#### Randomized algorithm

Tutorial 6 Solution for Assignment 3 Hint for Assignment 4

#### Solution for assignment 3

- Let X be a Poisson random variable with mean  $\mu$ .
  - a) What is the most likely value of X when  $\lambda$  is an integer?
    - $\lambda$  is not an integer?

• [Sol]  $Pr(X = k) = \frac{\lambda^k e^{-\lambda}}{k!}$  $\Pr(X = k + 1) = \frac{\lambda^{k+1} e^{-\lambda}}{(k+1)!}$  $\frac{\Pr(X=k+1)}{\Pr(X=k)} = \frac{\lambda}{k+1}$ 

$$k < \lambda \rightarrow \Pr(X = k + 1) > \Pr(X = k)$$
$$k = \lambda \rightarrow \Pr(X = k + 1) = \Pr(X = k)$$
$$k > \lambda \rightarrow \Pr(X = k + 1) < \Pr(X = k)$$

The most likely value of *X* is when  $X = \lfloor \lambda \rfloor$ . If  $\lambda$  is an integer, both  $\lambda$  and  $\lambda - 1$  are the most likely values.

- We define the median of X to be the b) least number *m* such that  $Pr(X \le m) \ge 1/2$ . What is the median of X when  $\lambda = 3.9$ ? [Sol] Pr(X = 0) = 0.020 $Pr(X = 1) = 0.079 \implies 0.099$  $Pr(X = 2) = 0.154 \implies 0.253$  $Pr(X = 3) = 0.200 \Rightarrow 0.453$ 
  - $Pr(X = 4) = 0.195 \Longrightarrow 0.648$



Are Y and Z independent?
[Sol] Pr(Y = k)

$$= \sum_{m=k}^{\infty} \Pr(X = m) \Pr(Y = k \mid X = m)$$
$$= \sum_{m=k}^{\infty} \frac{\lambda^m e^{-\lambda}}{m!} {m \choose k} p^k (1-p)^{m-k}$$
$$= \frac{(\lambda p)^k e^{-p\lambda}}{k!}$$

$$Pr(Y = k \text{ and } Z = j)$$

$$= Pr(X = k + j) Pr(Y = k | X = k + j)$$

$$= \frac{\lambda^{k+j} e^{-\lambda}}{(k+j)!} {k \choose k} p^k (1-p)^j$$

$$= \frac{(\lambda p)^k e^{-p\lambda}}{k!} \cdot \frac{(\lambda (1-p))^j e^{-(1-p)\lambda}}{j!}$$

$$= Pr(Y = k) Pr(Z = j)$$

a) Now, b balls are in play.
 f(b): the expected number of balls that survive to the subsequent round.
 Given an explicit formula for f(b).

Pr(*i*th bin has exactly 1 ball) = 
$$b \frac{1}{n} (1 - \frac{1}{n})^{b-1}$$

E[number of bins with 1 ball] =  $n \operatorname{Pr}(i \operatorname{th...}) = b(1 - \frac{1}{n})^{b-1}$ 

$$f(b) = b - E[\text{number of bins with 1 ball}] = b(1 - (1 - \frac{1}{n})^{b-1})$$

b) Show that  $f(b) \leq b^2/n$ . By Bernoulli's inequality:  $(1 - \frac{1}{n})^{b-1} \geq 1 - \frac{b-1}{n}$   $f(b) = b(1 - (1 - \frac{1}{n})^{b-1})$  $\leq b - b(1 - \frac{b-1}{n})$ 

$$=\frac{b(b-1)}{n} \le \frac{b^2}{n}$$

Suppose that every round the number of balls served was exactly the expected value.
 Show that all the balls would be served in O(log log *n*) rounds.

# Exercise 3 : Solution

Suppose we have n / k balls initially, for some fixed constant k > 1.

From part (b),

 $\mathsf{f}(n/k) \leq n/k^2.$ 

After r rounds,

 $f^{(r)}(n/k) \le n/k^s$  where  $s = 2^r$ 

When  $r = \log_k \log_2 n = O(\log \log n)$ ,  $f^{(r)}(n/k) \le 1$ 

Now, consider about the case that we have *n* balls initially.

