

CS 3331 Numerical Methods

Lecture 1: Introduction

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About this course

- Text: *Applied Numerical Analysis using Matlab, 2nd edition*
– Laurene V. Fausett. (LVF)

- TA: TBA

- Website:

<http://www.cs.nthu.edu.tw/~cherung/teaching/cs3331/cs3331.html>

- Office hours: Tuesday 2:00-3:00, Friday 3:00-4:00 (or by appointment).

Tentative agenda

- 1 Introduction (chap 1)
- 2 Functions of one variable (chap 2)
- 3 Linear systems (chap 3, 4)
- 4 Linear least square problems (chap 4,9)
- 5 Eigenvalues/eigenvectors (chap 5)
- 6 Iterative methods for solving linear systems (chap 6)
- 7 Interpolation (chap 8)
- 8 Approximation (chap 9)
- 9 Fourier methods (chap 10)
- 10 Numerical differentiation and integration (chap 11)
- 11 Numerical optimization (chap 2, 7)

Pre-requirements

- Calculus: mean value theory, Taylor expansion ...
- Linear algebra: symmetric matrix, orthogonal matrix, eigenvalues/eigenvectors ...
- Computer science: floating-point arithmetic, algorithm ...
- Programming: Matlab, c/c++

Grading

- Quiz (50%)
 - every 1-2 weeks
- Assignment (50%)
 - 4-5 (programming) projects

QUESTIONS?

Introduction

Numerical methods

- Numerical vs. Analytical
- Continuous vs. Discrete
- Examples: solving nonlinear equations, linear systems, numerical integration ...

Nonlinear equations LVF pp.4-5

- Solve $f(x) = x^2 - 3 = 0$ ($x = \pm\sqrt{3}$).

- Fixed point iterations:

– rewrite $x^2 - 3 = 0$ as $x = \frac{1}{2}\left(x + \frac{3}{x}\right)$

$$x_0 = 1$$

$$x_1 = \frac{1}{2}\left(1 + \frac{3}{1}\right) = 2$$

$$x_2 = \frac{1}{2}\left(2 + \frac{3}{2}\right) = \frac{7}{4}$$

$$\vdots = \vdots$$

Linear system LVF pp.6-7

$$\begin{aligned}L_1 &: 4x_1 + x_2 = 6 \\M_1 &: -x_1 + 5x_2 = 9\end{aligned}$$

- Gaussian elimination.

$$\text{– Solve } \begin{pmatrix} 4 & 1 \\ -1 & 5 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 6 \\ 9 \end{pmatrix}$$

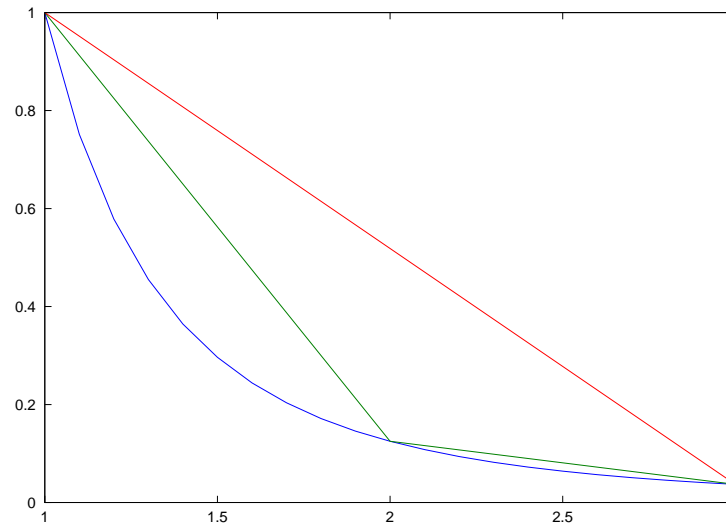
$$\Rightarrow \begin{pmatrix} 1 & 0 \\ 1/4 & 1 \end{pmatrix} \begin{pmatrix} 4 & 1 \\ -1 & 5 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 1/4 & 1 \end{pmatrix} \begin{pmatrix} 6 \\ 9 \end{pmatrix}$$

$$\Rightarrow \begin{pmatrix} 4 & 1 \\ 0 & 5.25 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 6 \\ 10.5 \end{pmatrix}$$

$$\Rightarrow \text{Back-substitution: } x_2 = 2, x_1 = \frac{1}{4}(6 - 1 * 2) = 1$$

Numerical integration LVF pp.8-9

- Compute $I = \int_1^3 \frac{1}{x^3} dx$
- Trapezoid rule.



| Method | Formula | Result |
|---------------------|--------------------------------------|----------|
| analytical solution | $\left. \frac{-1}{2x^2} \right _1^3$ | 0.444444 |
| one subdivision | $2/2[1 + 1/27]$ | 1.037037 |
| two subdivisions | $1/2[1 + 1/8] + 1/2[1/8 + 1/27]$ | 0.643518 |

Basic issues of numerical methods

- Accuracy: (errors)
- Speed: (cost)
 - Time complexity (operation counts).
 - Converge rate.
 - Machine/software properties.

