer" called deterministic finite automaton (DFA)
- Show that DFA can solve some string decision problem
- Then, we give slightly change the definition of DFA to give another computer called non-deterministic finite automaton (NFA)

DFA



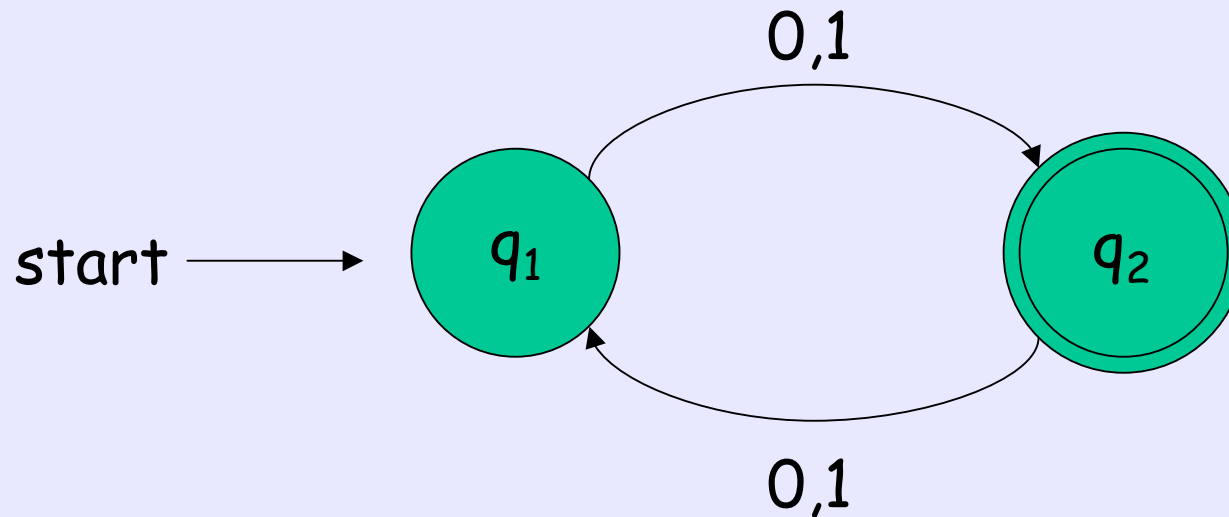
- A machine with finite number of **states**, some states are **accepting** states, others are **rejecting** states
- At any time, it is in one of the states
- It reads an input string, one character at a time

DFA



- After reading each character, it moves to another state depending on **what is read** and **what is the current state**
- If reading all characters, the DFA is in an accepting state, the input string is **accepted**.
- Otherwise, the input string is **rejected**.

Example of DFA

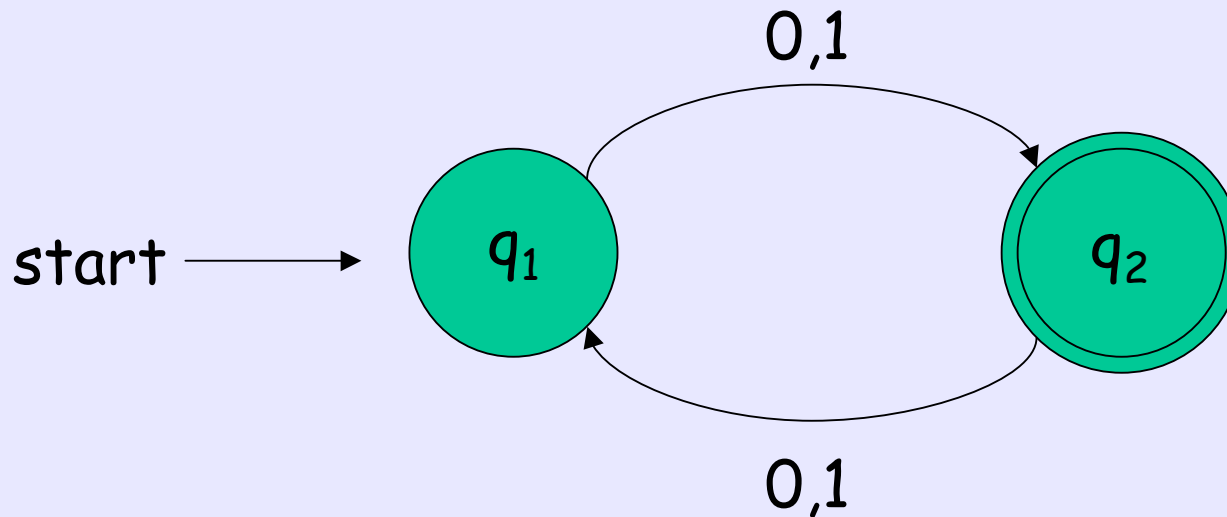


- The circles indicates the states
- If **accepting** state is marked with double circle
- The arrows pointing from a state q indicates how to move on reading a character when current state is q

