

CS2422 Assembly Language & System Programming

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Today's Topics

- 3.1 Basic Loader Functions
 - Design of an Absolute Loader
 - A Simple Bootstrap Loader
- 3.2 Machine-Dependent Loader Features
 - Relocation
 - Program Linking

Introduction

- To execute an object program, we need
 - Relocation, which modifies the object program so that it can be loaded at an address different from the location originally specified
 - Linking, which combines two or more separate object programs and supplies the information needed to allow references between them
 - Loading and Allocation, which allocates memory location and brings the object program into memory for execution

Overview of Chapter 3

- Type of loaders
 - Assemble-and-go loader
 - Absolute loader (bootstrap loader)
 - Relocating loader (relative loader)
 - Direct linking loader
- Design options
 - Linkage editors
 - Dynamic linking
 - Bootstrap loaders

Assemble-and-go Loader

- Characteristic
 - The object code is stored in memory after assembly
 - Single JUMP instruction
- Advantage
 - Simple, developing environment
- Disadvantage
 - Whenever the assembly program is to be executed, it has to be assembled again
 - Programs have to be coded in the same language

Design of an Absolute Loader

- Absolute Program
 - Advantage
 - Simple and efficient
 - Disadvantage
 - The need for programmer to specify the actual address
 - Difficult to use subroutine libraries
- Program Logic

Fig. 3.2 Algorithm for an absolute loader

```
begin
  read Header record
  verify program name and length
  read first Text record
  while record type <> 'E' do
    begin
      {if object code is in character form, convert into
       internal representation}
      move object code to specified location in memory
      read next object program record
    end
    jump to address specified in End record
  end
```

Loading of an absolute program (1/2)

Fig. 3.1(a)

```
HCOPY 00100000107A
T0010001E1410334820390010362810303010154820613C100300102A0C103900102D
T00101E150C10364820610810334C0000454F46000003000000
T0020391E041030001030E0205D30203FD8205D2810303020575490392C205E38203F
T0020571C1010364C0000F1001000041030E02079302064509039DC20792C1036
T002073073820644C000005
E001000
```

(a) Object program

Loading of an absolute program (2/2)

Fig. 3.1(b)

Memory address	Contents			
0000	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
0010	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
:	:	:	:	:
0FF0	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
1000	14103348	20390010	36281030	30101548
1010	20613C10	0300102A	0C103900	102D0C10
1020	36482061	0810334C	0000454F	46000003
1030	000000xx	xxxxxxxx	xxxxxxxx	xxxxxxxx
:	:	:	:	:
2030	xxxxxxxx	xxxxxxxx	xx041030	001030E0
2040	205D3020	3FD8205D	28103030	20575490
2050	392C205E	38203F10	10364C00	00F10010
2060	00041030	E0207930	20645090	39DC2079
2070	2C103638	20644C00	0005xxxx	xxxxxxxx
2080	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
:	:	:	:	:

← COPY

(b) Program loaded in memory

Figure 3.1 Loading of an absolute program.

Object Code Representation

- Fig. 3.1(a)
 - Each byte of assembled code is given using its hexadecimal representation in character form
 - Easy to read by human beings
- In general
 - Each byte of object code is stored as a single byte
 - Most machine store object programs in a binary form
 - We must be sure that our file and device conventions do not cause some of the program bytes to be interpreted as control characters

A Simple Bootstrap Loader

- Bootstrap Loader
 - When a computer is first tuned on or restarted, a special type of absolute loader, called *bootstrap loader* is executed
 - This bootstrap loads the first program to be run by the computer -- usually an operating system
- Example (SIC bootstrap loader)
 - The bootstrap itself begins at address 0
 - It loads the OS starting address 0x80
 - No header record or control information, the object code is consecutive bytes of memory

Bootstrap loader for SIC/XE

Begin

X=0x80 (the address of the next memory location to be loaded)

Loop

A←GETC (and convert it from the ASCII character code to the value of the hexadecimal digit)

save the value in the high-order 4 bits of S

A←GETC

combine the value to form one byte A← (A+S)

store the value (in A) to the address in register X

X←X+1

End

GETC A←read one character
if A=0x04 then jump to 0x80
if A<48 then GETC
A ← A-48 (0x30)
if A<10 then return
A ← A-7
return

