

# Problems of Eigenvalues/Eigenvectors

- ♣ Reveiw of Eigenvalues and Eigenvectors
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## Definition and Examples

Let  $A \in R^{n \times n}$ . If  $\exists \mathbf{v} \neq \mathbf{0}$  such that  $A\mathbf{v} = \lambda\mathbf{v}$ ,  $\lambda$  is called an eigenvalue of matrix  $A$ , and  $\mathbf{v}$  is called an eigenvector corresponding to (or belonging to) the eigenvalue  $\lambda$ . Note that  $\mathbf{v}$  is an eigenvector implies that  $\alpha\mathbf{v}$  is also an eigenvector for all  $\alpha \neq 0$ . We define the Eigenspace( $\lambda$ ) as the vector space spanned by all of the eigenvectors corresponding to the eigenvalue  $\lambda$ .

$$A\mathbf{x} = \lambda\mathbf{x} \Rightarrow (\lambda I - A)\mathbf{x} = \mathbf{0}, \mathbf{x} \neq \mathbf{0} \Rightarrow \det(\lambda I - A) = P(\lambda) = 0.$$

*Examples:*

$$\begin{aligned} 1. \quad A &= \begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix}, \lambda_1 = 2, \mathbf{u}_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \lambda_2 = 1, \mathbf{u}_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}. \\ 2. \quad A &= \begin{bmatrix} 2 & 1 \\ 0 & 1 \end{bmatrix}, \lambda_1 = 2, \mathbf{u}_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \lambda_2 = 1, \mathbf{u}_2 = \begin{bmatrix} -1 \\ 1 \end{bmatrix}. \\ 3. \quad A &= \begin{bmatrix} 3 & 1 \\ 1 & 3 \end{bmatrix}, \lambda_1 = 4, \mathbf{u}_1 = \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \lambda_2 = 2, \mathbf{u}_2 = \begin{bmatrix} -1 \\ 1 \end{bmatrix}. \\ 4. \quad A &= \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}, \lambda_1 = j, \mathbf{u}_1 = \begin{bmatrix} 1 \\ j \end{bmatrix}, \lambda_2 = -j, \mathbf{u}_2 = \begin{bmatrix} j \\ 1 \end{bmatrix}, j = \sqrt{-1}. \\ 5. \quad B &= \begin{bmatrix} 3 & 0 \\ 8 & -1 \end{bmatrix}, \text{ then } \lambda_1 = 3, \mathbf{u}_1 = \begin{bmatrix} \frac{1}{\sqrt{5}} \\ \frac{2}{\sqrt{5}} \end{bmatrix}; \lambda_2 = -1, \mathbf{u}_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}. \\ 6. \quad C &= \begin{bmatrix} 3 & -1 \\ -1 & 3 \end{bmatrix}, \text{ then } \tau_1 = 4, \mathbf{v}_1 = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix}; \tau_2 = 2, \mathbf{v}_2 = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix}. \end{aligned}$$

Note that  $\|\mathbf{u}_i\|_2 = 1$  and  $\|\mathbf{v}_i\|_2 = 1$  for  $i = 1, 2$ . Denote  $U = [\mathbf{u}_1, \mathbf{u}_2]$  and  $V = [\mathbf{v}_1, \mathbf{v}_2]$ , then

$$U^{-1}BU = \begin{bmatrix} 3 & 0 \\ 0 & -1 \end{bmatrix}, \quad V^{-1}CV = \begin{bmatrix} 4 & 0 \\ 0 & 2 \end{bmatrix}$$

Note that  $V^t = V^{-1}$  but  $U^t \neq U^{-1}$ .

$$\sum_{j=1}^n \lambda_j = \sum_{i=1}^n a_{ii} \text{ and } \prod_{j=1}^n \lambda_j = \det(A)$$

Let  $A \in R^{n \times n}$ , then  $P(\lambda) = \det(\lambda I - A)$  is called the *characteristic polynomial* of matrix  $A$ .