Formal Definition of DFA

- A DFA is a 5-tuple $(Q, \Sigma, \delta, q_{\text{start}}, F)$, where
 - Q is a set consisting finite number of **states**
 - Σ is an **alphabet** consisting finite number of characters
 - $\delta: Q \times \Sigma \rightarrow Q$ is the **transition function**
 - q_{start} is the **start state**
 - F is the set of **accepting states**
- Note: only 1 start state, and can have many accepting states

Formal Definition of DFA



$$Q = \{q_1, q_2\}, \Sigma = \{0, 1\}, q_{\text{start}} = q_1, F = \{q_2\}$$

$$\delta(q_1, 0) = q_2, \quad \delta(q_1, 1) = q_2$$

$$\delta(q_2, 0) = q_1, \quad \delta(q_2, 1) = q_1$$

Formal Definition of DFA's Computation

- Recall how a DFA performs its computation (to decide whether an input string is accepted or rejected):
 - If after reading the whole string, DFA is in an accepting state, then the input string is accepted; otherwise the input string is rejected
- How to define this formally??

Formal Definition of Computation

- Let $M = (Q, \Sigma, \delta, q_{\text{start}}, F)$ be a DFA
- Let $w = w_1 w_2 \dots w_n$ be a string with each w_i a member of the alphabet Σ
- Then, M **accepts** w if a sequence of states r_0, r_1, \dots, r_n in Q exists with the three conditions:
 - $r_0 = q_{\text{start}}$
 - $\delta(r_i, w_{i+1}) = r_{i+1}$
 - $r_n \in F$

Some Terminology

Let M be a DFA

- Among all possible strings, M will accept some of them, and M will reject the remaining
- The set of strings which M accepts is called the language **recognized** by M
- That is, M **recognizes** A if
$$A = \{ w \mid M \text{ accepts } w \}$$

Quick Quiz

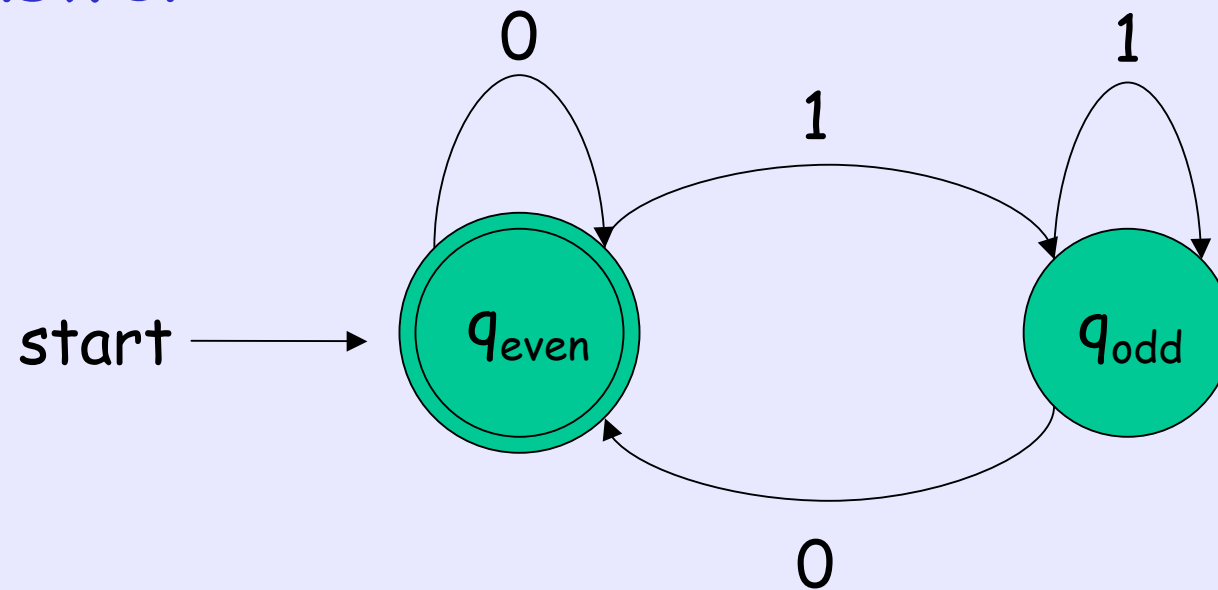
- What is an **alphabet**?
- What is a **language**?
- What are the min and max numbers of strings a DFA can **accept**?
- What are the min and max numbers of languages a DFA can **recognize**?

Designing a DFA [Example 1]

- How to design a DFA that accepts all binary strings that represent an even number? (e.g., accepts 110, 010, but rejects 111, 010101)
- Let's pretend ourselves as a DFA...
 - After reading a character, we must decide whether the string seen so far is in the language (what is the language we're talking now??) because we don't know if it is the last character...