If *n* is large, #balls after 1st round is :

$$\lim_{n \to \infty} n(1 - (1 - \frac{1}{n})^{n-1})$$
  
= 
$$\lim_{n \to \infty} n(1 - (1 - \frac{1}{n-1})^{n-1})$$
  
= 
$$n(1 - \frac{1}{e})$$

a) Argue that the maximum load in this case is only  $O(\log \log n / \log \log \log n)$  with probability that approaches 1 as  $n \rightarrow \infty$ .

 Among all *n* bins, we choose log *n* bins (evenly) as representatives.



#### **Exercise 4 : Solution**

 $Pr(\exists bin \ge 2M balls)$   $\leq Pr(\exists Rep \ge M)$   $\leq (\log n)(\frac{1}{M!})$   $\leq \frac{1}{\log n} \qquad (if M = \frac{\log \log n}{\log \log \log n})$  $\rightarrow 0 \qquad (when n \rightarrow \infty)$ 

Consider *n* balls thrown randomly into *n* bins

Let  $X=X_1+X_2+...+X_n$ , where  $X_i = 1$  if *i*-th bin is empty ; 0 otherwise.

• Let 
$$Y=Y_1+Y_2+...+Y_n$$
, where  
each  $Y_i$  is an independent Bernoulli  
random variables with

 $Pr(Y_i = 1) = (1-1/n)^n$ .

a) Show that  $E[X_1X_2...X_k] \leq E[Y_1Y_2...Y_k]$ .

$$E[X_1 X_2 \dots X_k] = \Pr(X_1 = 1 \cap X_2 = 1 \cap \dots \cap X_k = 1) = (\frac{n-k}{n})^n$$
$$E[Y_1 Y_2 \dots Y_k] = \Pr(Y_1 = 1 \cap Y_2 = 1 \cap \dots \cap Y_k = 1) = (1 - \frac{1}{n})^{kn}$$

$$\left(\frac{n-k}{n}\right)^n \le \left(1-\frac{1}{n}\right)^{kn}$$

## Exercise 5(a) : Alternative solution

• By induction  $E[X_{1}X_{2}...X_{k}X_{k+1}]$   $= E[X_{1}X_{2}...X_{k}|X_{k+1}=1]Pr(X_{k+1}=1)$   $\leq E[X_{1}X_{2}...X_{k}]Pr(X_{k+1}=1)$   $\leq (1 - \frac{1}{n})^{(k+1)n}$   $= E[Y_{1}Y_{2}...Y_{k}Y_{k+1}]$ 

b) Show that  $X_{i1}^{k1} X_{i2}^{k2} \cdots X_{ij}^{kj} = X_{i1} X_{i2} \cdots X_{ij}$ 

 $X_i$  is an indicator

Thus,

$$X_{i1}^{k1}X_{i2}^{k2}...X_{ij}^{kj} = 1 = X_{i1}X_{i2}...X_{ik}$$

or

$$X_{i1}^{k1}X_{i2}^{k2}...X_{ij}^{kj} = 0 = X_{i1}X_{i2}...X_{ik}$$

c) Show that  $E[e^{tX}] \leq E[e^{tY}]$   $E[e^{tX}]$   $= E[1+tX + \frac{(tX)^2}{2!} + ...]$   $= E[1] + tE[X] + t^2 E[\frac{X^2}{2!}] + ...$  $E[X] \leq E[Y]$  (by (a))

 $E[X^{2}] = E[(X_{1} + X_{2} + ... + X_{n})^{2}]$ =  $E[X_{1}^{2} + 2(X_{1}X_{2} + X_{1}X_{3} + ... + X_{1}X_{n}) + X_{2}^{2} + ... + X_{n}^{2}]$  $\leq E[Y_{1}^{2} + 2(Y_{1}Y_{2} + Y_{1}Y_{3} + ... + Y_{1}Y_{n}) + Y_{2}^{2} + ... + Y_{n}^{2}]$ =  $E[Y^{2}]$ 