□ *Fundamental Theorem of Algebra*

A real polynomial  $P(\lambda) = \lambda^n + a_{n-1}\lambda^{n-1} + \dots + a_0$  of degree  $n$  has  $n$  roots  $\{\lambda_1, \lambda_2, \dots, \lambda_n\}$  such that

$$P(\lambda) = (\lambda - \lambda_1)(\lambda - \lambda_2) \cdots (\lambda - \lambda_n) = \lambda^n - \left( \sum_{i=1}^n \lambda_i \right) \lambda^{n-1} + \dots + (-1)^n \left( \prod_{i=1}^n \lambda_i \right)$$

- $\sum_{i=1}^n \lambda_i = \sum_{i=1}^n a_{ii} = \text{tr}(A)$  (called the trace of  $A$ )
- $\prod_{i=1}^n \lambda_i = \det(A)$

□ *Gershgorin's Disk Theorem*

Every eigenvalue of matrix  $A \in R^{n \times n}$  lies in at least one of the following disks

$$D_i = \{x \mid |x - a_{ii}| \leq \sum_{j \neq i} |a_{ij}|\}, \quad 1 \leq i \leq n$$

*Example:*  $B = \begin{bmatrix} 3 & 1 & 1 \\ 0 & 4 & 1 \\ 2 & 2 & 5 \end{bmatrix}$ ,  $\lambda_1, \lambda_2, \lambda_3 \in D_1 \cup D_2 \cup D_3$ , where

$$D_1 = \{z \mid |z - 3| \leq 2\}, \quad D_2 = \{z \mid |z - 4| \leq 1\}, \quad D_3 = \{z \mid |z - 5| \leq 4\}.$$

Note that  $\lambda_1 = 6.5616$ ,  $\lambda_2 = 3.0000$ ,  $\lambda_3 = 2.4383$ .

□ A matrix is said to be *diagonally dominant* if  $|a_{ii}| > \sum_{j \neq i} |a_{ij}|$ ,  $\forall 1 \leq i \leq n$ .

◇ A diagonally dominant matrix is invertible.

**Theorem:** Let  $A, P \in R^{n \times n}$ , with  $P$  nonsingular, then  $\lambda$  is an eigenvalue of  $A$  with eigenvector  $\mathbf{x}$  iff  $\lambda$  is an eigenvalue of  $P^{-1}AP$  with eigenvector  $P^{-1}\mathbf{x}$ .

**(Proof)** Let  $\mathbf{x}$  be an eigenvector of  $A$  corresponding to the eigenvalue  $\lambda$ , that is,  $A\mathbf{x} = \lambda\mathbf{x}$ . Then, we have

$$(P^{-1}AP)(P^{-1}\mathbf{x}) = P^{-1}A(PP^{-1})\mathbf{x} = P^{-1}A\mathbf{x} = P^{-1}(\lambda\mathbf{x}) = \lambda(P^{-1}\mathbf{x})$$

Thus,  $P^{-1}\mathbf{x}$  is an eigenvector corresponding to the eigenvalue  $\lambda$  of the matrix  $P^{-1}AP$  (according to the definition).

On the other hand,

$$(P^{-1}AP)(P^{-1}\mathbf{x}) = \lambda(P^{-1}\mathbf{x})$$

implies that  $A\mathbf{x} = \lambda\mathbf{x}$  could be achieved based on simple matrix operations.

**Theorem:** Let  $A \in R^{n \times n}$  and let  $\lambda$  be an eigenvalue of  $A$  with eigenvector  $\mathbf{x}$ . Then

- (a)  $\alpha\lambda$  is an eigenvalue of matrix  $\alpha A$  with eigenvector  $\mathbf{x}$
- (b)  $\lambda - \mu$  is an eigenvalue of matrix  $A - \mu I$  with eigenvector  $\mathbf{x}$
- (c) If  $A$  is nonsingular, then  $\lambda \neq 0$  and  $\lambda^{-1}$  is an eigenvalue of  $A^{-1}$  with eigenvector  $\mathbf{x}$

Let  $\mathbf{x}$  be an eigenvector of  $A$  corresponding to the eigenvalue  $\lambda$ , that is,  $A\mathbf{x} = \lambda\mathbf{x}$ . Then

**Proof of (a)**  $(\alpha A)\mathbf{x} = \alpha(A\mathbf{x}) = \alpha(\lambda\mathbf{x}) = (\alpha\lambda)\mathbf{x}$ .

**Proof of (b)**  $(A - \mu I)\mathbf{x} = A\mathbf{x} - \mu\mathbf{x} = \lambda\mathbf{x} - \mu\mathbf{x} = (\lambda - \mu)\mathbf{x}$ .

