CS 5319 Advanced Discrete Structure

Lecture 5: Recurrence Relations I

Outline

- Introduction
- Linear Recurrence Relations with Constant Coefficients
- Solution by Generating Functions
- ***** Special Non-Linear Recurrences
- Recurrence with Two Indices

This Lecture

- Consider the sequence (1, 3, 9, 27, 81, ...)
- We can describe the sequence by :

$$a_n = 3^n$$

Alternatively, we can describe the sequence by expressing the *n*th term with the (*n* – 1)st term, together with the specification of the 1st term :

$$a_n = 3a_{n-1}$$
, $a_0 = 1$

- For a sequence (a₀, a₁, a₂, ...), an equation that relates a_n to some of its predecessor is called :
 recurrence relation or difference equation
- To start the computation, one must know several numbers in the sequence. These numbers are called the boundary conditions
- Ex : $a_n = 3a_{n-1} \iff$ recurrence relation $a_0 = 1 \iff$ boundary condition

- Ex: Consider the Fibonacci sequence (1, 1, 2, 3, 5, 8, 13, 21, ...)
- The sequence can be described by the recurrence relation:

$$a_n = a_{n-1} + a_{n-2}$$

boundary conditions:

$$a_0 = 1$$
, $a_1 = 1$

- Ex: Let there be *n* ovals in the plane.
 Suppose that any two ovals intersect at exactly two points, and no three ovals intersect at the same point
 - Q: How many regions do they divide the plane ?



- Let a_n = number of regions divided by *n* ovals
- Suppose we have already drawn n 1 ovals
- The *n*th oval will be divided into 2(n-1) arcs

→
$$a_n = a_{n-1} + 2(n-1)$$

• Together with the boundary condition $a_1 = 2$, we can further compute

$$a_2 = 4$$
, $a_3 = 8$, $a_4 = 14$, $a_5 = 22$, $a_6 = 32$, ...

- In many cases obtaining the recurrence is a big step towards the solution to a counting problem
 - This is because even if we do not know the general expression, we can still compute the desired term a_n by a step-by-step approach
- Of course it is better if we can really obtain a general expression for a_n; in the following, we shall discuss a few ways of doing so

Linear Recurrence Relations with Constant Coefficients

• A recurrence relation of the form :

$$C_0 a_n + C_1 a_{n-1} + \ldots + C_r a_{n-r} = f(n)$$

is called a linear recurrence relation with constant coefficients, where all *C*'s are constants

• Ex :
$$3a_n - 5a_{n-1} + 2a_{n-2} = n^2 + 5$$

- If *r* consecutive terms, say $a_{k+1}, a_{k+2}, \dots, a_{k+r}$, are known, then any a_n can be calculated recursively
- This implies that :
 - Solution to the linear recurrence is determined uniquely by these *r* boundary conditions
- Indeed, we shall see that :

General form of the solution (when boundary conditions not fixed) has *r* unknown constants

- Let us begin with a simple example
- Suppose we know that

$$a_n + a_{n-1} = 2n + 1$$

• By observation, we may see that

 $a_n = n + 1$

satisfies the above relation

• However, if we consider a similar recurrence

$$a_n + a_{n-1} = 0$$

we may notice that for any constant A $a_n = A \times (-1)^n$

satisfies the above relation

• Consequently, we see that for any *A*,

$$a_n = n + 1 + A \times (-1)^n$$

will satisfy the original recurrence

$$a_n + a_{n-1} = 2n + 1$$

- If somehow we know the boundary condition, we can determine the value of *A* exactly
- Ex : If $a_0 = 8$, then A = 7

- Let us review what we have done
- In the above computation, we first find out a particular case where

$$a_n + a_{n-1} = 2n + 1$$

to obtain

$$a_n = n + 1$$

• This is called a particular solution to the recurrence

• Next, we set the right side of the recurrence to 0.

$$a_n + a_{n-1} = 0$$

By solving this new recurrence, we obtain

$$a_n = A \times (-1)^n$$

• This is called the homogeneous solution to the recurrence

• The particular solution and the homogeneous solution together allow us to obtain a general (total) solution to the original recurrence :

$$a_n = n + 1 + A \times (-1)^n$$

• Finally, we resolve the unknowns based on the boundary conditions, and obtain the exact solution to the recurrence relation

