

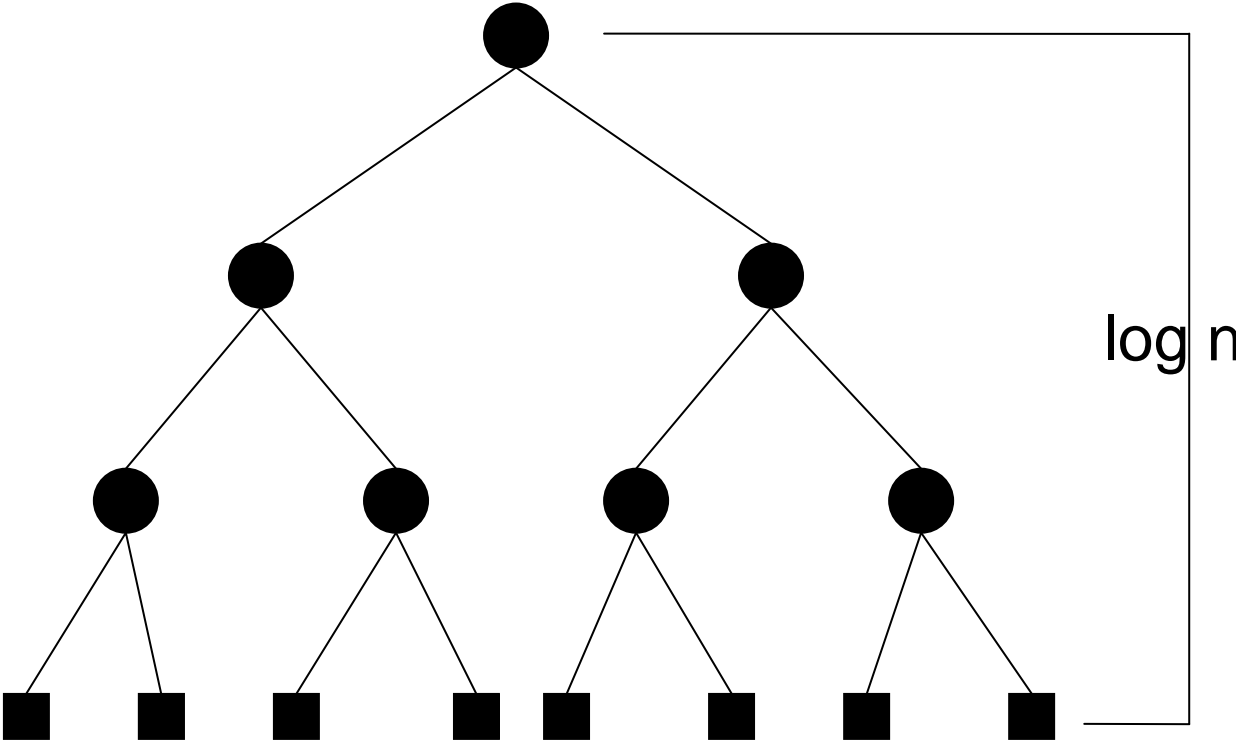
# Red Black Tree

A balanced binary search tree

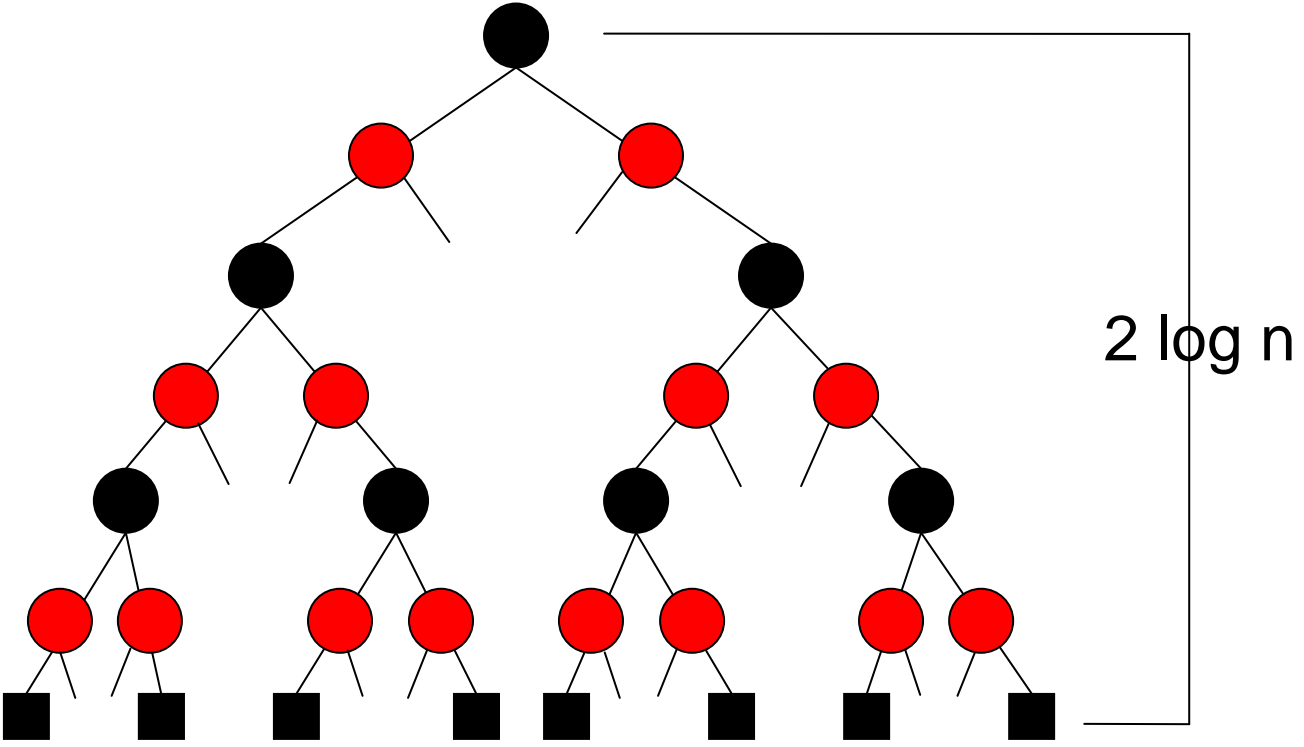
# Red Black Tree

1. Every node is either red or black
2. For each node, all paths from the node to descendant leaves contain the same number of black nodes
3. If a node is red, then both its children are black
4. The root is black
5. Every dummy leaf is black

