

CS5371

Theory of Computation

Lecture 2: Mathematics Review II
(Proof Techniques)

Some Updates

- Our consultation hours are as follows:
Kai (資電館741): Tue 1500—1600
Fri 1400—1500
Yu-Han (紅樓315): Wed 1500—1700
- There is a link from my homepage to access the course homepage:
www.cs.nthu.edu.tw/~wkhon

Objectives

- This time, we will look at some examples to demonstrate the following common proof techniques
 - By contradiction
 - By construction
 - By induction
- These techniques often occur in proving theorems in the theory of computation

By Contradiction

- One common way to prove a theorem is to assume that the theorem is false, and then show that this assumption leads to an obviously false consequence (also called a **contradiction**)
- This type of reasoning is used frequently in everyday life, as shown in the following example

By Contradiction

- Jack sees Jill, who just comes in from outdoor
- Jill looks completely dry
- Jack knows that it is not raining
- Jack's proof:
 - If it *were* raining (the assumption that the statement is false), Jill will be wet.
 - The consequence is: "Jill is wet" AND "Jill is dry", which is obviously false
 - Therefore, it must not be raining

By Contradiction [Example 1]

- Let us define a number is **rational** if it can be expressed as p/q where p and q are integers; if it cannot, then the number is called **irrational**
- E.g.,
 - 0.5 is rational because $0.5 = 1/2$
 - 2.375 is rational because $2.375 = 2375 / 1000$

By Contradiction

- Theorem: $\sqrt{2}$ (the square-root of 2) is irrational.
- How to prove?
- First thing is ...
 - Assume that $\sqrt{2}$ is rational

By Contradiction

- Proof: Assume that $\sqrt{2}$ is rational. Then, it can be written as p/q for some positive integers p and q .
- In fact, we can further restrict that p and q does not have common factor.
 - If D is a common factor of p and q , we use $p' = p/D$ and $q' = q/D$ so that $p'/q' = p/q = \sqrt{2}$ and there is no common factor between p' and q'
- Then, we have $p^2/q^2 = 2$, or $2q^2 = p^2$.

By Contradiction

- Since $2q^2$ is an even number, p^2 is also an even number
 - This implies that p is an even number (why?)
- So, $p = 2r$ for some integer r
- $2q^2 = p^2 = (2r)^2 = 4r^2$
 - This implies $2r^2 = q^2$
- So, q is an even number
- Something wrong happens... (what is it?)

By Contradiction

- We now have: "p and q does not have common factor" AND "p and q have common factor"
 - This is a contradiction
- Thus, the assumption is wrong, so that $\sqrt{2}$ is irrational

By Contradiction [Example 2]

- Theorem (Pigeonhole principle): A total of $n+1$ balls are put into n boxes. At least one box containing 2 or more balls.
- Proof: Assume "at least one box containing 2 or more balls" is false
 - That is, each has at most 1 or fewer ball

Consequence: total number of balls $\leq n$

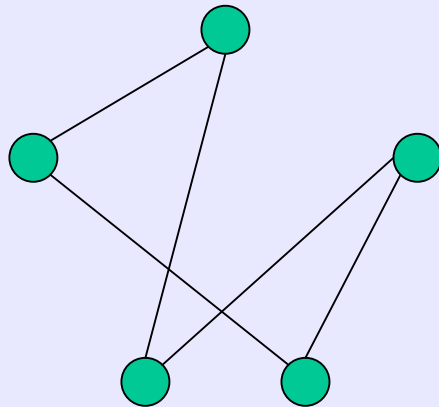
Thus, there is a contradiction (what is that?)

Proof By Construction

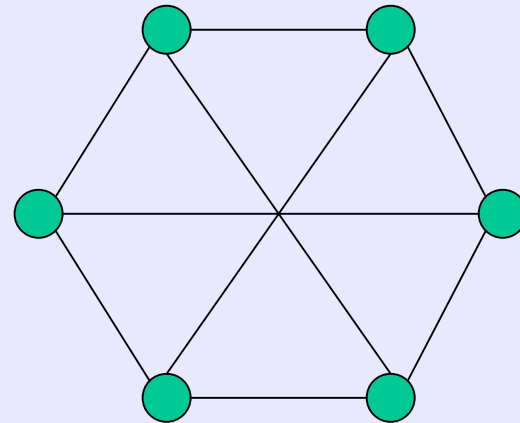
- Many theorem states that a particular type of object exists
- One way to prove is to find a way to construct one such object
- This technique is called **proof by construction**

By Construction [Example 1]

- Let us define a graph to be **k-regular** if every vertex of the graph has degree k
- E.g.,



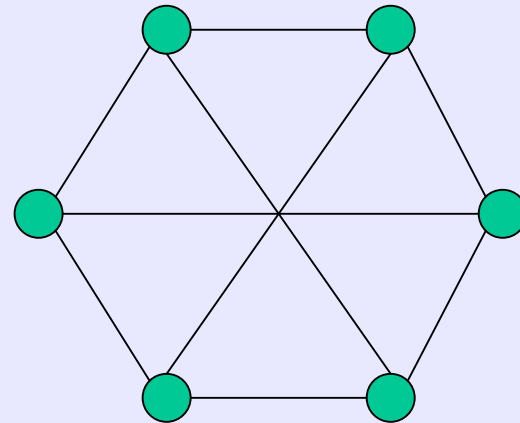
2-regular



3-regular

By Construction

- Theorem: For each even number $n \geq 4$, there exists a 3-regular graph with n vertices.
- How to prove it?