- The particular solution is usually found by clever "guessing"
- Ex : To solve

$$a_n + 2a_{n-1} = n+3$$

a reasonable guess is that

 $a_n = Bn + D$

for some constants B and D

• Ex : To solve

$$a_n - a_{n-1} = n + 3$$

a reasonable guess is that

$$a_n = Bn^2 + Dn + E$$

for some constants B, D, and E

• Ex : To solve

$$a_n + 2a_{n-1} + a_{n-2} = 2^n$$

a reasonable guess is that

$$a_n = B \star 2^n$$

for some constant B

- The homogeneous solution can be found by a systematic way
- To solve

$$C_0 a_n + C_1 a_{n-1} + \ldots + C_r a_{n-r} = 0$$

we first create a characteristic equation :

$$C_0 x^r + C_1 x^{r-1} + \ldots + C_r = 0$$

• Suppose α is a root to the characteristic equation. Then we see that for any constant *A*

 $a_n = A\alpha^n$

is a homogeneous solution, because

 $C_0(A\alpha^n) + C_1(A\alpha^{n-1}) + \ldots + C_r(A\alpha^{n-r}) = 0$

- Extending the previous arguments, if the characteristic equation has *r* distinct roots, denoted by α₁, α₂, ..., α_r.
- Then we see that for any $A_1, A_2, ..., A_r$

$$a_n = A_1 \alpha_1^n + A_2 \alpha_2^n + \ldots + A_r \alpha_r^n$$

is a homogeneous solution

• Ex : The recurrence relation for Fibonacci sequence is :

$$a_n = a_{n-1} + a_{n-2}$$

This gives the characteristic equation :

$$x^2 - x - 1 = 0$$

• The two roots of the characteristic equation are :

$$\alpha = \frac{1 + \sqrt{5}}{2}$$
, $\beta = \frac{1 - \sqrt{5}}{2}$

Thus the homogeneous solution is :

$$a_n = A \alpha^n + B \beta^n$$

This is also the total solution since the particular solution is 0

• Because $a_0 = 1$ and $a_1 = 1$, we have

$$1 = A + B$$
, $1 = A\alpha + B\beta$

Thus, we can solve A and B, and finally obtain

$$a_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^{n+1}$$

• Ex : Evaluate the $n \times n$ determinant

1	1	0	0	0	•••	0	0	0	0	0
1	1	1	0	0	•••	0	0	0	0	0
0	1	1	1	0	•••	0	0	0	0	0
0	0	1	1	1	•••	0	0	0	0	0
••	••	••	••	••	•••	••	••	••	••	••
 0	 0	 0	 0	 0	•••	 0	 1	 1	 1	 0
 0 0	 0 0	 0 0	 0 0	 0 0	•••	 0 0	 1 0	 1 1	 1 1	 0 1

- Let a_n = the value of $n \times n$ determinant
- By expanding the determinant with respect to the first column, we get

$$a_{n} = a_{n-1} - \begin{vmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & \dots & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & \dots & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \dots & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & \dots & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & \dots & 0 & 0 & 0 & 1 & 1 \end{vmatrix} \right\} n-1$$

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• Next by expanding the remaining determinant with respect to the first row, we get

$$a_n = a_{n-1} - a_{n-2}$$

This gives the characteristic equation :

$$x^2 - x + 1 = 0$$

• The two roots of the characteristic equation are :

$$\alpha = \frac{1 + \sqrt{3}i}{2} , \quad \beta = \frac{1 - \sqrt{3}i}{2}$$

Thus the homogeneous solution is :

$$a_n = A \alpha^n + B \beta^n$$

This is also the total solution since the particular solution is 0

• Since we are dealing with complex numbers, it is more convenient to express it in polar form :

$$\alpha = \cos\frac{\pi}{3} + i\sin\frac{\pi}{3}$$
, $\beta = \cos\frac{\pi}{3} - i\sin\frac{\pi}{3}$

Thus the homogeneous solution becomes :

$$A\left(\cos\frac{n\pi}{3}+i\sin\frac{n\pi}{3}\right)+B\left(\cos\frac{n\pi}{3}-i\sin\frac{n\pi}{3}\right)$$

Because a₁ = 1 and a₂ = 0, we can solve A and B, and finally obtain

$$a_n = \frac{1}{2} \cos \frac{n\pi}{3} + \frac{1}{\sqrt{3}} \sin \frac{n\pi}{3}$$