Designing a DFA

Answer:



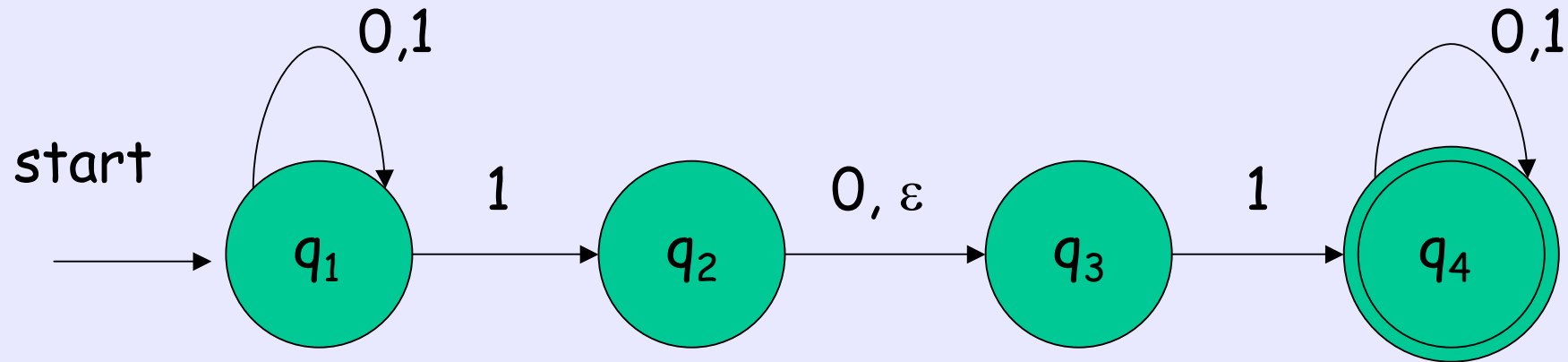
Designing a DFA (Quick Quiz)

- How to design a DFA that accepts all binary strings representing a multiple of 5? (E.g., 101, 1111, 11001, ...)

NFA

- Every step of DFA's computation follows in a unique way from the preceding step
 - When a machine is in a given state, and reads the next input character, we know what the next state is
- In an NFA, several choices may exist for the next state

Example of NFA



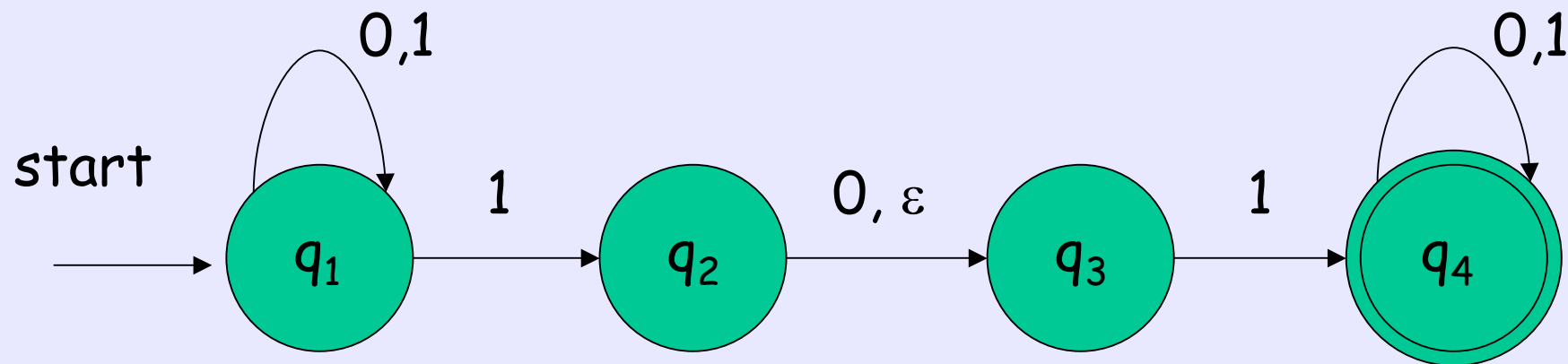
- Difference with DFA

- Can move to **more than 1 states**, or **nowhere**, after reading a character (E.g., at q_1 , on reading 1, we can move to q_1 or q_2 ; at q_2 , on reading 1, nowhere to go!)
- Can move to another state **without reading anything** (E.g., at q_2 , the symbol ϵ on arrow pointing at q_3 indicates that we can move to q_3 without reading anything)

What is accepted by NFA?

- Let M be an NFA
- (Informally,) M accepts a string w if there is at least one way to move from the start state to a final state according to the transition arrows, such that the concatenation of the true characters (that is, ignoring ε) used by the transition is equal to w

What is accepted by this?



Next time

- Continue the discussion on NFA
- Properties of language recognized by DFA or NFA