# Balance

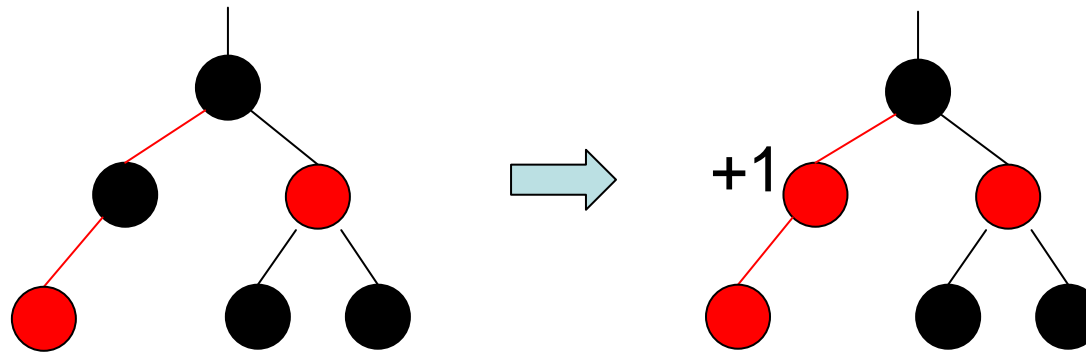


# Balance



# Notation

- +1 means need one more black for the node



# Insertion

- Set inserted node to be red
- Fix the violation of Properties 3 and 4

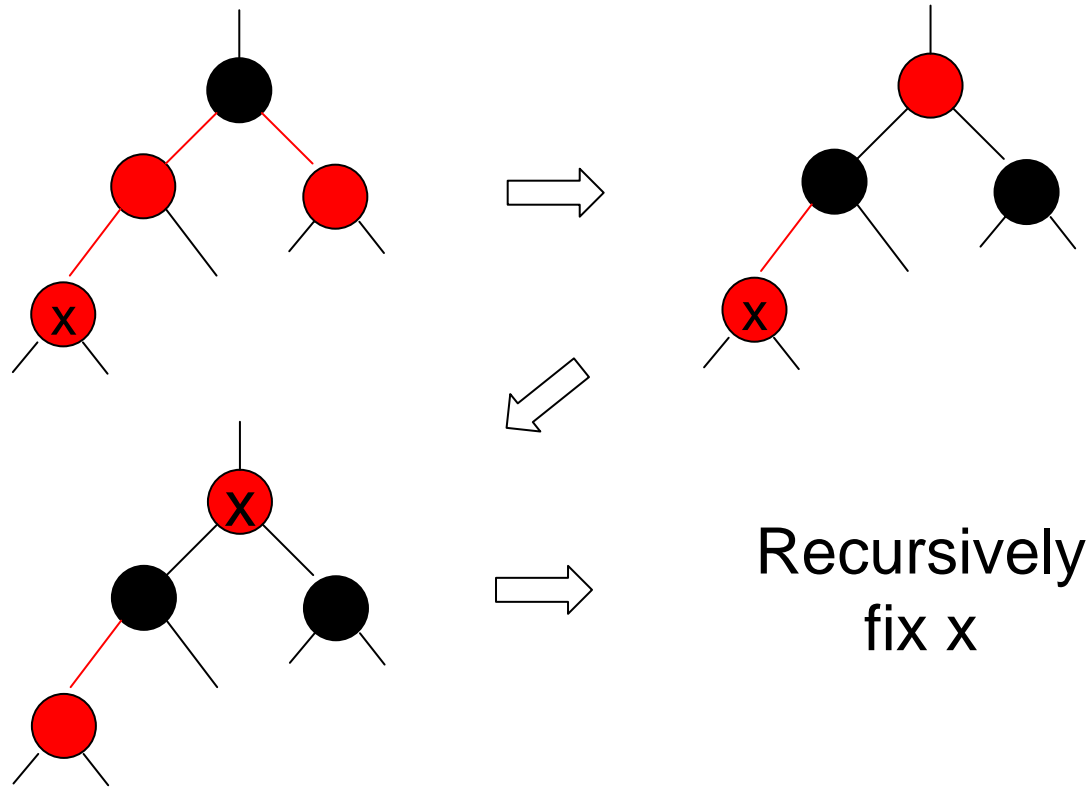
# Insertion

To insert x:

- Uncle is red (Case 1)
- Uncle is black
  - Both x and parent are left (or right) child (Case 2)
  - Others (Case 3)

