

On Unfolding Trees and Polygons on Various Lattices

Author: Sheung-Hung Poon,

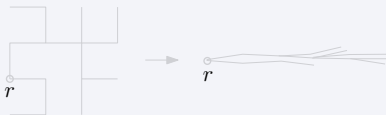
Presented by: Elena Mumford

Technical University of Eindhoven (TU/e),
The Netherlands

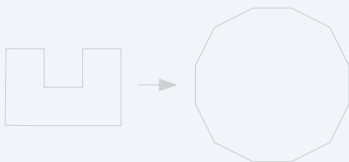
*The 19th Canadian Conference on Computational
Geometry (CCCG), August 20-22, 2007*

Introduction

- Rigid edges can rotate around vertex w/o crossing other edges.
- “*Straightening*” = move rigid edges to lie on a straight line.



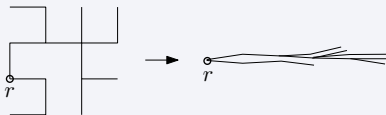
- “*Convexifying*” = move rigid edges until polygon becomes convex.



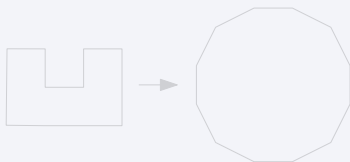
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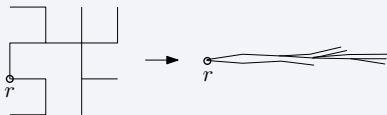
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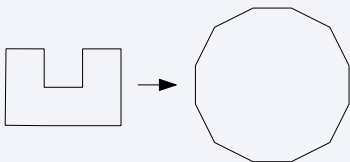
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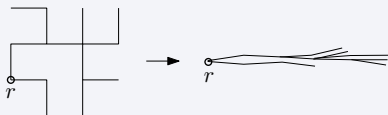
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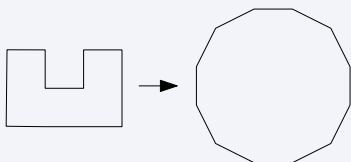
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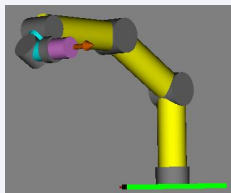


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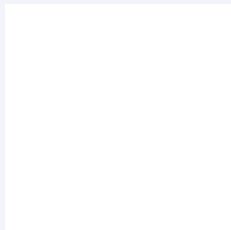


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- Applications:
 - movements of robot arms,



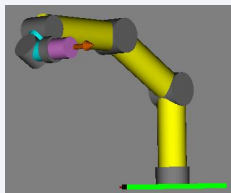
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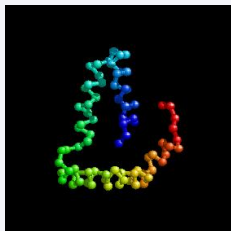
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- wire bending,
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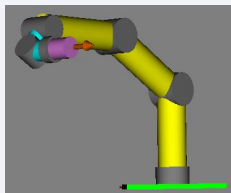
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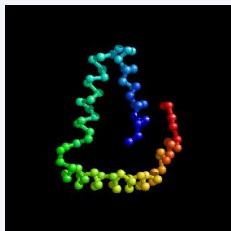
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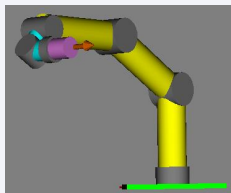
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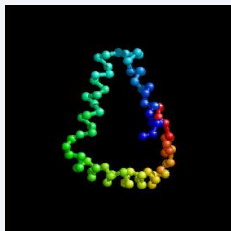
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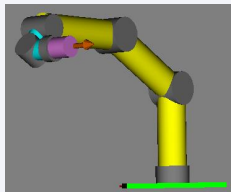
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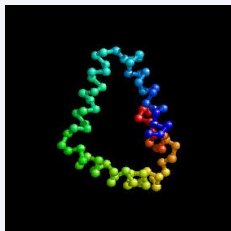
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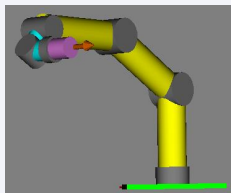
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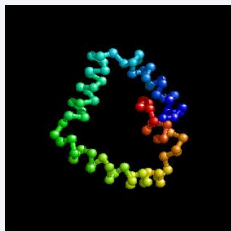
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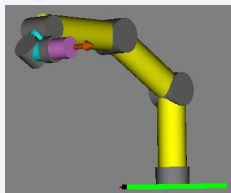
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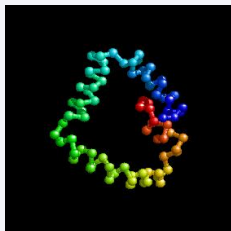
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Definition of Lattice Polygons/Trees

In this talk, we mainly consider lattice polygons/trees.