- Observe: There is no *i* in the expression for a_n
 - Reason: All coefficients in the recurrence, and the boundary conditions, are real
- Thus we may at the beginning assume that

$$a_n = D \cos \frac{n\pi}{3} + E \sin \frac{n\pi}{3}$$

for some real numbers D and E, and solve for them

- So far, we have only considered the case where the characteristic equation has distinct roots.
- What if a root α is a double root ?
- We shall see that apart from $a_n = A\alpha^n$,

 $a_n = B n \alpha^n$

is also a homogeneous solution

• First, we observe that α is a root to

$$C_0 n x^{n-1} + C_1 (n-1) x^{n-2} + \ldots + C_r (n-r) x^{n-r-1} = 0$$

so that

$$C_0 n \alpha^{n-1} + C_1 (n-1) \alpha^{n-2} + \ldots + C_r (n-r) \alpha^{n-r-1} = 0$$

• Consequently, this implies $a_n = B n \alpha^{n-1}$ (or $B n \alpha^n$) is a homogeneous solution

• Extending the previous arguments, if α is a *k*-multiple root of the characteristic equation, then

$$a_n = (A_1 n^k + A_2 n^{k-1} + \dots + A_k) \alpha^n$$

is a homogeneous solution

• Further, if $\alpha_1, \alpha_2, \dots$ are all roots of the equation such that α_i is a k_i -multiple root, then we have

$$k_1 + k_2 + \ldots = r ,$$

and the homogenous solution has *r* unknowns which can be found by boundary conditions

• Ex : To solve

$$a_n + 6a_{n-1} + 12a_{n-2} + 8a_{n-3} = 0$$

• The homogeneous solution is :

$$(A_1 n^2 + A_2 n + A_3) (-2)^n$$

so that we can solve for A_1, A_2, A_3 by boundary conditions

• Ex : Evaluate the $n \times n$ determinant

2	1	0	0	0	•••	0	0	0	0	0
1	2	1	0	0	•••	0	0	0	0	0
0	1	2	1	0	•••	0	0	0	0	0
0	0	1	2	1	•••	0	0	0	0	0
••	••	••	••	••	•••	••	••	••	••	••
 0	 0	 0	 0	 0	•••	 0	 1	 2	 1	 0
 0 0	 0 0	 0 0	 0 0	 0 0	•••	 0 0	 1 0	 2 1	 1 2	 0 1

• Ex : The Tower of Hanoi

We are given a tower of *n* discs, initially stacked in decreasing size on one of the 3 pegs :



- Objective : Move entire tower to another peg
- Restrictions :
 - 1. Move only one disc at a time, and
 - 2. Never move a large one onto a smaller one
- Question : What is the minimum # of moves ?



Solution by Generating Functions

• Let us revisit an old problem

Consider *n* ovals in the plane, where any two ovals intersect at exactly two points, and no three ovals intersect at the same point

Q: How many regions do they divide the plane ?



- Let a_n = number of regions divided by *n* ovals
- Then we have $a_1 = 2$, and for any n > 1

$$a_n = a_{n-1} + 2(n-1)$$

• Now, let A(x) be the GF of the values of a_n 's:

$$A(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots$$

and for consistency, we assume $a_0 = 2$

• Next, we use of A(x) to obtain a formula for a_n

• Firstly, we have

$$a_n x^n = a_{n-1} x^n + 2(n-1) x^n$$

which implies

$$\sum_{n=1}^{\infty} a_n x^n = \sum_{n=1}^{\infty} a_{n-1} x^n + \sum_{n=1}^{\infty} 2(n-1) x^n$$

• Thus we have :

$$A(x) - a_0 = x A(x) + 2x^2 / (1 - x)^2$$

• Next by rearranging terms, we get

$$A(x) = 2 / (1-x) + 2x^2 / (1-x)^3$$

• It follows that

$$a_n = 2 + 2 \times C(-3, n-2) (-1)^{n-2}$$

= 2 + 2 × n (n - 1) / 2
= 2 + n (n - 1)

• In the previous derivation, we have set $a_0 = 2$ so that the formula

$$a_n = a_{n-1} + 2(n-1)$$

can be applied even for n = 1

- Indeed, we can set a_0 to any arbitrary constant and the above method for solving a_n still work
 - But we will need some adjustments since the formula is no longer valid for *n* = 1

