

Matrices and Linear Systems of Equations

- ♣ Matrix notations and operations
- ♣ A Linear System of Equations
- ♣ Matlab Solutions for A Linear System of Equations
- ♣ Elementary row operations
- ♣ LU-decomposition
- ♣ Gaussian Elimination with Partial Pivoting
- ♣ Row-echelon form
- ♣ Matrix inverse and Transpose of a Matrix
- ♣ Computing An Inverse Matrix By Elementary Row Operations
- ♣ Some special matrices
- ♣ Analysis of Gaussian Elimination and Back Substitution

Matrix Notations and Operations

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdot & \cdot & a_{1n} \\ a_{21} & a_{22} & \cdot & \cdot & a_{2n} \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ a_{m1} & a_{m2} & \cdot & \cdot & a_{mn} \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} b_1 \\ b_2 \\ \cdot \\ \cdot \\ b_m \end{bmatrix}$$

- $A = [a_{ij}]$ or (a_{ij}) , $i = 1, 2, \dots, m$; $j = 1, 2, \dots, n$ or denote $A \in R^{m \times n}$
- An $m \times 1$ matrix is called a column vector, denote $\mathbf{b} \in R^m$
- A $1 \times n$ matrix is called a row vector, denote \mathbf{y}^t or $\mathbf{y}' \in R^n$, where $\mathbf{y} = [y_1, y_2, \dots, y_n]$
- Let $X, Y \in R^{m \times n}$, $X = [x_{ij}]$, $Y = [y_{ij}]$, define $X + Y = [x_{ij} + y_{ij}]$
- $C = AB$ for $A \in R^{m \times n}$, $B \in R^{n \times q}$ is well-defined only when $n = p$

$$C = [c_{ij}] \in R^{m \times q}, \text{ where } c_{ij} = \sum_{k=1}^n a_{ik}b_{kj}$$

- $C = A[\mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_q] = [A\mathbf{b}_1, A\mathbf{b}_2, \dots, A\mathbf{b}_q]$
- Matrix multiplication is associative $(AB)C = A(BC)$ but not commutative $AB \neq BA$

◇ The following linear system of equations can be written as $A\mathbf{x} = \mathbf{b}$ with the *augmented* matrix $[A \mid \mathbf{b}]$.

$$a_{11}x_1 + a_{12}x_2 + \cdot \cdot + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \cdot \cdot + a_{2n}x_n = b_2$$

$$\cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot$$

$$\cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot$$

$$a_{m1}x_1 + a_{m2}x_2 + \cdot \cdot + a_{mn}x_n = b_m$$

The system is *homogeneous* if $b_1 = b_2 = \dots = b_m = 0$, *overdetermined* if $m > n$, and *underdetermined* if $m < n$

Overdetermined, Underdetermined, Homogeneous Systems

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2$$

$$\cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot$$

$$\cdot \quad \cdot \quad \cdot \quad \cdot \quad \cdot$$

$$a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n = b_m$$

$$A\mathbf{x} = \mathbf{b}$$

Definition: A linear system is said to be *overdetermined* if there are more equations than unknowns ($m > n$), *underdetermined* if $m < n$, *homogeneous* if $b_i = 0, \forall 1 \leq i \leq m$.

$$\begin{array}{lll} x + y = 1 & x + y = 3 & x + y = 2 \\ (A) \quad x - y = 3 & (B) \quad x - y = 1 & (C) \quad 2x + 2y = 4 \\ -x + 2y = -2 & 2x + y = 5 & -x - y = -2 \end{array}$$

(A) has no solution, (B) has unique solution, (C) has infinitely many solutions

$$\begin{array}{ll} (D) \quad \begin{array}{l} x + 2y + z = -1 \\ 2x + 4y + 2z = 3 \end{array} & (E) \quad \begin{array}{l} x + 2y + z = 5 \\ 2x - y + z = 3 \end{array} \end{array}$$

(D) has no solution, (E) has infinitely many solutions

A Direct Solution of Linear Systems

A linear system

$$\begin{aligned}2x + y + z &= 5 \\4x - 6y &= -2 \\-2x + 7y + 2z &= 9\end{aligned}$$

A matrix representation

$$A\mathbf{x} = \mathbf{b}, \text{ or } \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \end{bmatrix} = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix}.$$

