Randomized algorithm

Randomized algorithm Tutorial 6

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Solution to assignment 4

Question 1 Question 2 Question 3 Question 4

Hint for assignment 5

Question 1 Question 2 Randomized algorithm

Solution to assignment 4

Solution to assignment 4

[Question 1]: Let G be a random graph drawn from the $G_{n,p}$ model.

- 1. (10%) What is the expected number of 5-clique in G?
- 2. (10%) What is the expected number of 5-cycle in G?

[Solution]:

- 1. Number of 5-vertex set: $\binom{n}{5}$ Probability of a 5-clique : p^{10} E[number of 5-clique in G] = $\binom{n}{5} \cdot p^{10}$.
- 2. Each 5-vertex set can compose 5! cycles but each cycle can be represented in 10 different ways. Ex: {1,2,3,4,5} {2,3,4,5,1} {3,4,5,1,2} {4,5,1,2,3} {5,1,2,3,4} The total number of 5-cycle is $\binom{n}{5} \cdot \frac{5!}{10} = 12\binom{n}{5}$

E[number of 5-cycle in
$$G$$
] = $12 \binom{n}{5} p^5$

[Question 2]: Suppose we have a set of *n* vectors, v_1, v_2, \ldots, v_n , in R^m . Each vector is of unit-length, i.e., $||v_i|| = 1$ for all *i*. In this question, we want to show that, there exists a set of values, $\rho_1, \rho_2, \ldots, \rho_n$, each $\rho_i \in \{-1, +1\}$, such that

$$\|\rho_1\mathbf{v}_1+\rho_2\mathbf{v}_2+\cdots+\rho_n\mathbf{v}_n\|\leq\sqrt{n}.$$

Intuitively, if we are allowed to "reflect" each v_i as we wish (i.e., by replacing v_i by $-v_i$), then it is possible that the vector formed by the sum of the *n* vectors is at most \sqrt{n} long.

1. Let $V = \rho_1 v_1 + \rho_2 v_2 + \dots + \rho_n v_n$, and recall that $\|V\|^2 = V \cdot V = \sum_{i,j} \rho_i \rho_j v_i \cdot v_j.$

Suppose that each ρ_i is chosen uniformly at random to be -1 or +1. Show that

$$\mathbf{E}[\|V\|^2] = n.$$

- 2. (5%) Argue that there exists a choice of $\rho_1, \rho_2, \ldots, \rho_n$ such that $\|V\| \leq \sqrt{n}$.
- 3. (5%) Your friend, Peter, is more ambitious, and asks if it is possible to to choose $\rho_1, \rho_2, \ldots, \rho_n$ such that

$$\|V\| < \sqrt{n}$$

instead of $||V|| \le \sqrt{n}$ we have just shown. Give a counter-example why this may not be possible.

Randomized algorithm Solution to assignment 4 Question 2

> [Solution 1]: Given that

$$\|V\|^2 = V \cdot V,$$

we have

$$\mathbf{E}[\|V\|^2] = \mathbf{E}[V \cdot V] = \mathbf{E}\left[\sum_{i,j} \rho_i \rho_j v_i \cdot v_j\right].$$

When $i \neq j$, ρ_i and ρ_j are independent random variables, so that

$$\mathbf{E}[\rho_i \rho_j] = \mathbf{E}[\rho_i] \mathbf{E}[\rho_j] = \mathbf{0}$$

On the other hand, when i = j, we have

$$E[\rho_i \rho_i] = 0.5(1)^2 + 0.5(-1)^2 = 1.$$

Also, $v_i \cdot v_i = v_i^2 = 1$ since v_i is a unit vector. Combining the above, we obtain

$$E[||V||^{2}] = E\left[\sum_{i,j} \rho_{i}\rho_{j}\mathbf{v}_{i}\cdot\mathbf{v}_{j}\right]$$
$$= E\left[\sum_{i} \rho_{i}\rho_{i}\mathbf{v}_{i}\cdot\mathbf{v}_{i}\right] + E\left[\sum_{i\neq j} \rho_{i}\rho_{j}\mathbf{v}_{i}\cdot\mathbf{v}_{j}\right]$$
$$= \sum_{i} E\left[\rho_{i}\rho_{i}\right] + \sum_{i\neq j} E\left[\rho_{i}\rho_{j}\mathbf{v}_{i}\cdot\mathbf{v}_{j}\right]$$
$$= n + 0 = n.$$

[Solution 2]: Since $E[||V||^2] = n$, there must be a choice of reflection ρ 's such that $||V||^2$ is at most n. With that particular choice, $||V|| \le \sqrt{n}$. [Solution 3]: Set $v_1 = (1,0)$ and $v_2 = (0,1)$. No matter what ρ_1 and ρ_2 is, $\rho_1 \rho_2 v_1 \cdot v_2 = \sqrt{2}$. In general in \mathbb{R}^m , we can set

$$v_k = (\underbrace{0,\ldots,0}_{k-1 \text{ zeroes}}, 1, 0, \ldots, 0)$$

for k = 1 to m. Then under any choice of the reflections, the resulting vector V will have length exactly \sqrt{m} .