- A *unit polygon/tree* = with all its edges of unit-length.



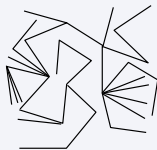
- A *lattice polygon/tree* = with all its edges from a lattice.



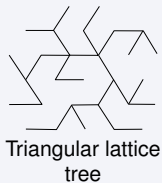
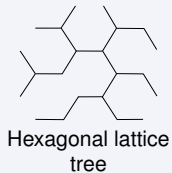
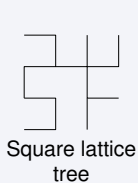
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Previous Work

- *Carpenter's Rule Conjecture* (solved):
Chains (polygons) in 2D can be straightened (convexified).
[Connelly, Demaine and Rote '00][Strienu '00]
- Trees (polygons) in 4D+ can be straightened (convexified).
[Cocan and O'Rourke '01]
- A tree in 2D & a 5-chain in 3D can lock. [Biedl et al '01]



(a) A locked tree in 2D.



(b) A locked 5-chain in 3D

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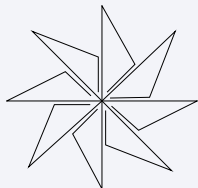
(a) A locked tree in 2D.



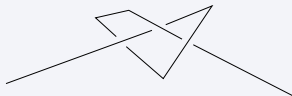
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- PSPACE-complete: to reconfigure 2D-trees or 3D-chains.
[Alt et al. '03]
- A unit tree *of diameter 4* can always be straightened.
[Poon '05]
- A 2D/3D lattice tree can always be straightened, and
A 2D lattice polygon can always be convexified. [Poon '06]

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Our Results

A “*move*” = a monotonic increase/decrease of angle at a vertex.

Theorem

- A hexagonal/triangular *lattice chain* can be straightened in $O(n)$ moves and time (n = no. of edges).

Theorem

- A hexagonal/triangular *lattice tree* can be straightened in $O(n^2)$ moves and time.

Theorem

- A hexagonal/triangular *lattice polygon* can be convexified in $O(n^2)$ moves and time.

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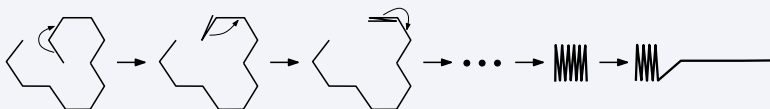
Theorem

- A hexagonal/triangular *lattice polygon* can be convexified in $O(n^2)$ moves and time.

Hexagonal Lattice Chains

Algorithm:

- Fold up the end edges of the chain iteratively.
- Unfold the final folded spring/zig-zag path.



PS. The algorithm is similar to that for *square lattice*.

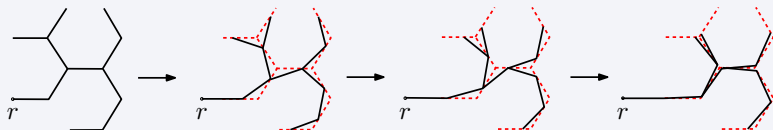
Hexagonal Lattice Trees

- Given a hexagonal lattice tree P .
- Let r be root = the *leftmost* vertex of P .

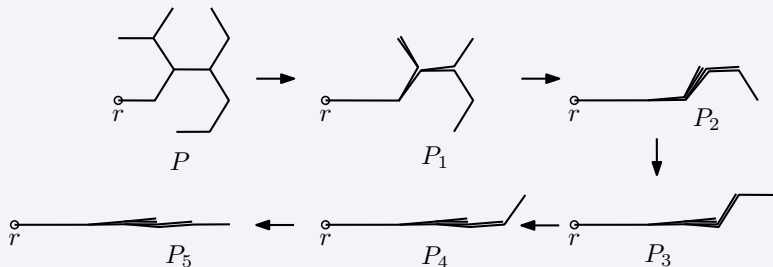
Algorithm:

- Our algorithm proceeds by pulling P to the left successively until the whole tree is straightened.
- In each pulling step:
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- The pulling step is repeated n times so that, finally, tree P is straightened.