♣ Solution using MATLAB

```
>> A = [2, 1, 2; 4, -6, 0; -2, 7, 2];  
>> b = [5, -2, 9]';  
>> x = A\b (x = [1; 1; 2])
```

Elementary Row Operations

- (1) Interchange two rows: $A_r \leftrightarrow A_s$
- (2) Multiply a row by a nonzero real number: $A_r \leftarrow \alpha A_r$
- (3) Replace a row by its sum with a multiple of another row: $A_s \leftarrow \alpha A_r + A_s$

$$E_1 = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad E_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -2 \end{bmatrix}, \quad E_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix},$$

♣ *Example*

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}, \quad E_1 A = \begin{bmatrix} 4 & -6 & 0 \\ 2 & 1 & 1 \\ -2 & 7 & 2 \end{bmatrix}, \quad E_2 A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ 4 & -14 & -4 \end{bmatrix}, \quad E_3 A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ 0 & 8 & 3 \end{bmatrix}$$

Let

$$L_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \quad L_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad L_1 = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}$$

Then

$$L_3 L_2 L_1 A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = U \quad (\text{Upper } - \Delta)$$

$$A = (L_1^{-1} L_2^{-1} L_3^{-1}) U = LU, \quad \text{where } L \text{ is unit lower } - \Delta$$

Matrix Inverse and Transpose

A matrix $A \in R^{n \times n}$ is invertible iff $\exists B$ such that $AB = BA = I_n$. If the inverse B exists, B is unique and is denoted by A^{-1}

(1) A matrix is *nonsingular* if it is invertible

(2) If there exists two inverses B and C , then $B=C$

(3) The inverse of A^{-1} is A

(4) A diagonal matrix is invertible if none of its diagonal entries is 0

(5) If $A = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$, then $A^{-1} = \frac{1}{ad-bc} \begin{bmatrix} d & -b \\ -c & a \end{bmatrix}$

(6) If both A and B are invertible, then $(AB)^{-1} = B^{-1}A^{-1}$

(7) The transpose of $A = [a_{ij}]$ is $A^t = [a_{ji}]$

(8) $A \in R^{m \times n} \Rightarrow A^t \in R^{n \times m}$, and $(AB)^t = B^t A^t$

LU-Decomposition

If a matrix A can be decomposed into the product of a unit lower- Δ matrix L and an upper- Δ matrix U , then the linear system $A\mathbf{x} = \mathbf{b}$ can be written as $LU\mathbf{x} = \mathbf{b}$. The problem is reduced to solving two simpler triangular systems $L\mathbf{y} = \mathbf{b}$ and $U\mathbf{x} = \mathbf{y}$ by forward and back substitutions.

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}, \quad L_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \quad L_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad L_1 = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Then

$$L_3 L_2 L_1 A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = U \Rightarrow A = L_1^{-1} L_2^{-1} L_3^{-1} U = LU$$

where

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = LU$$

◇ If $\mathbf{b} = [5, -2, 9]^t$, then $\mathbf{y} = [5, -12, 2]^t$ and $\mathbf{x} = [1, 1, 2]^t$

$$B = \begin{bmatrix} 2 & -2 & 3 \\ -2 & 3 & -4 \\ 4 & -3 & 7 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 1 \\ 0 \\ 5 \end{bmatrix} \Rightarrow \mathbf{x} = \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix}$$

Gaussian Elimination with Partial Pivoting

Not every matrix A (even if A is nonsingular) can be decomposed into the product of a unit lower- Δ matrix L and an upper- Δ matrix U by directly using Gaussian elimination. Whereas, for a nonsingular matrix A , Gaussian Elimination with Partial Pivoting will be introduced to overcome this problem later.