**Proof of (c)** If  $A$  is nonsingular, none of its eigenvalues is zero, otherwise,  $A\mathbf{x} = \lambda\mathbf{x} = 0 \cdot \mathbf{x} = \mathbf{0}$  and  $\mathbf{x} = A^{-1}\mathbf{0} = \mathbf{0}$  which implies that  $\mathbf{x} = \mathbf{0}$  that contradicts that  $\mathbf{x}$  is an eigenvector (of  $A$ ). Then,  $A\mathbf{x} = \lambda\mathbf{x}$  implies that  $\frac{1}{\lambda}\mathbf{x} = A^{-1}\mathbf{x}$ . Therefore,  $\frac{1}{\lambda}$  is an eigenvalue of matrix  $A^{-1}$  with eigenvector  $\mathbf{x}$ .

**Definition:** A matrix  $A$  is similar to  $B$ , denote by  $A \sim B$ , iff there exists an invertible matrix  $U$  such that  $U^{-1}AU = B$ . Furthermore, a matrix  $A$  is *orthogonally similar* to  $B$ , iff there exists an orthogonal matrix  $Q$  such that  $Q^tAQ = B$ .

**Theorem:** Two similar matrices have the same eigenvalues, i.e.,  $A \sim B \Rightarrow \lambda(A) = \lambda(B)$ .

**Proof** Since  $A \sim B$ , we have  $B = U^{-1}AU$  for some  $U$ , then

$$|\lambda I - B| = |U^{-1}(\lambda I)U - U^{-1}AU| = |U^{-1}(\lambda I - A)U| = |U^{-1}| \cdot |\lambda I - A| \cdot |U| = |U|^{-1} \cdot |\lambda I - A| \cdot |U|$$

## Diagonalization of Matrices

**Theorem:** Suppose  $A \in R^{n \times n}$  has  $n$  linearly independent eigenvectors  $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n$  corresponding to eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_n$ . Let  $V = [\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n]$ , then  $V^{-1}AV = \text{diag}[\lambda_1, \lambda_2, \dots, \lambda_n]$ .

- ◇ If  $A \in R^{n \times n}$  has  $n$  distinct eigenvalues, then their corresponding eigenvectors are linearly independent. Thus, any matrix with distinct eigenvalues can be diagonalized.
- ◇ Not all matrices have distinct eigenvalues, therefore not all matrices are diagonalizable.

### Nondiagonalizable Matrices

$$A = \begin{bmatrix} 2 & 1 & 0 \\ 0 & 2 & 1 \\ 0 & 0 & 2 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 2 & 0 \\ -3 & 5 & 2 \end{bmatrix}$$

### Diagonalizable Matrices

$$C = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}, \quad D = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}, \quad E = \begin{bmatrix} 0 & 0 & -2 \\ 1 & 2 & 1 \\ 1 & 0 & 3 \end{bmatrix}, \quad K = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$$

**Spectrum Decomposition Theorem:** Every real symmetric matrix can be orthogonally diagonalized.

- ◇  $U^tAU = \Lambda$  or  $A = U\Lambda U^t = \sum_{i=1}^n \lambda_i \mathbf{u}_i \mathbf{u}_i^t$ , where  $U$  is an orthogonal matrix, and  $\Lambda = \text{diag}[\lambda_1, \lambda_2, \dots, \lambda_n]$ .

## Similarity transformation and triangularization

**Schur's Theorem:**  $\forall A \in R^{n \times n}$ ,  $\exists$  an orthogonal matrix  $U$  such that  $U^t A U = T$  is upper- $\Delta$ . The eigenvalues must be shared by the similarity matrix  $T$  and appear along its main diagonal.