[Question 3]: Let F be a family of subsets of $N = \{1, 2, ..., n\}$. F is called an antichain if there are no $A, B \in F$ satisfying $A \subset B$. Ex: $N = \{1, 2, 3\}$

All subsets of N :{}, {1}, {2}, {3}, {1,2}, {1,3}, {2,3}, {1,2,3} {{1,3}, {1,2}, {2,3}} is an antichain. {{1,2,3}, {1,2}} is not an antichain. ({1,2} \subset {1,2,3})

We want to prove that $|F| \leq {n \choose \lfloor n/2 \rfloor}$. How? Let $\sigma \in S_n$ be a random permutation of the elements of N and consider the random variable

$$X = |\{i : \{\sigma(1), \sigma(2), ... \sigma(i)\} \in F\}|$$

What is X?

Let us set
$$F = \{\{2\}, \{1,3\}\}$$
.
For $\sigma = \{2,3,1\}, X = 1$. $\because \{2\} \in F$.
For $\sigma = \{3,2,1\}, X = 0$. \because None of $\{3\}, \{3,2\}, \{3,2,1\}$ are in F .
For $\sigma = \{3,1,2\}, X = 1$. $\because \{1,3\} \in F$.
Try to find the relation between X and F .

[Solution]: Partition F into n sets, k_1, k_2, \dots, k_n where k_i is the number of those sets whose size is i. Therefore, $|F| = k_1 + k_2 + \dots + k_n$. The expected number of X is the can be considered as all the possible case in all permutations.

$$E[X] = (k_1 \cdot (n-1)! \cdot 1! + k_2 \cdot (n-2)! \cdot 2! + \dots + k_n \cdot n!)/n!$$

= $\frac{k_1}{\binom{n}{1}} + \frac{k_2}{\binom{n}{2}} + \dots + \frac{k_n}{\binom{n}{n}}$
 $\geq \frac{(k_1 + k_2 + \dots + k_n)}{\binom{n}{n/2}}$
= $\frac{|F|}{\binom{n}{n/2}}$

As we have already known that X is an indicator random variable, $E[X] \le 1$. Thus, $\frac{|F|}{\binom{n}{n/2}} \le 1$.

[Question 4]: Consider a graph in $G_{n,p}$, with p = 1/n. Let X be the number of triangles in the graph, where a triangle is a clique with three edges. Show that

$$Pr(X \ge 1) \le 1/6$$

and that

 $\lim_{n\to\infty}\Pr(X\geq 1)\geq 1/7$

[Solution 1]:

$$Pr(X \ge 1) = \binom{n}{3}p^3$$
$$= \frac{n(n-1)(n-2)}{6n^3}$$
$$\le \frac{1}{6}$$

-Solution to assignment 4

Question 4

[Solution 2]: Consider all possible size-3 subsets of the vertices of the random graph generated by $G_{n,p}$. Let X_i be an indicator that denote *i*th size-3 subset forms a triangle in G.

Now we consider $E[X|X_i = 1]$.

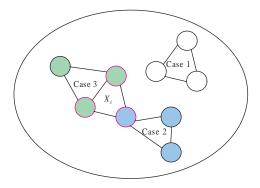


Figure: 4 cases. Case 1: X_i ; Case 2: p^3 ; Case 3: p^3 ; Case 4: p^2 .



By conditional expectation inequality, $Pr(X) = \sum_{i=1}^{\binom{n}{3}} \frac{Pr(X_i=1)}{E[X|X_i=1]}$

$$\begin{split} \mathbf{E}[X|X_i = 1] &= 1 + \binom{n-3}{3} p^3 + 3\binom{n-3}{2} p^3 + 3\binom{n-3}{1} p^2 \\ \Pr(X \ge 1) &= \frac{\binom{n}{3} p^3}{1 + \binom{n-3}{3} p^3 + 3\binom{n-3}{2} p^3 + 3\binom{n-3}{1} p^2} \\ &= \frac{1/6}{1 + 1/6 + 0 + 0} = \frac{1}{7} \end{split}$$

 $\lim_{n\to\infty}\Pr(X\geq 1)\geq 1/7$

Randomized algorithm

Hint for assignment 5

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└─Hint for assignment 5
Question 1

[Question 1]: Consider the Markov chain in the following figure. Suppose that $n \ge 4$. Find the expected number of moves to reach n starting from position i, when (a) i = n - 1, (b) i = n - 2, (c) i = n - 3, and (d) i = n - 4. (Express the answers exactly in terms of n.)

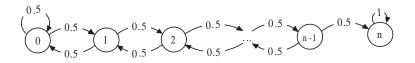


Figure: The modeling Markov chain for 2-SAT problem.

[Question 2]: Consider the Markov chain in the following figure.

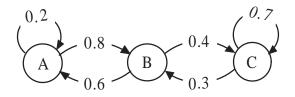
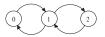


Figure: The modeling Markov chain of this question.

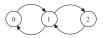
- 1. (20%) Argue that the Markov chain is aperiodic and irreducible.
- 2. (20%) Find the stationary probability.
- 3. (20%) Find the probability of being in state 0 after 16 steps if the chain begins at state chosen uniformly at random from the 3 states.

[Hint]:

• Irreducible: Every state j can reach every state k



Periodic: Once we start at state j, we can only return to j after a multiple of d steps.



Stationary distribution: A probability distribution that p(n) remains a certain distribution when n > t.



Thank you