A linear system

$$\begin{aligned} y + 2z &= 0 \\ x + y + 2z &= 1 \\ 2x + 2y + z &= -1 \end{aligned}$$

A matrix representation

$$A\mathbf{x} = \mathbf{b}, \text{ or } \begin{bmatrix} 0 & 1 & 2 \\ 1 & 6 & 2 \\ 2 & 2 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ -1 \end{bmatrix}.$$

A corresponding augmented matrix for $A\mathbf{x} = \mathbf{b}$ is

$$[A|\mathbf{b}] = \left[\begin{array}{ccc|c} 0 & 1 & 2 & 0 \\ 1 & 6 & 2 & 1 \\ 2 & 2 & 1 & -1 \end{array} \right]$$

or equivalently

$$E_{13} \times [A|\mathbf{b}] = \left[\begin{array}{ccc|c} 2 & 2 & 1 & -1 \\ 1 & 6 & 2 & 1 \\ 0 & 1 & 2 & 0 \end{array} \right]$$

♣ Solution by using MATLAB

```
>> A = [0, 1, 2; 1, 1, 2; 2, 2, 1];
>> b = [0, 1, -1]';
>> x = A\b (x = [1; -2; 1])
```


Row Echelon Form

Definition: A matrix is said to be in row echelon form if

- (a) The first nonzero entry in each row is 1
- (b) If row k does not consist entirely of 0s, the number of leading zero entries in row $k+1$ is greater than that in row k
- (c) If there are rows whose entries are all zero, they are below the rows having nonzero entries

♣ *Example:* Matrices are in row echelon form

$$\begin{bmatrix} 1 & 4 & 2 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{bmatrix}, \quad \begin{bmatrix} 1 & 2 & 3 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}, \quad \begin{bmatrix} 1 & 3 & 0 & 1 \\ 0 & 0 & 1 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

♣ *Example:* Matrices which are not in row echelon form

$$\begin{bmatrix} 2 & 4 & 6 \\ 0 & 3 & 5 \\ 0 & 0 & 4 \end{bmatrix}, \quad \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix}, \quad \begin{bmatrix} 1 & 0 & 2 & 0 \\ 0 & 0 & 1 & 2 \\ 0 & 1 & 0 & 1 \end{bmatrix}$$

Computing An Inverse Matrix By Elementary Row Operations

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 1 & 3 & 3 \\ 2 & 4 & 7 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = I$$

$$E_1A = \begin{bmatrix} 1 & 2 & 3 \\ 0 & 1 & 0 \\ 2 & 4 & 7 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = E_1I$$

$$E_2E_1A = \begin{bmatrix} 1 & 2 & 3 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ -2 & 0 & 1 \end{bmatrix} = E_2E_1I$$

$$E_3E_2E_1A = \begin{bmatrix} 1 & 0 & 3 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 3 & -2 & 0 \\ -1 & 1 & 0 \\ -2 & 0 & 1 \end{bmatrix} = E_3E_2E_1I$$

$$E_4E_3E_2E_1A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 9 & -2 & -3 \\ -1 & 1 & 0 \\ -2 & 0 & 1 \end{bmatrix} = E_4E_3E_2E_1I = A^{-1}$$

where the elementary matrices are

$$E_1 = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad E_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -2 & 0 & 1 \end{bmatrix}, \quad E_3 = \begin{bmatrix} 1 & -2 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad E_4 = \begin{bmatrix} 1 & 0 & -3 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Computing Matrix Inverse

◇ Gauss-Jordan Method for Computing A^{-1} with $O(n^3)$

◇ By solving $A[\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n] = [\mathbf{e}_1, \mathbf{e}_2, \dots, \mathbf{e}_n]$

Example

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} \Rightarrow A^{-1} = \begin{bmatrix} \frac{3}{4} & \frac{-5}{16} & \frac{-3}{8} \\ \frac{1}{2} & \frac{-3}{8} & \frac{-1}{4} \\ -1 & 1 & 1 \end{bmatrix}$$

◇ The inverse of a unit lower- Δ matrix is also unit lower- Δ

$$L = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \frac{1}{4} & 1 & 0 & 0 \\ \frac{1}{3} & \frac{1}{3} & 1 & 0 \\ \frac{1}{2} & \frac{1}{2} & \frac{1}{2} & 1 \end{bmatrix} \Rightarrow L^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ -\frac{1}{4} & 1 & 0 & 0 \\ -\frac{1}{4} & -\frac{1}{3} & 1 & 0 \\ -\frac{1}{4} & -\frac{1}{2} & -\frac{1}{2} & 1 \end{bmatrix}$$