Real Numbers in Computer

Real number in computer LVF pp.13

- Binary representation

$$\begin{aligned} N &= (d_k d_{k-1} \cdots d_1 d_0 . d_{-1} \cdots d_{-p})_b \\ &= d_k 2^k + d_{k-1} 2^{k-1} \cdots 2d_1 + d_0 + d_{-1} \frac{1}{2} + d_{-2} \frac{1}{4} + \cdots + d_{-p} \frac{1}{2^p} \end{aligned}$$

– d_1, d_2, \dots, d_p are in $\{0, 1\}$.

– Alternative representation $(d_k . d_{k-1} \cdots d_1 d_0 d_{-1} \cdots d_{-p})_b \times 2^k$

- Floating-point vs. fixed-point
- Multiprecision and arbitrary precision

IEEE 754 LVF pp.13-15



$$a = (-1)^s \times 2^{\text{exponent} - \text{exponent bias}} \times 1.\text{mantissa}$$

| | single (32bits) | double (64bits) |
|----------|--------------------|----------------------|
| s | 1 bit | 1 bit |
| exponent | 8 bits ($e = 8$) | 11 bits ($e = 11$) |
| mantissa | 23 bits | 52 bits |

- Normalization: the leading digit is 1.
 - Subnormal: when the exponent is the smallest number, the leading digit is allowed to be zero
- Exponent bias: exponent are shifted by $2^{e-1} - 1$.

IEEE 754–continue

- Special numbers
 - Inf: (Infinite) exponent= $2^e - 1$, mantissa=0.
 - NaN: (Not a Number) exponent= $2^e - 1$, mantissa $\neq 0$.
 - Zeros: exponent=0, mantissa=0.
- Representable ranges:

| absolute values | single (32bits) | double (64bits) |
|-----------------|------------------------|---------------------------|
| Min. normal | 2^{-126} | 2^{-1022} |
| Min. subnormal | 2^{-149} | 2^{-1074} |
| Max. finite | $(1 - 2^{-24})2^{128}$ | $(1 - (2)^{-53})2^{1024}$ |

- Overflow and underflow

Multiple precision arithmetic

- Double-double approach: use two doubles for a number.
- Extend floating-point format: use more bits for exponent and mantissa.
- Software:
 - vpa (Matlab Symbolic Math Toolbox),
 - GNU Multi-Precision Library (c/c++),
 - ARPREC and MPFUN (Fortran),
 - Bignum and BigInteger (Java).

Errors

Source of errors

- From measurement/sampling.
- From modeling.
- From number representation.
- From algorithm.

Measuring errors LVF pp.16

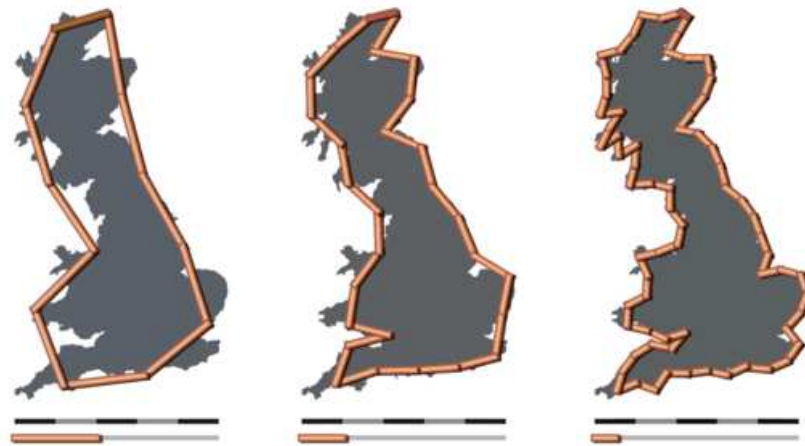
- Let x^* be the exact value.
 - Absolute error: $\text{Error}(x) = |x - x^*|$.
 - Relative error: $\text{Rel Error}(x) = |x - x^*| / |x^*|$.
- *Significant digits*: The number x is said to have t significant digits if t is the largest nonnegative integer for which

$$\frac{|x - x^*|}{|x^*|} < 5 \times 10^{-t}.$$

- Big-Oh notation: $f(h) = O(g(h))$ if $f(h) \leq c|g(h)|$ for some positive constant c when $h \rightarrow 0$. LVF pp.21

Errors from modeling

- Benoit Mandelbrot, *How Long Is the Coast of Britain? Statistical Self-Similarity and Fractional Dimension*, Science Vol 156, 1967.