Fig. 3.3 Bootstrap loader for SIC/XE

```
BOOT    START    0        BOOTSTRAP LOADER FOR SIC/XE
.
. THIS BOOTSTRAP READS OBJECT CODE FROM DEVICE F1 AND ENTERS IT
. INTO MEMORY STARTING AT ADDRESS 80 (HEXADECIMAL). AFTER ALL OF
. THE CODE FROM DEVF1 HAS BEEN SEEN ENTERED INTO MEMORY, THE
. BOOTSTRAP EXECUTES A JUMP TO ADDRESS 80 TO BEGIN EXECUTION OF
. THE PROGRAM JUST LOADED. REGISTER X CONTAINS THE NEXT ADDRESS
. TO BE LOADED.
.
      CLEAR      A        CLEAR REGISTER A TO ZERO
      LDX        #128     INITIALIZE REGISTER X TO HEX 80
LOOP    JSUB     GETC     READ HEX DIGIT FROM PROGRAM BEING LOADED
      RMO       A,S      SAVE IN REGISTER S
      SHIFTL    S,4      MOVE TO HIGH-ORDER 4 BITS OF BYTE
      JSUB     GETC     GET NEXT HEX DIGIT
      ADDR     S,A      COMBINE DIGITS TO FORM ONE BYTE
      STCH     0,X      STORE AT ADDRESS IN REGISTER X
      TIXR    X,X      ADD 1 TO MEMORY ADDRESS BEING LOADED
      J        LOOP     LOOP UNTIL END OF INPUT IS REACHED
.
. SUBROUTINE TO READ ONE CHARACTER FROM INPUT DEVICE AND
. CONVERT IT FROM ASCII CODE TO HEXADECIMAL DIGIT VALUE. THE
. CONVERTED DIGIT VALUE IS RETURNED IN REGISTER A. WHEN AN
. END-OF-FILE IS READ, CONTROL IS TRANSFERRED TO THE STARTING
. ADDRESS (HEX 80).
.
GETC    TD       INPUT    TEST INPUT DEVICE
      JEQ      GETC     LOOP UNTIL READY
      RD       INPUT    READ CHARACTER
      COMP     #4       IF CHARACTER IS HEX 04 (END OF FILE),
      JEQ      80      JUMP TO START OF PROGRAM JUST LOADED
      COMP     #48     COMPARE TO HEX 30 (CHARACTER '0')
      JLT     GETC     SKIP CHARACTERS LESS THAN '0'
      SUB      #48     SUBTRACT HEX 30 FROM ASCII CODE
      COMP     #10     IF RESULT IS LESS THAN 10, CONVERSION IS
      JLT     RETURN   COMPLETE. OTHERWISE, SUBTRACT 7 MORE
      SUB      #7      (FOR HEX DIGITS 'A' THROUGH 'F')
RETURN  RSUB    RETURN  TO CALLER
INPUT  BYTE    X'F1'    CODE FOR INPUT DEVICE
      END     LOOP
```

Figure 3.3 Bootstrap loader for SIC/XE.

Relocating Loaders

- Motivation
 - Efficient sharing of the machine with larger memory and when several independent programs are to be run together
 - Support the use of subroutine libraries efficiently
- Two methods for specifying relocation
 - Modification record (Fig. 3.4, 3.5)
 - Relocation bit (Fig. 3.6, 3.7)
 - Each instruction is associated with one relocation bit
 - These relocation bits in a Text record is gathered into bit masks

Modification Record

- For complex machines
- Also called RLD specification
 - Relocation and Linkage Directory

Modification record col 1: M col 2-7: relocation address col 8-9: length (halfbyte) col 10: flag (+/-) col 11-17: segment name

Fig. 3.4: An SIC/XE Example (from Fig. 2.6)