Some Special Matrices

$$A = [a_{ij}] \in R^{n \times n}$$

- *Diagonal* if $a_{ij} = 0 \forall i \neq j$
- *Lower - Δ* if $a_{ij} = 0$ if $j > i$
- *Unit lower - Δ* if A is lower- Δ with $a_{ii} = 1$
- *Lower Hessenberg* if $a_{ij} = 0$ for $j > i + 1$
- *Band matrix with bandwidth $2k + 1$* if $a_{ij} = 0$ for $|i - j| > k$

- A band matrix with *bandwidth 1* is *diagonal*
- A band matrix with *bandwidth 3* is *tridiagonal*
- A band matrix with *bandwidth 5* is *pentadiagonal*
- A lower and upper Hessenberg matrix is *tridiagonal*

$$A_1 = \begin{bmatrix} 7 & 0 & 0 \\ 1 & 8 & 0 \\ 2 & 3 & 9 \end{bmatrix}, \quad A_2 = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 3 & 4 & 1 \end{bmatrix}, \quad A_3 = \begin{bmatrix} 5 & 2 & 0 & 0 \\ 1 & 6 & 4 & 0 \\ 2 & 3 & 7 & 3 \\ 1 & 2 & 0 & 8 \end{bmatrix}, \quad A_4 = \begin{bmatrix} 5 & 2 & 0 & 0 \\ 1 & 6 & 4 & 0 \\ 0 & 3 & 7 & 3 \\ 0 & 0 & 2 & 8 \end{bmatrix}$$

Elementary Matrices (Row Operations)

(1) Interchange two rows $E_{ij} : R_i \longleftrightarrow R_j$

(2) Multiply a row by a nonzero real number $E_k(\alpha) : \alpha R_k$

(3) Replace a row by its sum with a multiple of another row $E_{ij}(\gamma) : \gamma R_i + R_j \longrightarrow R_j$

$$E_{12} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad E_3(-2) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -2 \end{bmatrix}, \quad E_{13}(1) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix},$$

♣ *Example*

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}, \quad E_{12}A = \begin{bmatrix} 4 & -6 & 0 \\ 2 & 1 & 1 \\ -2 & 7 & 2 \end{bmatrix}, \quad E_3(-2)A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ 4 & -14 & -4 \end{bmatrix}, \quad E_{13}(1)A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ 0 & 8 & 3 \end{bmatrix}$$

For

$$E_{23}(1) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \quad E_{13}(1) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad E_{12}(-2) = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}$$

We have

$$E_{23}(1)E_{13}(1)E_{12}(-2)A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = U \text{ (Upper-}\Delta) \Rightarrow A = [E_{12}(-2)]^{-1}[E_{13}(1)]^{-1}[E_{23}(1)]^{-1}U$$

LU-Decomposition

If a matrix A can be decomposed into the product of a unit lower- Δ matrix L and an upper- Δ matrix U , then the linear system $A\mathbf{x} = \mathbf{b}$ can be written as $LU\mathbf{x} = \mathbf{b}$. The problem is reduced to solving two simpler triangular systems $L\mathbf{y} = \mathbf{b}$ and $U\mathbf{x} = \mathbf{y}$ by forward and back substitutions.

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}, \quad E_{23}(1) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \quad E_{13}(1) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad E_{12}(-2) = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Then

$$E_{23}(1)E_{13}(1)E_{12}(-2)A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = U \Rightarrow A = [E_{12}(-2)]^{-1}[E_{13}(1)]^{-1}[E_{23}(1)]^{-1}U = LU$$

where

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = LU$$

◇ If $\mathbf{b} = [5, -2, 9]^t$ and use $B = [A|\mathbf{b}]$ instead of A , the above processes will obtain $\mathbf{y} = [5, -12, 2]^t$ and $\mathbf{x} = [1, 1, 2]^t$

- The inverse of a (unit) *lower* - Δ matrix is (unit) *lower* - Δ .
- The product of (unit) *lower* - Δ matrices is (unit) *lower* - Δ .