# Insertion

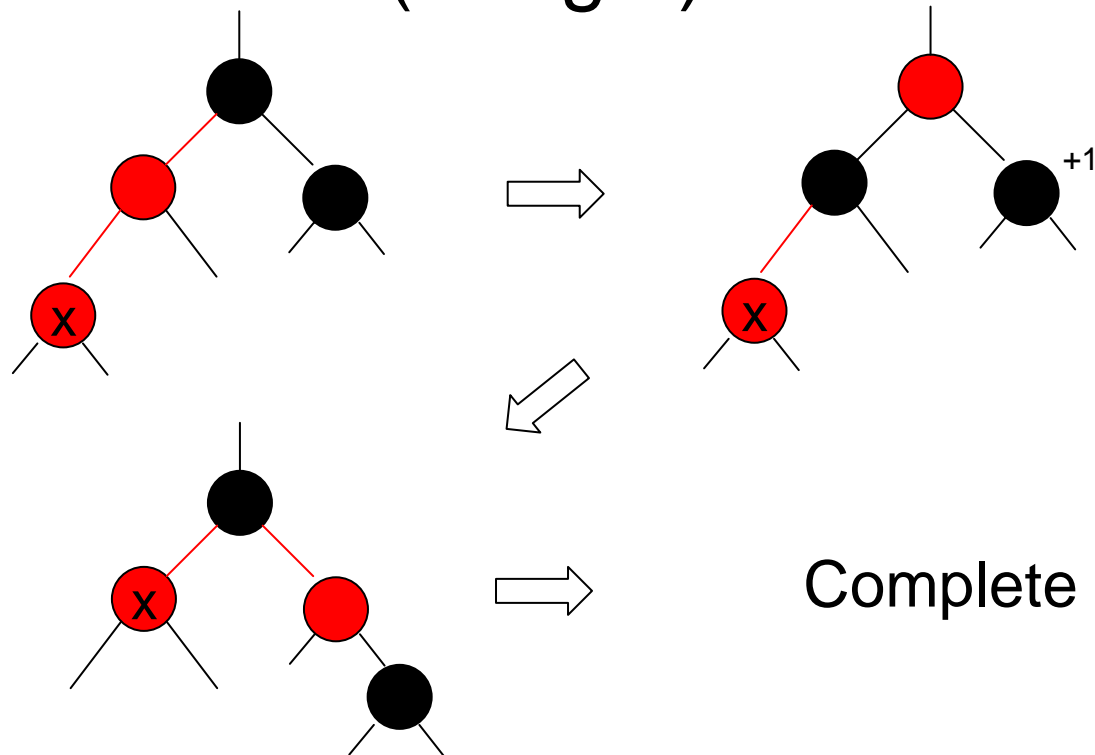
- Case 1: Uncle is red





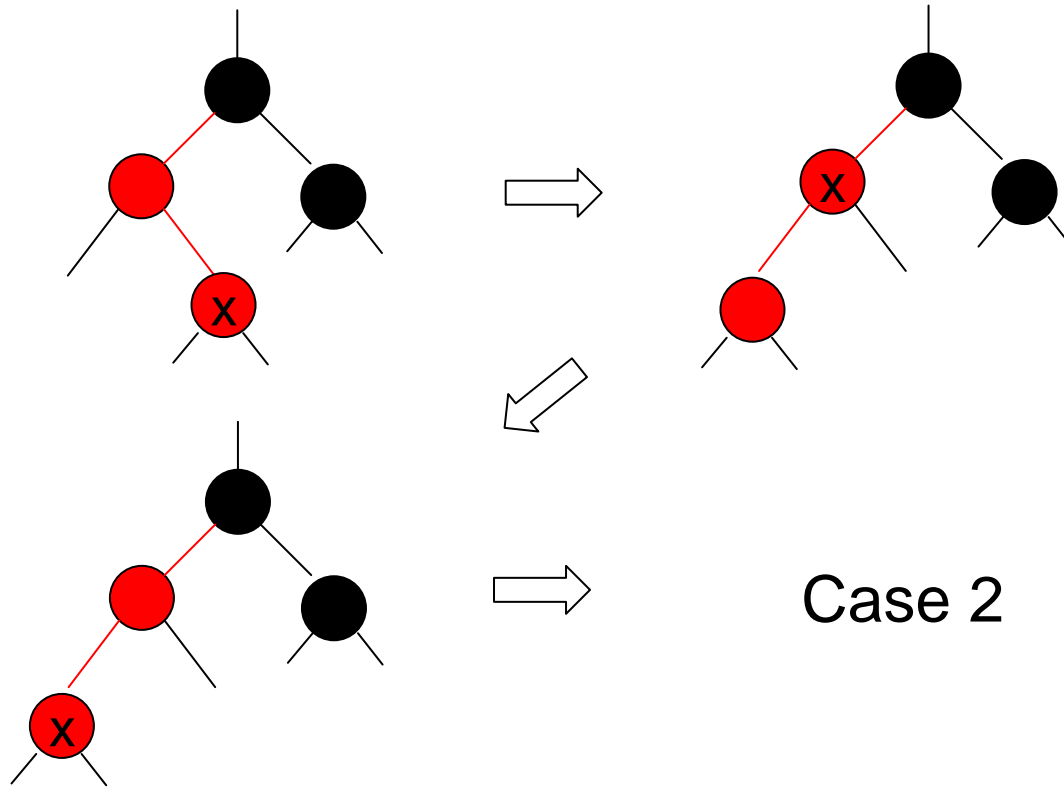
# Insertion

- Case 2: Uncle is black. Both x and parent are left (or right) child



# Insertion

- Case 3: Others



# Delete vs Remove

- To delete  $z$ , node  $z$  may not be removed in the tree
- Denote  $y$  as the removed node
- Let  $x$  be the child of  $y$