Line	Loc	Source statement	Object code
5	0000	COPY START 0	
10	0000	FIRST STL RETADR	17202D
12	0003	LDB #LENGTH	69202D
13		BASE LENGTH	
15	0006	CLOOP +JSUB RDREC	4B101036
20	000A	LDA LENGTH	032026
25	000D	COMP #0	290000
30	0010	JEQ ENDFIL	332007
35	0013	+JSUB WRREC	4B10105D
40	0017	J CLOOP	3F2FEC
45	001A	ENDFIL LDA EOF	032010
50	001D	STA BUFFER	0F2016
55	0020	LDA #3	010003
60	0023	STA LENGTH	0F200D
65	0026	+JSUB WRREC	4B10105D
70	002A	J @RETADR	3E2003
80	002D	EOF BYTE C' EOF'	454F46
95	0030	RETADR RESW 1	
100	0033	LENGTH RESW 1	
105	0036	BUFFER RESB 4096	

```

115      .      READ RECORD INTO BUFFER
120      .
125 1036 RDREC  CLEAR X      B410
130 1038      CLEAR A      B400
132 103A      CLEAR S      B440
133 103C      +LDT #4096    75101000
135 1040 RLOOP  TD      INPUT  E32019
140 1043      JEQ      RLOOP  332FFA
145 1046      RD      INPUT  DB2013
150 1049      COMPR  A,S    A004
155 104B      JEQ      EXIT   332008
160 104E      STCH   BUFFER,X 57C003
165 1051      TIXR   T      B850
170 1053      JLT    RLOOP  3B2FEA
175 1056 EXIT   STX    LENGTH 134000
180 1059      RSUB   4F0000
185 105C INPUT  BYTE   X'F1'   F1
195      .
200      .      WRITE RECORD FROM BUFFER
205      .
210 105D WRREC  CLEAR X      B410
212 105F      LDT    LENGTH 774000
215 1062 WLOOP  TD      OUTPUT E32011
220 1065      JEQ    WLOOP  332FFA
225 1068      LDCH   BUFFER,X 53C003
230 106B      WD     OUTPUT DF2008
235 106E      TIXR   T      B850
      ... (omitted)

```

Fig. 3.5

```

^H^C^O^P^Y  ^0^0^0^0^0^0^0^1^0^7^7
T^0^0^0^0^0^1^D^1^7^2^0^2^D^6^9^2^0^2^D^4^B^1^0^1^0^3^6^0^3^2^0^2^6^2^9^0^0^0^0^3^3^2^0^0^7^4^B^1^0^1^0^5^D^3^F^2^F^E^C^0^3^2^0^1^0
T^0^0^0^0^1^D^1^3^0^F^2^0^1^6^0^1^0^0^0^3^0^F^2^0^0^D^4^B^1^0^1^0^5^D^3^E^2^0^0^3^4^5^4^F^4^6
T^0^0^1^0^3^6^1^D^B^4^1^0^B^4^0^0^B^4^4^0^7^5^1^0^1^0^0^0^E^3^2^0^1^9^3^3^2^F^F^A^D^B^2^0^1^3^A^0^0^4^3^3^2^0^0^8^5^7^C^0^0^3^B^8^5^0
T^0^0^1^0^5^3^1^D^3^B^2^F^E^A^1^3^4^0^0^0^4^F^0^0^0^0^F^1^B^4^1^0^7^7^4^0^0^0^E^3^2^0^1^1^3^3^2^F^F^A^5^3^C^0^0^3^D^F^2^0^0^8^B^8^5^0
T^0^0^1^0^7^0^0^7^3^B^2^F^E^F^4^F^0^0^0^0^0^5
M^0^0^0^0^0^7^0^5^+^C^O^P^Y
M^0^0^0^0^1^4^0^5^+^C^O^P^Y
M^0^0^0^0^2^7^0^5^+^C^O^P^Y
E^0^0^0^0^0^0

```

Figure 3.5 Object program with relocation by Modification records.

Relocation Bit

- For simple machines
- Relocation bit
 - 0: No modification is necessary
 - 1: Modification is needed

Text record
col 1: T
col 2-7: starting address
col 8-9: length (byte)
col 10-12: relocation bits
col 13-72: object code

- Twelve-bit mask is used in each Text record
 - Since each text record contains less than 12 words
 - Unused words are set to 0
 - Any value that is to be modified during relocation must coincide with one of these 3-byte segments
 - e.g. line 210