Analysis of Gaussian Elimination

♣ Algorithm

for $i = 1, 2, \dots, n - 1$

 for $k = i + 1, i + 2, \dots, n$

$m_{ki} \leftarrow a_{ki}/a_{ii}$ if $a_{ii} \neq 0$

$a_{ki} \leftarrow m_{ki}$

 for $j = i + 1, i + 2, \dots, n$

$a_{kj} \leftarrow a_{kj} - m_{ki} * a_{ij}$

 endfor

endfor

endfor

- The Worst Computational Complexity is $O(\frac{2}{3}n^3)$

1. # of divisions are $(n - 1) + (n - 2) + \dots + 1 = \frac{n(n-1)}{2}$

2. # of multiplications are $(n - 1)^2 + (n - 2)^2 + \dots + 1^2 = \frac{n(n-1)(2n-1)}{6}$

3. # of subtractions are $(n - 1)^2 + (n - 2)^2 + \dots + 1^2 = \frac{n(n-1)(2n-1)}{6}$

The Analysis of Gaussian Elimination and Back Substitution to solve $Ax=b$

$$\frac{2}{3}n^3 + \frac{3}{2}n^2 - \frac{7}{6}n$$

$$R_1 : a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1$$

$$R_2 : a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2$$

$$\vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots$$

$$R_i : a_{i1}x_1 + a_{i2}x_2 + \cdots + a_{in}x_n = b_i$$

$$\vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots$$

$$R_n : a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n$$

By Gaussian Elimination, we need $C_1 = [\sum_{k=1}^n (k+1)(k-1) + \sum_{k=1}^n k(k-1)]$ flops to reduce the above linear system of equations equivalent to the following upper triangular system.

$$R_1 : u_{11}x_1 + u_{12}x_2 + \cdots + u_{1n}x_n = c_1$$

$$\quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots$$

$$R_i : \quad \quad \quad u_{ii}x_i + \cdots + u_{in}x_n = c_i$$

$$\quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots$$

$$R_n : \quad \quad \quad \quad \quad \quad \quad u_{nn}x_n = c_n$$

We need $C_2 = \sum_{k=1}^n (2k-1) = n^2$ flops to solve an upper triangular linear system of equations. Therefore, the total number of flops of solving $Ax = b$ is summarized as

$$C_1 + C_2 = \frac{2}{3}n^3 + \frac{3}{2}n^2 - \frac{7}{6}n$$

PA=LU

Let $A \in R^{4 \times 4}$, by Gaussian Elimination with Partial Pivoting, we might have

$$L_3 P_3 L_2 P_2 L_1 P_1 A = U$$

where P_1 corresponds to $R_1 \longleftrightarrow R_4$, P_2 corresponds to $R_2 \longleftrightarrow R_4$, P_3 corresponds to $R_3 \longleftrightarrow R_4$, and

$$L_1 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \alpha_1 & 1 & 0 & 0 \\ \alpha_2 & 0 & 1 & 0 \\ \alpha_3 & 0 & 0 & 1 \end{bmatrix}, \quad L_2 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & \alpha_4 & 1 & 0 \\ 0 & \alpha_5 & 0 & 1 \end{bmatrix}, \quad L_3 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & \alpha_6 & 1 \end{bmatrix}$$

Here

$$P_1 = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}, \quad P_2 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}, \quad P_3 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

If we rewrite $L_3 P_3 L_2 P_2 L_1 P_1 A = U$ as

$$U = L_3 P_3 L_2 P_2 L_1 (P_2^{-1} P_2) P_1 A = L_3 P_3 L_2 (P_2 L_1 P_2^{-1}) P_2 P_1 A,$$

we have

$$P_2 L_1 P_2^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ \alpha_1 & 1 & 0 & 0 \\ \alpha_2 & 0 & 1 & 0 \\ \alpha_3 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \alpha_3 & 1 & 0 & 0 \\ \alpha_2 & 0 & 1 & 0 \\ \alpha_1 & 0 & 0 & 1 \end{bmatrix} = L_1^{(1)}$$

Note that $P_2^{-1} = P_2$. Similarly, $P_3^{-1} = P_3$, and thus we have

$$U = L_3 P_3 L_2 L_1^{(1)} P_2 P_1 A = L_3 (P_3 L_2 P_3) (P_3 L_1^{(1)} P_3) P_3 P_2 P_1 A = L_3 L_2^{(1)} L_1^{(2)} P_3 P_2 P_1 A$$

where

$$L_2^{(1)} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & \alpha_5 & 1 & 0 \\ 0 & \alpha_4 & 0 & 1 \end{bmatrix}, \quad L_1^{(2)} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \alpha_3 & 1 & 0 & 0 \\ \alpha_1 & 0 & 1 & 0 \\ \alpha_2 & 0 & 0 & 1 \end{bmatrix},$$