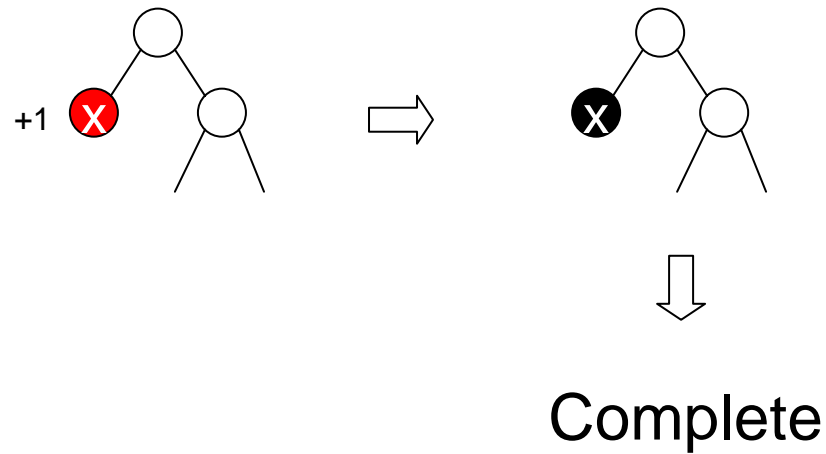
# Deletion

- No violation when the removed node  $y$  is red
- Otherwise, fix the violation of property 2 and 4

- x is red (Case 1)
- x is black
  - Sibling is red (Case 2)
  - Sibling, denoted as s, is black
    - Both s's children are black (Case 3)
    - The children of s, left are black, right are red (Case 4)
    - Right children of s are red
      - Parent are black (Case 5)
      - Parent are red (Case 6)