Figure 3.6: Relocatable Program for SIC

Line	Loc	Source statement	Object code
5	0000	COPY START 0	
10	0000	FIRST STL RETADR	140033
15	0003	CLOOP JSUB RDREC	481039
20	0006	LDA LENGTH	000036
25	0009	COMP ZERO	280030
30	000C	JEQ ENDFIL	300015
35	000F	JSUB WRREC	481061
40	0012	J CLOOP	3C0003
45	0015	ENDFIL LDA EOF	00002A
50	0018	STA BUFFER	0C0039
55	001B	LDA THREE	00002D
60	001E	STA LENGTH	0C0036
65	0021	JSUB WRREC	481061
70	0024	LDL RETADR	080033
75	0027	RSUB	4C0000
80	002A	EOF BYTE C' EOF'	454F46
85	002D	THREE WORD 3	000003
90	0030	ZERO WORD 0	000000
95	0033	RETADR RESW 1	
100	0036	LENGTH RESW 1	
105	0039	BUFFER RESB 4096	

```

110      .
115      .          SUBROUTINE TO READ RECORD INTO BUFFER
120      .
125 1039 RDREC   LDX    ZERO      040030
130 103C      LDA    ZERO      000030
135 103F      RLOOP  TD     INPUT   E0105D
140 1042      JEQ   RLOOP  30103D
145 1045      RD     INPUT   D8105D
150 1048      COMP  ZERO      280030
155 104B      JEQ   EXIT   301057
160 104E      STCH  BUFFER,X  548039
165 1051      TIX   MAXLEN  2C105E
170 1054      JLT   RLOOP  38103F
175 1057      EXIT  STX    LENGTH  100036
180 105A      RSUB  4C0000
185 105D      INPUT BYTE   X'F1'   F1
190 105E      MAXLEN WORD   4096    001000
195      .
200      .          SUBROUTINE TO WRITE RECORD FROM BUFFER
205      .
210 1061 WRREC  LDX    ZERO      040030
215 1064 WLOOP  TD     OUTPUT  E01079
220 1067      JEQ   WLOOP  301064
225 106A      LDCH  BUFFER,X  508039
230 106D      WD     OUTPUT  DC1079
235 1070      TIX   LENGTH  2C0036
240 1073      JLT   WLOOP  381064
245 1076      RSUB  4C0000
250 1079      OUTPUT BYTE   X'05'   05
255      END    FIRST

```

Fig. 3.7

```

HCOPY 00000000107A
T0000001EFFC1400334810390000362800303000154810613C000300002A0C003900002D
T00001E15E000C00364810610800334C0000454F46000003000000
T0010391EFFC040030000030E0105D30103FD8105D2800303010575480392C105E38103F
T0010570A8001000364C0000F1001000
T00106119FE0040030E01079301064508039DC10792C00363810644C000005
E000000

```

Figure 3.7 Object program with relocation by bit mask.

Program Linking

- Goal
 - Resolve the problems with EXTREF and EXTDEF from different control sections (sec 2.3.5)
- Example
 - Program in Fig. 3.8 and object code in Fig. 3.9
 - Use modification records for both relocation and linking
 - Address constant
 - External reference

Fig. 3.8-1

Loc		Source statement	Object code
0000	PROGA	START 0	
		EXTDEF LISTA, ENDA	
		EXTREF LISTB, ENDB, LISTC, ENDC	
		.	
		.	
0020	REF1	LDA LISTA	03201D
0023	REF2	+LDT LISTB+4	77100004
0027	REF3	LDX #ENDA-LISTA	050014
		.	
		.	
0040	LISTA	EQU *	
		.	
		.	
0054	ENDA	EQU *	
0054	REF4	WORD ENDA-LISTA+LISTC	000014
0057	REF5	WORD ENDC-LISTC-10	FFFFFF6
005A	REF6	WORD ENDC-LISTC+LISTA-1	00003F
005D	REF7	WORD ENDA-LISTA- (ENDB-LISTB)	000014
0060	REF8	WORD LISTB-LISTA	FFFFFFC0
		END REF1	

Fig. 3.9-1

```

HPROGA 000000000063
DLISTA 000040ENDA 000054
RLISTB ENDB LISTC ENDC
.
.
T0000200A03201D77100004050014
.
.
T0000540F000014FFFFF600003F000014FFFFC0
M00002405+LISTB
M00005406+LISTC
M00005706+ENDC
M00005706-LISTC
M00005A06+ENDC
M00005A06-LISTC
M00005A06+PROGA
M00005D06-ENDB
M00005D06+LISTB
M00006006+LISTB
M00006006-PROGA
E000020

```

Figure 3.9 Object programs corresponding to Fig. 3.8.

Fig. 3.8-2

Loc	Source statement	Object code
0000	PROGB START 0 EXTDEF LISTB, ENDB EXTREF LISTA, ENDA, LISTC, ENDC . . .	
0036	REF1 +LDA LISTA	03100000
003A	REF2 LDT LISTB+4	772027
003D	REF3 +LDX #ENDA-LISTA . . .	05100000
0060	LISTB EQU * . .	
0070	ENDB EQU *	
0070	REF4 WORD ENDA-LISTA+LISTC	000000
0073	REF5 WORD ENDC-LISTC-10	FFFFFF6
0076	REF6 WORD ENDC-LISTC+LISTA-1	FFFFFFF
0079	REF7 WORD ENDA-LISTA- (ENDB-LISTB)	FFFFFF0
007C	REF8 WORD LISTB-LISTA END	000060

Figure 3.8 Sample programs illustrating linking and relocation.

