CS4311 Design and Analysis of Algorithms

Tutorial: Assignment 1

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Outline

- What's a good answer
 - A sample question and solution

- Hint for the assignment 1
 - Due date: 3/12

Sample Question

Question:

Given a sequence of numbers, design an algorithm to find the maximum number. Show your algorithm is correct and analyze the time complexity.

How to write the answer?

Bad Answer

· Algorithm:

```
public class FindMax {
00
         public static void main(String[] args) {
01
              int data[] = new int[] {17, 26, 14, 19, 10, 23, 12, 5};
02
              int count = 0;
03
              int max = Integer.MIN_VALUE;
04
              for(int i=0; i<data.length; i++) {
05
                  if(data[i] > max)
                       max = data[i];
06
                  count++;
07
08
              System.out.println("Max Number: " + max);
09
              System.out.println("Count: " + count);
10
11
```

Bad Answer

C:\>java FindMax

Max Number: 26

Count: 8

C:\>

· Correctness:

26 is actually the maximum of the input data, so the algorithm is correct.

Time complexity:

Since count = 8 is equal to the number of input data, therefore, the running time is O(n).

Good Answer

·Algorithm:

- Given a sequence of number,
- Step 1: Pick up the first number and suppose it is maximum.
- Step 2: Pick up a number from the remaining numbers.
- Step 3: Compare the picked number and current maximum.
- Step 4: Throw out the smaller and suppose the larger one as the new maximum.
- Step 5: Repeat 2~4 until there's no number remaining.
- Step 6: Output the current maximum.

Good Answer

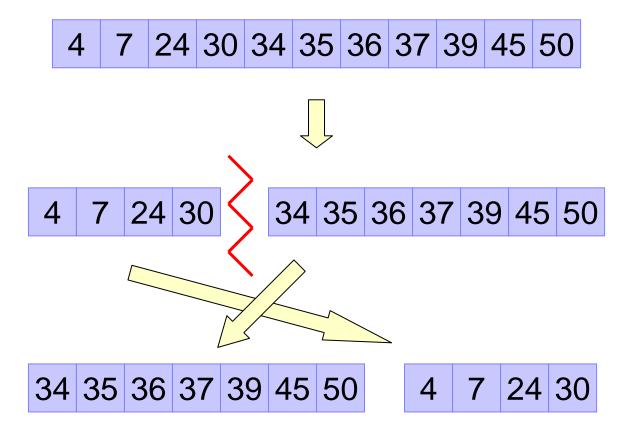
· Correctness:

Suppose our algorithm doesn't output the maximum number m, then m must be thrown out. It means m is smaller than some number. Since some number is greater than m, m is not maximum. We got a contradiction, therefore, our algorithm will output the maximum number.

Good Answer

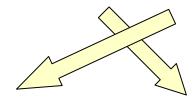
· Time complexity:

The operations involved are picking or throwing out numbers. Since each number is only picked once and thrown out once, and there are totally n numbers, the running time is n * O(1) = O(n).







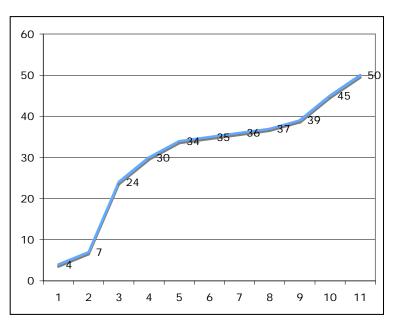


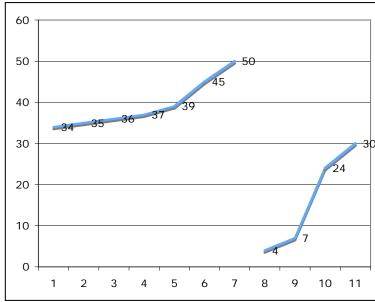
4 7 24 30 34 35 36 37 39 45 50

 Design an O(log n)-time algorithm to find the minimum item.

Show that your algorithm is correct.

· Hint:





ComputeCount()

- 1. Input a positive integer *n*
- 2. Set count = 0
- 3. **for** j = 1, 2, ..., n4.
- 4. **if** j is a factor of n
- { Update count to be faster?
- Output count

What's the running time?

 $\Theta(n)$

Can it be

- Design a faster algorithm that can compute count.
 - $O(\sqrt{n})$, $O(\log n)$, O(1)...
 - Note: You can only use RAM operations.
 - E.g., you cannot assume 2^x or $\log x$ can be computed in O(1) time.
- Explain why your algorithm is correct.

· Hint:

What's the physical meaning of count?

```
BubbleSort(A)

1. for Phase k = 1, 2, ..., n

2. for Position j = 1, 2, ..., n - 1

3. if A[j] > A[j + 1]

4. { Swap the entries A[j] and A[j + 1] }
```

Show the algorithm is correct.

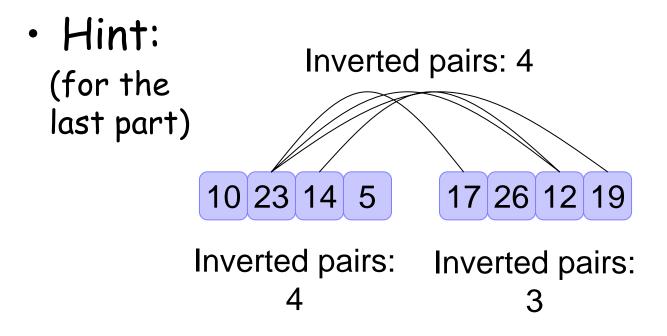
Inverted Pair:

Given $(A_0, A_1, A_2, ..., A_n)$ if i < j and $A_i > A_j$, then (A_i, A_i) is inverted.

Example:

Given (2, 3, 6, 4, 0)(2, 0), (3, 0), (6, 4), (6, 0), (4, 0) are inverted.

- Show that the number of inverted pairs is exactly equal to the number of swaps when we perform BubbleSort.
- Design an algorithm that counts the number of inverted pairs in O(n log n) time.



Total inverted pairs = 4 + 4 + 3 = 11

Question 4 (No marks)

Give asymptotic upper bound for each T(n)

a)
$$T(n) = 9T(n/2) + n^3$$

b)
$$T(n) = 7T(n/2) + n^3$$

c)
$$T(n) = T(\sqrt{n}) + \log n$$

d)
$$T(n) = 0.5T(n/2) + n$$

e)
$$T(n) = 3T(n/3) + n/3$$

Good Luck

 Please try to write your answer in large font. Thanks a lot!