It follows that

$$PA = (P_3 P_2 P_1) A = (L_3 L_2^{(1)} L_1^{(2)})^{-1} U = (L_1^{(2)})^{-1} (L_2^{(1)})^{-1} L_3^{-1} U = LU$$

Here L is the *unit lower triangular matrix* (unit lower- Δ) and P is a permutation matrix given below.

$$L = \begin{bmatrix} 1 & 0 & 0 & 0 \\ -\alpha_3 & 1 & 0 & 0 \\ -\alpha_1 & -\alpha_5 & 1 & 0 \\ -\alpha_2 & -\alpha_4 & -\alpha_6 & 1 \end{bmatrix}, \quad P = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

Theorem: For any $A \in R^{n \times n}$, there exists a permutation matrix P such that $PA = LU$, where L is unit *lower* - Δ and U is *upper* - Δ .

Gaussian Elimination with partial Pivoting

♣ Algorithm

for $i = 1, 2, \dots, n$

$p(i) = i$

endfor

for $i = 1, 2, \dots, n - 1$

(a) select a pivotal element $a_{p(j),i}$ such that $|a_{p(j),i}| = \max_{i \leq k \leq n} |a_{p(k),i}|$

(b) $p(i) \longleftrightarrow p(j)$

(c) for $k = i + 1, i + 2, \dots, n$

$m_{p(k),i} = a_{p(k),i} / a_{p(i),i}$

 for $j = i + 1, i + 2, \dots, n$

$a_{p(k),j} = a_{p(k),j} - m_{p(k),i} * a_{p(i),j}$

 endfor

endfor

endfor

• An example

$$A = \begin{bmatrix} 0 & 9 & 1 \\ 1 & 2 & -2 \\ 2 & -5 & 4 \end{bmatrix} \Rightarrow P_{23}P_{13}A = LU = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ \frac{1}{2} & \frac{1}{2} & 1 \end{bmatrix} \begin{bmatrix} 2 & -5 & 4 \\ 0 & 9 & 1 \\ 0 & 0 & \frac{-9}{2} \end{bmatrix}$$

Matlab Codes for Gaussian Elimination with Partial Pivoting

```

% function [x,P]=GaussPP(A,b) - Solving Ax=b by PA=LUx=b,
%           Gaussian Elimination with partial pivoting
function [x,P]=GaussPP(A,b)
[m n]=size(A);
if (m~=n)
    error('matrix A must be square');
end
P=1:n;
Aug=[A, b];          % the augmented matrix
% Forward Elimination
for i=1:n-1
    kmax=i;
    t=abs(Aug(i,i));
    for k=i+1:n
        if (abs(Aug(k,i))>t)
            t=abs(A(k,i)); kmax=k;
        end
    end
    if (kmax~=i)
        tv=Aug(i,i:n+1);
        Aug(i,i:n+1)=Aug(kmax,i:n+1);
        Aug(kmax,i:n+1)=tv;
        kr=P(i); P(i)=P(kmax); P(kmax)=kr;
    end
    for k=i+1:n
        r=Aug(k,i)/Aug(i,i);
        Aug(k,i+1:n+1)=Aug(k,i+1:n+1)-r*Aug(i,i+1:n+1);
        Aug(k,i)=r;
    end
end
% Back Substitution
x=zeros(n,1);
x(n)=Aug(n,n+1)/Aug(n,n);
for j=n-1:-1:1
    x(j)=(Aug(j,n+1)-Aug(j,j+1:n)*x(j+1:n))/Aug(j,j);
end