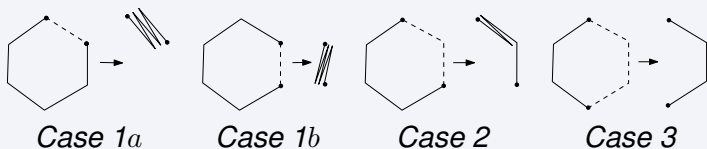
- Each pulling step takes $O(n)$ moves.
Thus n pulling steps take $O(n^2)$ moves in total.

Hexagonal Lattice Polygons

- Use similar technique of *block-collapsing* as square lattice.
- New definition: A *block* = a hexagonal cell.

Algorithm:

Collapse leftmost collapsible block iteratively.



- Observe: operation in *Case 3*, no edges are reduced.
- After at most $O(n)$ operations of *Case 3*, we reach one *Case 1* or *Case 2*.
- Thus whole algorithm takes $O(n^2)$ moves and time.

Triangular Lattice Chains

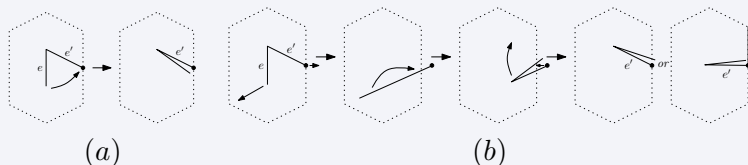
- Use similar technique as for *square & hexagonal lattice*.

Algorithm:

- Fold up the end edges of the chain iteratively.
- Unfold the final folded spring/zig-zag path.

Cases to handle:

- 1 When angle α % end edge & its adjacent edge is $\pi/3$:



- 2 When angle $\alpha = 2\pi/3$ or π , we can handle similarly.

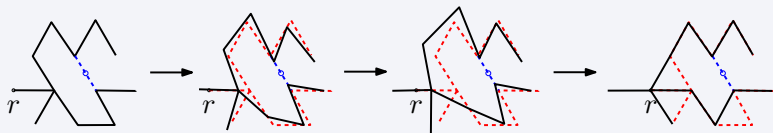
Triangular Lattice Trees

- Given a triangular lattice tree P .
- Let r be root = the *leftmost* vertex of P .

Algorithm:

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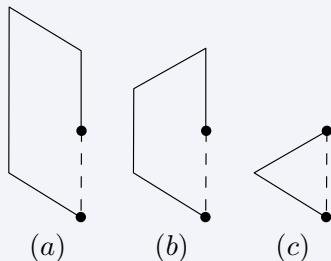
- During motion, a vertex never exceed middle point of an *extension edge* (dashed blue). So *no edge crossings occur*.
- Each pulling step takes $O(n)$ moves, and thus the whole algorithm takes $O(n^2)$ moves and time.

Triangular Lattice Polygons

- We extend *block-collapsing* technique for square/hexagonal lattice.

New definition:

A *block* = a (a) *parallelogram*, (b) *trapezoidal* or (c) *triangular* block.

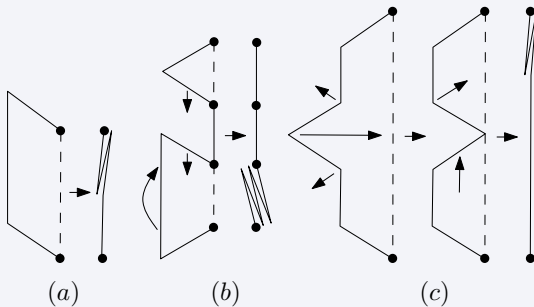


Algorithm:

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Cases to handle:

- (a) Collapse a parallelogram block;
- (b) Collapse two trapezoidal/triangular blocks; and
- (c) Collapse an extended triangular/trapezoidal block.



Conclusions

- We obtain results for chains/trees/polygons on hexagonal or triangular lattice can be straightened/covexified.

Conjecture:

A unit tree in two dimensions can always be straightened.



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