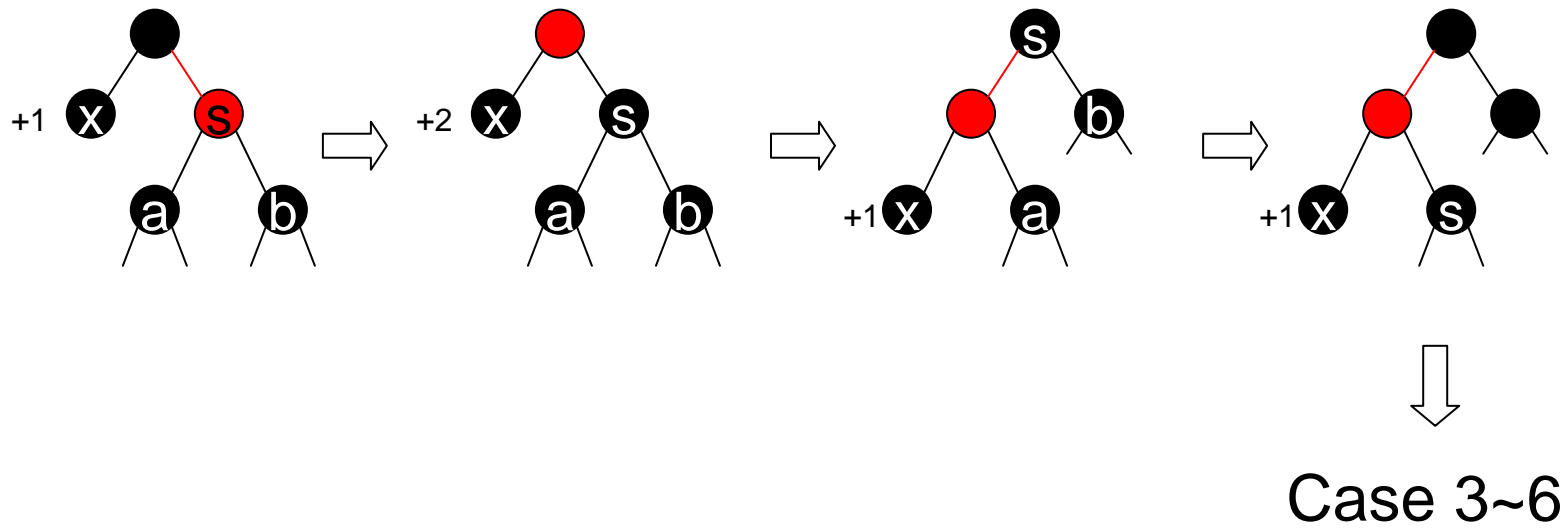
# Deletion

- Case 1: x is red



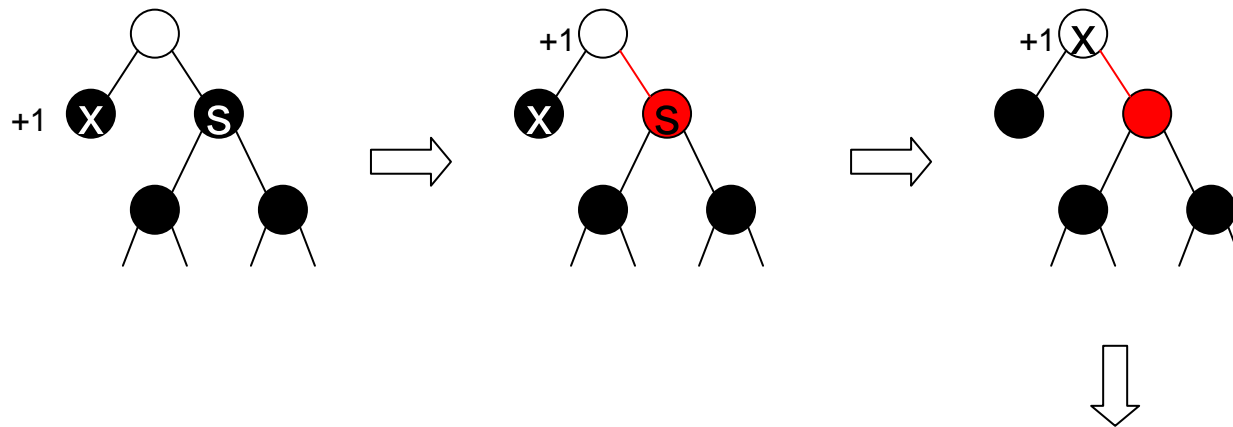
# Deletion

- Case 2: Sibling is red



# Deletion

- Case 3: Both sibling's children are black