**Hint:** By induction, suppose that the theorem has been proved for all matrices of order  $n - 1$ , and consider  $A \in R^{n \times n}$  with  $A\mathbf{x} = \lambda\mathbf{x}$  and  $\|\mathbf{x}\|_2 = 1$ , then  $\exists$  a Householder matrix  $H_1$  such that  $H_1\mathbf{x} = \beta\mathbf{e}_1$ , e.g.,  $\beta = -\|\mathbf{x}\|_2$ , hence

$$H_1 A H_1^t \mathbf{e}_1 = H_1 A (H_1^{-1} \mathbf{e}_1) = H_1 A (\beta^{-1} \mathbf{x}) = H_1 \beta^{-1} A \mathbf{x} = \beta^{-1} \lambda (H_1 \mathbf{x}) = \beta^{-1} \lambda (\beta \mathbf{e}_1) = \lambda \mathbf{e}_1$$

Thus,

$$H_1 A H_1^t = \left[ \begin{array}{c|c} \lambda & * \\ \hline - & - \\ O & A^{(1)} \end{array} \right]$$

**Spectrum Decomposition Theorem:** Every real symmetric matrix can be orthogonally diagonalized.

$\diamond U^t A U = \Lambda$  or  $A = U \Lambda U^t = \sum_{i=1}^n \lambda_i \mathbf{u}_i \mathbf{u}_i^t$ , where  $U$  is an orthogonal matrix, and  $\Lambda = \text{diag}[\lambda_1, \lambda_2, \dots, \lambda_n]$ .

**Definition:** A symmetric matrix  $A \in R^{n \times n}$  is nonnegative definite if  $\mathbf{x}^t A \mathbf{x} \geq 0 \forall \mathbf{x} \in R^n$ ,  $\mathbf{x} \neq \mathbf{0}$ .

**Definition:** A symmetric matrix  $A \in R^{n \times n}$  is positive definite if  $\mathbf{x}^t A \mathbf{x} > 0 \forall \mathbf{x} \in R^n$ ,  $\mathbf{x} \neq \mathbf{0}$ .

**Singular Value Decomposition Theorem:** Each matrix  $A \in R^{m \times n}$  can be decomposed as  $A = U \Sigma V^t$ , where both  $U \in R^{m \times m}$  and  $V \in R^{n \times n}$  are orthogonal. Moreover,  $\Sigma \in R^{m \times n} = \text{diag}[\sigma_1, \sigma_2, \dots, \sigma_k, 0, \dots, 0]$  is essentially diagonal with the singular values satisfying  $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_k > 0$ .

$\diamond A = U \Sigma V^t = \sum_{i=1}^k \sigma_i \mathbf{u}_i \mathbf{v}_i^t$

*Example:*

$$A = \begin{bmatrix} 2 & -1 \\ -1 & 2 \end{bmatrix}, \quad B = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad C = \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 0 & 0 \end{bmatrix}$$

## A Jacobi Transform (Givens Rotation)

$$J(i, k; \theta) = \begin{bmatrix} 1 & \cdot & \cdot & \cdots & \cdot & \cdot & 0 \\ 0 & \ddots & \cdot & \cdots & \cdot & \vdots & 0 \\ 0 & \cdot & c & \cdots & s & \cdot & 0 \\ \cdot & \vdots & \cdot & \ddots & \cdot & \vdots & \cdot \\ 0 & \cdot & -s & \cdots & c & \cdot & 0 \\ 0 & \vdots & \cdot & \cdots & \cdot & \ddots & 0 \\ \cdot & \cdot & 0 & \cdots & 0 & \cdot & 1 \end{bmatrix}$$

$$J_{hh} = 1 \text{ if } h \neq i \text{ or } h \neq k, \text{ where } i < k$$

$$J_{ii} = J_{kk} = c = \cos \theta$$

$$J_{ki} = -s = -\sin \theta, \quad J_{ik} = s = \sin \theta$$

Let  $\mathbf{x}, \mathbf{y} \in R^n$ , then  $\mathbf{y} = J(i, k; \theta)\mathbf{x}$  implies that

