

# CS4311 DESIGN AND ANALYSIS OF ALGORITHMS

## Homework 1

Due: 11:10 am, March 12, 2009 (before class)

1. You have just finished sorting an array  $A[1..n]$  of  $n$  distinct numbers into increasing order. When you go out to have a break, your mischievous friend, John, has divided your array into two parts  $A_{left} = A[1..i]$  and  $A_{right} = A[i + 1..n]$ , and re-arrange the array so that he puts  $A_{right}$  in front of  $A_{left}$ ; precisely, the array now becomes  $A[i + 1..n]A[1..i]$ . See Figure 1 for an example.

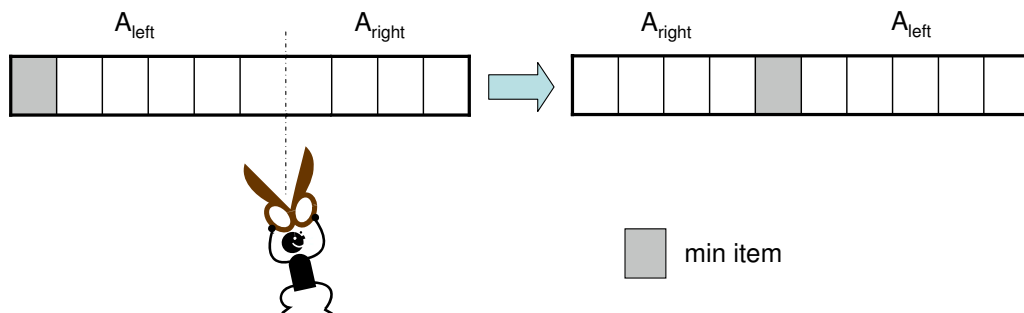


Figure 1: John's modification to the array.

After you come back, John tells you about what he has done, but without telling you the value of  $i$ . To reverse the change, you want to locate the entry with the minimum item, as this will be the boundary between  $A_{right}$  and  $A_{left}$ .

- (a) (15%) Design an  $O(\log n)$ -time algorithm to find the position of the minimum item.
  - (b) (15%) Show that your algorithm is correct.
2. Consider the following code `ComputeCount`:

```
ComputeCount()
1. Input a positive integer  $n$ ;
2. Set count = 0;
2. for  $j = 1, 2, \dots, n$ 
3.   if  $j$  is a factor of  $n$ 
4.     { Update count to become  $1 - \text{count}$ ; }
5. Output count;
```

- (a) (15%) The above code computes the value of `count` in  $\Theta(n)$  time. Design a faster algorithm that can compute `count`, and analyze its running time.
  - For this problem, the marks will also depend on the quality of your algorithm. At most 15% if your algorithm runs in  $O(\log n)$  time; otherwise, at most 5% if it runs in  $o(n)$  time, and 0% if it runs in  $\Theta(n)$  time.
- (b) (15%) Explain why your algorithm is correct.

3. The `BubbleSort` algorithm is a very simple algorithm for sorting an array of numbers. Given an input array  $A[1..n]$  with  $n$  distinct numbers, `BubbleSort` works by repeatedly swapping adjacent elements in  $A$  as follows:

```

BubbleSort(A)
1. for Phase  $k = 1, 2, \dots, n$ 
2.   for Position  $j = 1, 2, \dots, n - 1$ 
3.     if  $A[j] > A[j + 1]$ 
4.       { Swap the entries  $A[j]$  and  $A[j + 1]$ ; }

```

- (a) (15%) Show that `BubbleSort` is correct.
- (b) Consider the original array  $A[1..n]$ . We say a pair  $(A[i], A[j])$  is *inverted* if  $i < j$  and  $A[i] > A[j]$ . Intuitively,  $A[i]$  should be on the right of  $A[j]$  when the array is sorted, but it is currently on the left of  $A[j]$ .

- For example, if the array is  $\langle 2, 3, 6, 4, 0 \rangle$ , then the pair  $(3, 0)$  is inverted, and in total there are 5 inverted pairs.

(15%) Show that the number of inverted pairs in  $A$  is exactly equal to the number of swaps when we perform `BubbleSort` on  $A$ .

- \*\* (c) (10%) By using brute force approach, one can easily count the number of inverted pairs of  $A$  in  $\Theta(n^2)$  time. Design an algorithm that counts the number of inverted pairs in  $O(n \log n)$  time.

\*\* *Q3(c) is the hardest question. Spend more time and try your best to solve it!*

4. (No marks.) Give asymptotic upper bound for  $T(n)$  in each of the following recurrence. Make your bounds as tight as possible.

- (a)  $T(n) = 9 T(n/2) + n^3$
- (b)  $T(n) = 7 T(n/2) + n^3$
- (c)  $T(n) = T(\sqrt{n}) + \log n$
- (d)  $T(n) = 0.5 T(n/2) + n$
- (e)  $T(n) = 3 T(n/3) + n/3$