```

Matlab Codes for Gaussian Elimination with Partial Pivoting

```

%%-----%%
%% gausspp.m - drive of Gaussian Elimination wit Partial Pivoting      %%
%%-----%%
fin=fopen('gaussmat.dat','r');
n=fscanf(fin,'%d',1);
A=fscanf(fin,'%f',[n n]);  A=A';
b=fscanf(fin,'%f',n);
X=gausspivot(A,b,n)

%%-----%%
%% gausspivot.m - Gaussian elimination with Partial Pivoting          %%
%%-----%%
function X=gausspivot(A,b,n)

if (abs(det(A))<eps)
    disp(sprintf('A is singular with det=%f\n',det(A)))
end

C=[A, b];
%----- Gaussian Elimination with Partial Pivoting -----%
for i=1:n-1
    [pivot, k]=max(abs(C(i:n,i)));
    if (k>1)
        temp=C(i,:);
        C(i,:)=C(i+k-1,:);
        C(i+k-1,:)=temp;
    end
    m(i+1:n,i)= -C(i+1:n,i)/C(i,i);
    C(i+1:n,:)=C(i+1:n,:)+m(i+1:n,i)*C(i,:);
end
%----- Back substitution -----%
X=zeros(n,1); %% Let X be a column vector of size n
X(n)=C(n,n+1)/C(n,n);
for i=n-1:-1:1
    X(i)=(C(i,n+1)-C(i,i+1:n)*X(i+1:n))/C(i,i);
end

```

Matlab Codes for A=LU Decomposition

```

% function [L,U]=LUdecomp(A) - A=LU decomposition
% Compute the LU decomposition of A such that (L,U) appeared in output A
function [L,U]=LUdecomp(A)
% A = [ 2,  1, -2;
%      -4, -1,  5;
%      2,  2,  2];
% b = [-2;  3;  0];
% x = [ 1; -2;  1]; Solution for A*x=b
[m n]=size(A);
if (m~=n)
    error('matrix A must be square');
end
A
% Forward Elimination
for i=1:n-1
    for k=i+1:n
        r=A(k,i)/A(i,i);
        A(k,i+1:n)=A(k,i+1:n)-r*A(i,i+1:n);
        A(k,i)=r;
    end
end
L=zeros(m,n);
U=zeros(m,n);
for i=1:m
    for j=1:n
        if (i>j)
            L(i,j)=A(i,j);
        elseif (i==j)
            L(i,j)=1;
            U(i,j)=A(i,j);
        elseif (i<j)
            U(i,j)=A(i,j);
        end
    end
end
end
OutputMatrixA=A
% 2  1 -2          1  0  0          2  1 -2
% -2  1  1, where L = -2  1  0 ; U = 0  1  1
% 1  1  3          1  1  1          0  0  3

```

Matlab Codes for $[L,U,P]=PALU(A)$

```

% function [L,U,P]=PALU(A) for (solving Gaussian Elimination with
%                               Partial Pivoting)
function [L,U,P]=PALU(A)
[m n]=size(A);
if (m~=n)
    error('matrix A must be square');
end
P=eye(n); Q=1:n-1;
% Forward Elimination with Partial Pivoting
for i=1:n-1
    t=abs(A(i,i));
    for k=i+1:n
        if (abs(A(k,i))>t)
            t=abs(A(k,i)); kmax=k;
        end
    end
    Q(i)=kmax;
    tv=A(i,i:n);
    A(i,i:n)=A(kmax,i:n);
    A(kmax,i:n)=tv;
    rowvec=P(i,:); P(i,:)=P(kmax,:); P(kmax,:)=rowvec; % compute P
    for k=i+1:n
        r=A(k,i)/A(i,i);
        A(k,i+1:n)=A(k,i+1:n)-r*A(i,i+1:n);
        A(k,i)=r;
    end
end
% Switch Q(j+1) with L(j) in PA=LU
for i=2:n-1
    k=Q(i);
    if (k~=i)
        rv=A(i,1:i-1);
        A(i,1:i-1)=A(k,1:i-1);
        A(k,1:i-1)=rv;
    end
end
L=eye(n); U=zeros(n,n);
for i=1:n
    for j=1:n

```

```
        if (i==j)
            U(i,j)=A(i,j);
        elseif (i>j)
            L(i,j)=A(i,j);
        elseif (i<j)
            U(i,j)=A(i,j);
        end
    end
end
end
Q
% A=[ 0  1  2  3;
%     1  2  4  5;
%     2 -2  3 -1;
%     4  2  1  0];
% 4.0000  2.0000  1.0000  0.0000
% 0.2500 -3.2500  2.5000 -1.0000
% 0.5000 -0.5000  5.0000  4.5000
% 0.0000 -0.3333  0.5667  0.1167
%-----
% P =[ 0 0 0 1;
%      0 0 1 0;
%      0 1 0 0;
%      1 0 0 0];
```