Fig. 3.9-2

```

HPROGB 00000000007F
DLISTB 000060ENDB 000070
RLISTA ENDA LISTC ENDC
:
T0000360B0310000077202705100000
:
T0000700F000000FFFFFF6FFFFFFF0000060
M00003705+LISTA
M00003E05+ENDA
M00003E05-LISTA
M00007006+ENDA
M00007006-LISTA
M00007006+LISTC
M00007306+ENDC
M00007306-LISTC
M00007606+ENDC
M00007606-LISTC
M00007606+LISTA
M00007906+ENDA
M00007906-LISTA
M00007C06+PROGB
M00007C06-LISTA
E

```

Fig. 3.8-3

Loc	Source statement	Object code
0000	PROGC START 0 EXTDEF LISTC, ENDC EXTREF LISTA, ENDA, LISTB, ENDB . .	
0018	REF1 +LDA LISTA	03100000
001C	REF2 +LDT LISTB+4	77100004
0020	REF3 +LDX #ENDA-LISTA . .	05100000
0030	LISTC EQU *	
0042	ENDC EQU *	
0042	REF4 WORD ENDA-LISTA+LISTC	000030
0045	REF5 WORD ENDC-LISTC-10	000008
0048	REF6 WORD ENDC-LISTC+LISTA-1	000011
004B	REF7 WORD ENDA-LISTA-(ENDB-LISTB)	000000
004E	REF8 WORD LISTB-LISTA END	000000

Figure 3.8 (cont'd)

Fig. 3.9-3

```
HPROGC 000000000051
DLISTC 000030ENDC 000042
RLISTA ENDA LISTB ENDB
.
.
T0000180C0310000Q7710000405100000
.
.
T0000420F000030000008000011000000000000
M00001905+LISTA
M00001D05+LISTB
M00002105+ENDA
M00002105-LISTA
M00004206+ENDA
M00004206-LISTA
M00004206+PROGC
M00004806+LISTA
M00004B06+ENDA
M00004B06-LISTA
M00004B06-ENDB
M00004E06+LISTB
M00004E06-LISTA
E
```

Figure 3.9 (cont'd)

Program Linking Example

- Fig. 3.10
- Load address for control sections
 - PROGA 004000 63
 - PROGB 004063 7F
 - PROGC 0040E2 51
- Load address for symbols
 - LISTA: PROGA+0040=4040
 - LISTB: PROGB+0060=40C3
 - LISTC: PROGC+0030=4112
- REF4 in PROGA
 - ENDA-LISTA+LISTC=14+4112=4126
 - T0000540F000014FFFFFF600003F000014FFFFFFC0
 - M00005406+LISTC

Fig. 3.10(a)

Memory address	Contents			
0000	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
⋮	⋮	⋮	⋮	⋮
3FF0	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
4000
4010
4020	03201D77	1040C705	0014 ← PROGA
4030
4040
4050	00412600	00080040	51000004
4060	000083
4070
4080
4090	031040	40772027 ← PROGB
40A0	05100014
40B0
40C0
40D0	00412600	08004051	00000400
40E0	0083
40F0	0310	40407710 ← PROGC
4100	40C70510	0014
4110
4120	00412600	00080040	51000004
4130	000083xx	xxxxxxxx	xxxxxxxx	xxxxxxxx
4140	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
⋮	⋮	⋮	⋮	⋮

Figure 3.10(a) Programs from Fig. 3.8 after linking and loading.

Fig. 3.10(b)

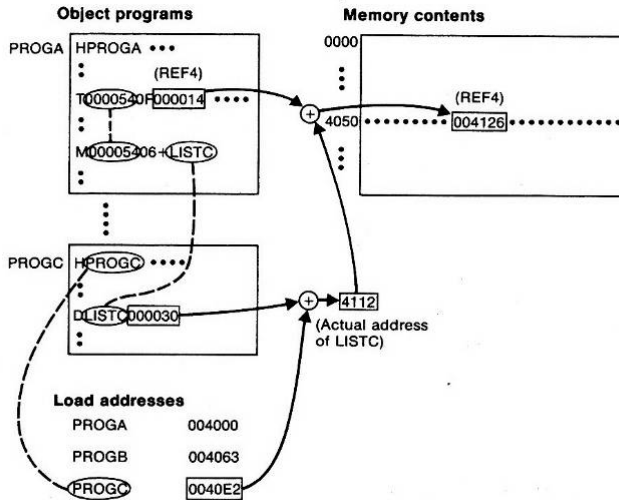


Figure 3.10(b) Relocation and linking operations performed on REF4 from PROGA.