 $E[e^{tX}]$ = E[1] + tE[X] + t<sup>2</sup>E[ $\frac{X^{2}}{2!}$ ] + ...  $\leq E[1] + tE[Y] + t^{2}E[\frac{Y^{2}}{2!}]$  + ... = E[ $e^{tY}$ ]

d) Derive a Chernoff bound for  $Pr(X \ge (1 + \delta)E[X])$ 

 $E[e^{tY}] = \Pi E[e^{tY_i}]$  $= [(1-p) + pe^t]^n$  $= [1+p(e^t - 1)]^n$  $\leq e^{np(e^t - 1)}$ 

 $\Pr(X \ge (1+\delta)E[X])$  $= \Pr(e^{tX} \ge e^{t(1+\delta)E[X]})$  $\leq \frac{E[e^{tX}]}{e^{t(1+\delta)\mu}} \leq \frac{E[e^{tY}]}{e^{t(1+\delta)\mu}} \quad (\text{Set E}[X] = \mu)$  $=\frac{e^{\mu(e^t-1)}}{e^{t(1+\delta)\mu}} \quad \text{(Choose } t = \ln(1+\delta)\text{)}$  $\leq \frac{e^{\mu\delta}}{(1+\delta)^{(1+\delta)\mu}}$ 

Hint for assignment 4

# *k*-uniform hypergraph V = {A,B,C,D,E,F,G} E = { {ABC}, {CDE}, {EFG} }



Given a *k*-uniform hypergraph with  $E = \{S_1, S_2, S_3, ..., S_{|C|}\},$   $|C| \leq 4^{k-1}-1,$  and  $k \geq 2.$ Show that there exists a 4-coloring such that no *k*-set is monochromatic.

[Hint] You can do it without any hint.

Anti-chain

F, a family of subsets of  $N=\{1,2,...,n\}$  is called *anti-chain* if there are no A, B in F satisfying A in B.

Ex: F={ {1,3,4} , {2,4} , {1,5} , {6} }

What if F={ {1,3,4}, {2,4}, {1,4}, {6} }?

- Let σ be a random permutation of the elements of N and consider the random variable
  - $\mathsf{X} = |\{i: \{\sigma(1), \sigma(2), ..., \sigma(i)\} \in F\}|$
  - $\begin{array}{ll} \mathsf{F=}\{\;\{1,3\},\;\{2\}\} & \sigma=(2,3,1) \\ \mathsf{X=1} & (\mathsf{why?}) \end{array} \end{array}$
  - F={ {1,3}, {2}}  $\sigma = (3,2,1)$ X= 0 (why?)

 Considering the expectation of X, prove that

$$\mid F \mid \leq \binom{n}{\lfloor n/2 \rfloor}$$

#### [Hint]

Separate *F* by the size of elements. Number of size-1 set:  $K_1$ Number of size-2 set  $K_2$ , ..., Number of size-*n* set  $K_n$ .

$$|F| = K_1 + K_2 + \dots + K_n$$
  
E[X] = ?

Tournament

A complete oriented graph i.e., a graph in which every pair of nodes is connected by a single uniquely directed edge.



Show that there is a tournament *T* with *n* vertices which contains at least *n*! 2<sup>-(n-1)</sup> Hamiltonian paths.

#### [Hint] You can do it without any hint.

• Consider a graph in  $G_{n,p}$ , with p = 1/n. Let X be the # of triangles in the graph. Show that

(a)  $\Pr(X \ge 1) \le 1/6$ (b)  $\lim_{n \to \infty} \Pr(X \ge 1) \ge 1/7$ 

[Hint] For (b), use conditional expectation inequality.

• Use the general form of the Lovasz local lemma to prove that the symmetric version can be improved where we can replace the condition  $4dp \le 1$  by the weaker condition  $ep(d+1) \le 1$ .

#### [Hint]

- 1. Set  $x_i = 1/(d+1)$  to the general case Lovasz local lemma.
- 2.  $Pr(E_i) \leq p$  (symmetric version) Try to prove

$$\Pr(E_i) \leq x_i \prod_{(i,j) \in E} (1=x_j).$$