Images from Wikipaida

Errors from inexact representation LVF pp.17

- Machine epsilon (*eps*): the difference between 1 and the smallest exactly representable number greater than one.
 - Single: $2^{-24} = 5.96 \times 10^{-8}$
 - Double: $2^{-53} = 1.1 \times 10^{-16}$
- Rounding
 - Round to nearest even
 - Example: $1/3$ (0.0101010**1**→0.0101010) and $1/5$ (0.0011001**1**→0.0011010)
 - Roundoff error: $1/3 + 1/6$

Errors from floating-point arithmetic L VF pp.17

- Cancellation: loss of significance
- Example: compute $2^1 \times 0.100 - 2^0 \times 0.111$
 - Alignment: $2^1 \times 0.100 - 2^1 \times 0.011$
 - Result: $2^1 \times 0.001 = 2^{-1} \times 0.100$
 - But the exact result is $2^{-2} \times 0.100$.
 - Relative error: $\frac{|2^{-1} \times 0.100 - 2^{-2} \times 0.100|}{|2^{-2} \times 0.100|} = 1.$

Errors from numerical algorithm LVF pp.18

- Example: Solve $ax^2 + bx + c = 0$ using $x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$.
 - Ex. $x^2 + 100x + 1 = 0$
 - * $\sqrt{100^2 - 4} = \sqrt{9996}$ is rounded to 100.
 - $\Rightarrow x_1 = (-100 + 100)/2 = 0, x_2 = (-100 - 100)/2 = -100.$
 - * Actual solution $x_1 = -0.01, x_2 = -99.99$
 - $\Rightarrow \text{RE}(x_1) = 1, \text{RE}(x_2) = 10^{-4}.$
 - Use $x = \frac{-2c}{b + \sqrt{b^2 - 4ac}}$ for x_1 . ($x_1 = -0.01$)

Forward and backward error analysis

- What we concern is the errors in the solutions (output).
- Example: evaluate $f(x)$.
 - Let $y = f(x)$ and \hat{y} be the computed result.
 - Forward error: error of the output: $|y - \hat{y}|$
 - Backward error: given the computed output \hat{y} , backward error is the smallest $|\Delta x|$ such that $f(x + \Delta x) = \hat{y}$.
 - Example: Evaluate $f(x) = \sqrt{x}$ at $x = 1/36$.

Condition number

- Let x be an input and $f(x)$ be its output. \tilde{x} is a perturbed x .
- Condition number = $\frac{|f(\tilde{x}) - f(x)|/|f(x)|}{|\tilde{x} - x|/|x|}$
- If $f(x)$ is continuously differentiable around x , the condition number is

$$\left| \frac{xf'(x)}{f(x)} \right|$$

- ex: $f(x) = x^{-1}$ for $x > 0$.

Last year's notes (<http://www.cs.nthu.edu.tw/~cchen>)

Performance Issues

Computational efforts LVF pp.30

- flops: floating-point operations
- Example: polynomial evaluation $P(x) = a_n x^n + \cdots + a_1 x + a_0$
 - Direct method: evaluate $a_i x^i$ one by one

$$\text{flops} = \left(\sum_{k=1}^n k \right) + n = \frac{n(n+3)}{2}$$

- Horner's algorithm: evaluate

$$P(x) = (\dots((a_n x) + a_{n-1})x + \cdots + a_1)x + a_0$$

$$\text{flops} = \sum_{k=1}^n 2 = 2n$$

Convergence LVF pp.22-23

- A sequence $\{y_k\}$ converges to y^* iff $\lim_{k \rightarrow \infty} |y_k - y^*| = 0$
- If there existing some $\lambda > 0, p > 0$ such that

$$\lim_{k \rightarrow \infty} \frac{|y_k - y^*|}{|y_{k-1} - y^*|^p} = \lambda,$$

the sequence $\{y_k\}$ is called converging to y^* with order p .
The number λ is called the asymptotic error constant.

| | | |
|-------------------------|----------------------|------------------|
| Sublinear convergence | $p = 1, \lambda = 1$ | $y_k = 1/k$ |
| Linear convergence | $p = 1, \lambda < 1$ | $y_k = 2^{-k}$ |
| Superlinear convergence | $p > 1$ | $y_k = k^{-k}$ |
| Quadratic convergence | $p = 2$ | $y_k = 2^{-2^k}$ |

Stopping criteria LVF pp.25

- In practice, we do not know y^* .
- For equations, $f(x) = b$, we can measure the *residual*
$$|f(x_k) - b|$$
- Some other stopping criteria (none guarantees convergence)
 - $|y_k - y_{k-1}|/|y_k| < \text{tol}$
 - $k > \text{maxIter}$
 - $|x_k - x_{k-1}| < \text{tol}^*$

*maxIter and tol are pre-specified constants.