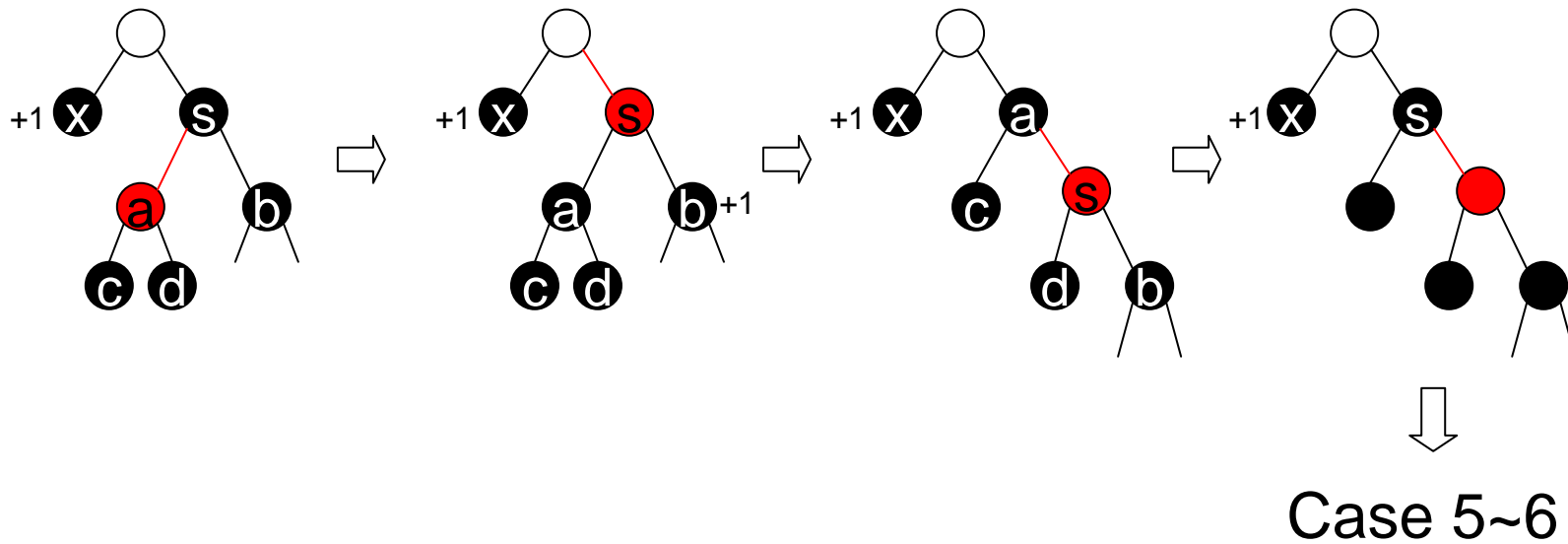


Recursively fix x



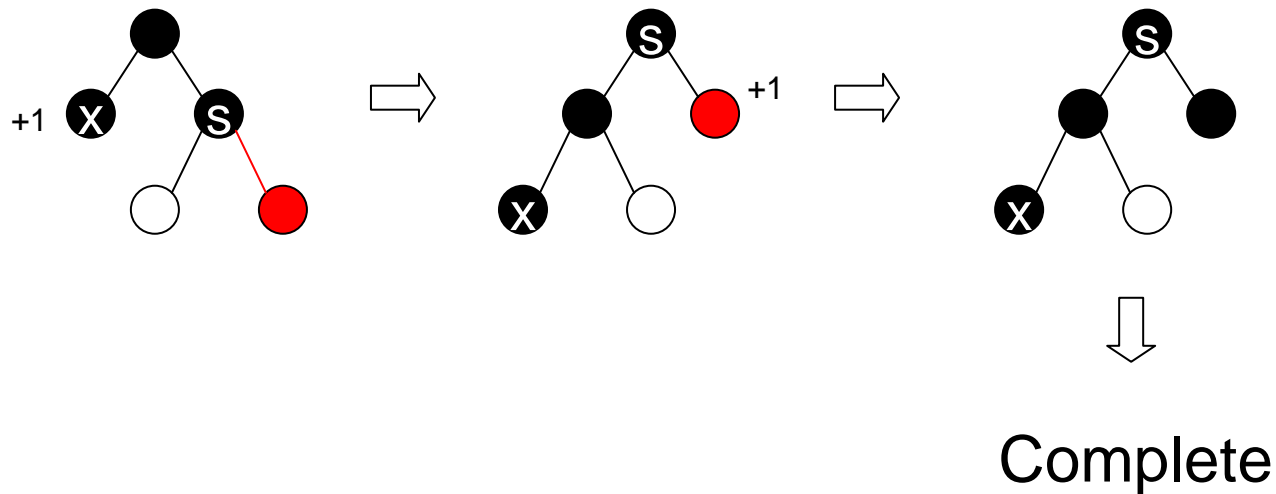
# Deletion

- Case 4: Sibling's left child is red, right child is black



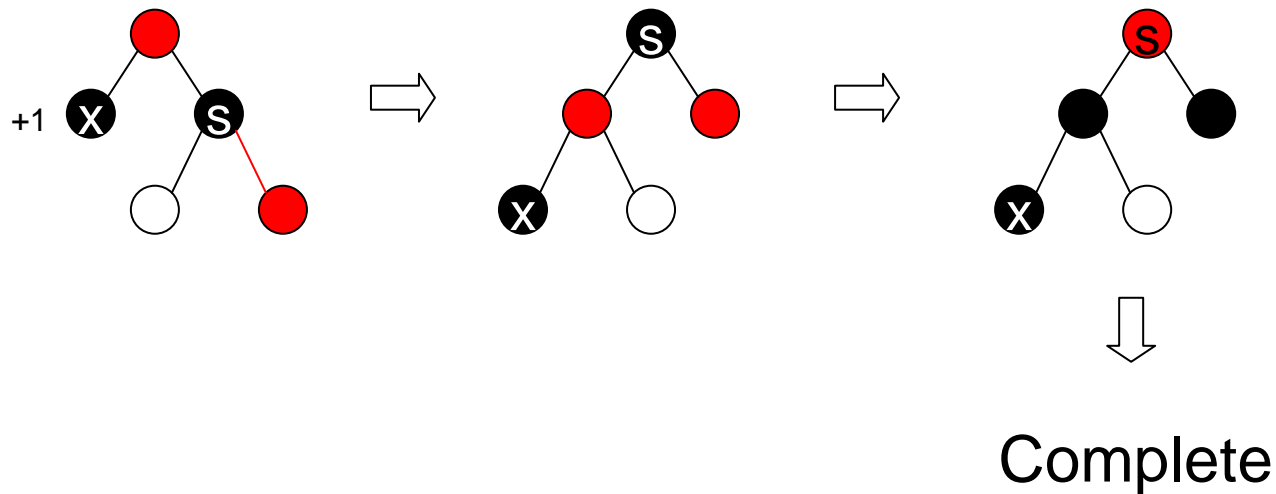
# Deletion

- Case 5: Sibling's right child is red



# Deletion

- Case 6: Sibling's right child is red



# 2 colors, why?

- Tree height is not that tight
- Reduce the cost to balance tree