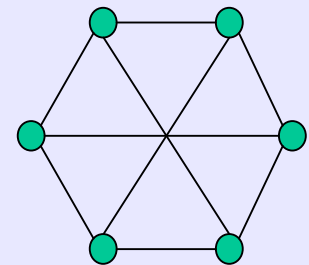
By Construction

- Proof Idea: Arrange the points evenly in a circle, for each vertex, form two edges one with its left neighbor and one with its right neighbor. Also, form an edge with the vertex opposite to it in the circle

- Formal Proof: Label the vertices by $1, 2, \dots, n$.

The edge set E is the union of

- $E_1 = \{ \{x, x+1\} \mid \text{for } x = 1, 2, \dots, n-1 \}$
- $E_2 = \{ \{1, n\} \}$
- $E_3 = \{ \{x, x + (n/2)\} \mid \text{for } x = 1, 2, \dots, n/2 \}$



Then, it is easy to check that the degree of each vertex is exactly 3.

By Construction [Example 2]

- Theorem: There exists a rational number p which can be expressed as q^r , with q and r both irrational.
- How to prove?
 - Find p, q, r satisfying the above condition
- What is the irrational number we just learnt? Can we make use of it?

By Construction

- What is the following value?

$$(\sqrt{2} \sqrt{2}) \sqrt{2}$$

- If $\sqrt{2} \sqrt{2}$ is rational, then $q = r = \sqrt{2}$ gives the desired answer
- Otherwise, $q = \sqrt{2} \sqrt{2}$ and $r = \sqrt{2}$ gives the desired answer

By Induction

- Normally used to show that all elements in an infinite set have a specified property
- The proof consists of proving two things: The **basis**, and the **inductive step**

By Induction

- To illustrate how induction works, let us consider the infinite set of natural numbers, $\{1,2,3,\dots\}$ and we want to show some property P holds for each element in the set
- One way to do so is:
 - Show P holds for 1 [shorthand: $P(1)$ is true]
 - Show for each $k \geq 1$, if $P(k)$ is true, then $P(k+1)$ is true [shorthand: $P(k) \rightarrow P(k+1)$ is true]

By Induction

- Then, we can conclude that $P(k)$ is true for all $k \geq 1$ (why?)
 - $P(1)$ is true
 - Because $P(1)$ is true and $P(k) \rightarrow P(k+1)$, then $P(2)$ is true
 - Because $P(2)$ is true and $P(k) \rightarrow P(k+1)$, then $P(3)$ is true
 - ...

By Induction

- There can be many other types of basis and inductive step, as long as by proving both of them, they can cover all the cases
- For example, to show P is true for all $k > 1$, we can show
 - Basis: $P(1)$ is true, $P(2)$ is true
 - Inductive step: $P(k) \rightarrow P(k+2)$
- Another example
 - Basis: $P(1)$ is true, $P(2)$ is true, ..., $P(2^i)$ is true for all i
 - Inductive step: $P(k) \rightarrow P(k-1)$

By Induction [Example 1]

- Let $F(k)$ be a sequence defined as follows:
- $F(1) = 1$
- $F(2) = 1$
- for all $k \geq 3$, $F(k) = F(k-1) + F(k-2)$
- Theorem: For all $n \geq 1$,
$$F(1)+F(2) + \dots + F(n) = F(n+2) - 1$$

By Induction

- Let $P(k)$ means "the theorem is true when $n = k$ "
- Basis: To show $P(1)$ is true.
 - $F(1) = 1, F(3) = F(1) + F(2) = 2$
 - Thus, $F(1) = F(3) - 1$
 - Thus, $P(1)$ is true
- Inductive Step: To show for $k \geq 1, P(k) \rightarrow P(k+1)$
 - $P(k)$ is true means: $F(1) + F(2) + \dots + F(k) = F(k+2) - 1$
 - Then, we have
$$\begin{aligned} & F(1) + F(2) + \dots + F(k+1) \\ &= (F(k+2) - 1) + F(k+1) \\ &= F(k+3) - 1 \end{aligned}$$
 - Thus, $P(k+1)$ is true if $P(k)$ is true

By Induction?

- CLAIM: In any set of h horses, all horses are of the same color.
- PROOF: By induction. Let $P(k)$ means "the claim is true when $h = k$ "
- Basis: $P(1)$ is true, because in any set of 1 horse, all horses clearly are the same color.

By Induction?

- Inductive step:
 - Assume $P(k)$ is true.
 - Then we take any set of $k+1$ horses.
 - Remove one of them. Then, the remaining horses are of the same color (because $P(k)$ is true).
 - Put back the removed horse into the set, and remove another horse
 - In this new set, all horses are of same color (because $P(k)$ is true).
 - Therefore, all horses are of the same color!
- What's wrong?

More on Pigeonhole Principle

- Theorem: For any graph with more than two vertices, there exists two vertices whose degree are the same.
- How to prove?

More on Pigeonhole Principle

- Theorem: There exists a number consisted by all 1's (such as 1, 11, 111, ...) which is divisible by 1997.
- How to prove?

Next

- Part I: Automata Theory