- Let us set $a_0 = 5$ (an arbitrary constant)
- Then we have :

$$\sum_{n=2}^{\infty} a_n x^n = \sum_{n=2}^{\infty} a_{n-1} x^n + \sum_{n=2}^{\infty} 2(n-1) x^n$$

• Thus we have :

$$A(x) - a_1 x - a_0 = x \left(A(x) - a_0 \right) + 2x^2 / (1 - x)^2$$

• Then by rearranging terms, we get

$$A(x) = 2x / (1-x) + 2x^2 / (1-x)^3 + 5$$

• It follows that

$$a_n = \begin{cases} 5 & n = 0 \\ n(n-1) + 2 & n = 1, 2, 3, \dots \end{cases}$$

• We next revisit another old problem

Consider all *n*-digit quaternary strings. How many of them contains even # of 0's?

- Let $a_n = \#$ of *n*-digit strings with even 0's
- Previously, we have used combinatorial arguments or exponential GF to find a_n
- We now find a_n by first deriving a recurrence relation for a_n , and solve the relation using GF

- Firstly, an *n*-digit string with even 0's can be obtained in one of the following ways :
 - 1. Obtain an (n 1)-digit string with even 0's, and append either 1, 2, 3 at its end ;
 - 2. Obtain an (n 1)-digit string with odd 0's, and append 0 at its end
- Thus for n > 1,

$$a_n = 3a_{n-1} + (4^{n-1} - a_{n-1})$$

• After simplification, we get

$$a_n - 2a_{n-1} = 4^{n-1}$$

• Since $a_1 = 3$, we set $a_0 = 1$ so that the recurrence also holds for n = 1, and thus obtain :

$$\sum_{n=1}^{\infty} a_n x^n - \sum_{n=1}^{\infty} 2a_{n-1} x^n = \sum_{n=1}^{\infty} 4^{n-1} x^n$$

• Consequently, we have

$$A(x) - 1 - 2x A(x) = x / (1 - 4x)$$

• By rearranging terms, we get

$$A(x) = \frac{1}{1 - 2x} \left(1 + \frac{x}{1 - 4x} \right)$$
$$= \frac{1/2}{1 - 2x} + \frac{1/2}{1 - 4x}$$

• This immediately shows that

$$a_n = \frac{1}{2} 2^n + \frac{1}{2} 4^n$$

• A challenging problem

Consider all *n*-digit quaternary strings. How many of them contains even # of 0's and even # of 1's?

Let b_n = # of n-strings with even 0's and even 1's
 c_n = # of n-strings with even 0's and odd 1's
 d_n = # of n-strings with odd 0's and even 1's

• Thus for n > 1, we have :

$$b_{n} = 2b_{n-1} + c_{n-1} + d_{n-1}$$

$$c_{n} = b_{n-1} + 2c_{n-1} + 4^{n-1} - b_{n-1} - c_{n-1} - d_{n-1}$$

$$d_{n} = b_{n-1} + 2d_{n-1} + 4^{n-1} - b_{n-1} - c_{n-1} - d_{n-1}$$

• After simplification, we have :

$$b_n = 2b_{n-1} + c_{n-1} + d_{n-1}$$

$$c_n = c_{n-1} - d_{n-1} + 4^{n-1}$$

$$d_n = d_{n-1} - c_{n-1} + 4^{n-1}$$

• Since we can choose b_0 , c_0 , d_0 arbitrarily without affecting the result, we shall set

$$b_0 = \frac{3}{4}$$
, $c_0 = \frac{1}{4}$, $d_0 = \frac{1}{4}$

so that the previous recurrences is also valid for n = 1 (In fact, there are other sets of b_0, c_0, d_0)

- Next, we shall multiply both sides by x^n and obtain the sum for all $n \ge 1$
 - That is, we sum all the valid cases

• Consequently, we obtain :

$$B(x) - \frac{3}{4} = 2x B(x) + x C(x) + x D(x)$$

$$C(x) - \frac{1}{4} = x C(x) - x D(x) + \frac{x}{1 - 4x}$$

$$D(x) - \frac{1}{4} = x D(x) - x C(x) + \frac{x}{1 - 4x}$$

• Solving the above, we get :

$$C(x) = D(x) = \frac{1}{4} / (1 - 4x)$$

$$B(x) = \frac{1}{4} / (1 - 4x) + \frac{1}{2} / (1 - 2x)$$

$$b_n = \frac{1}{4} \frac{4^n}{4^n} + \frac{1}{2} \frac{2^n}{4^n}$$