$$y_i = cx_i + sx_k$$

$$y_k = -sx_i + cx_k$$

$$c = \frac{x_i}{\sqrt{x_i^2 + x_k^2}}, \quad s = \frac{x_k}{\sqrt{x_i^2 + x_k^2}},$$

$$\mathbf{x} = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \end{bmatrix}, \quad \begin{bmatrix} \cos \theta \\ \sin \theta \end{bmatrix} = \begin{bmatrix} 1/\sqrt{5} \\ 2/\sqrt{5} \end{bmatrix}, \quad \text{then } J(2, 4; \theta)\mathbf{x} = \begin{bmatrix} 1 \\ \sqrt{20} \\ 3 \\ 0 \end{bmatrix}$$

## Jacobi Transforms (Givens Rotations)

The Jacobi method consists of a sequence of orthogonal similarity transformations such that

$$J_K^t J_{K-1}^t \cdots J_2^t J_1^t A J_1 J_2 \cdots J_{K-1} J_K = \Lambda$$

where each  $J_i$  is orthogonal, so is  $Q = J_1 J_2 \cdots J_{K-1} J_K$ .

Each Jacobi transform (Given rotation) is just a plane rotation designed to annihilate one of the off-diagonal matrix elements. Let  $A = (a_{ij})$  be symmetric, then

$B = J^t(p, q, \theta) A J(p, q, \theta)$ , where

$$b_{rp} = ca_{rp} - sa_{rq} \quad \text{for } r \neq p, r \neq q$$

$$b_{rq} = sa_{rp} + ca_{rq} \quad \text{for } r \neq p, r \neq q$$

$$b_{pp} = c^2 a_{pp} + s^2 a_{qq} - 2sca_{pq}$$

$$b_{qq} = s^2 a_{pp} + c^2 a_{qq} + 2sca_{pq}$$

$$b_{pq} = (c^2 - s^2)a_{pq} + sc(a_{pp} - a_{qq})$$

To set  $b_{pq} = 0$ , we choose  $c, s$  such that

$$\alpha = \cot(2\theta) = \frac{c^2 - s^2}{2sc} = \frac{a_{qq} - a_{pp}}{2a_{pq}} \quad (1)$$

For computational convenience, let  $t = \frac{s}{c}$ , then  $t^2 + 2\alpha t - 1 = 0$  whose smaller root (in absolute sense) can be computed by

$$t = \frac{\text{sgn}(\alpha)}{\sqrt{\alpha^2 + 1} + |\alpha|}, \quad \text{and } c = \frac{1}{\sqrt{1 + t^2}}, \quad s = ct, \quad \tau = \frac{s}{1 + c} \quad (2)$$

Remark

$$b_{pp} = a_{pp} - ta_{pq}$$

$$b_{qq} = a_{qq} + ta_{pq}$$

$$b_{rp} = a_{rp} - s(a_{rq} + \tau a_{rp})$$

$$b_{rq} = a_{rq} + s(a_{rp} - \tau a_{rq})$$



## Algorithm of Jacobi Transforms to Diagonalize A

$$A^{(0)} \leftarrow A$$

for  $k = 0, 1, \dots$ , until convergence

$$\text{Let } |a_{pq}^{(k)}| = \text{Max}_{i < j} \{|a_{ij}^{(k)}|\}$$

Compute

$$\alpha_k = \frac{a_{qq}^{(k)} - a_{pp}^{(k)}}{2a_{pq}^{(k)}}, \text{ solve } \cot(2\theta_k) = \alpha_k \text{ for } \theta_k.$$

$$t = \frac{\text{sgn}(\alpha)}{\sqrt{\alpha^2 + 1 + |\alpha|}}$$

$$c = \frac{1}{\sqrt{1+t^2}}, \quad s = ct$$

$$\tau = \frac{s}{1+c}$$

$$A^{(k+1)} \leftarrow J_k^t A^{(k)} J_k, \text{ where } J_k = J(p, q, \theta_k)$$

endfor

# Convergence of Jacobi Algorithm to Diagonalize A

**Proof:**

Since  $|a_{pq}^{(k)}| \geq |a_{ij}^{(k)}|$  for  $i \neq j$ ,  $p \neq q$ , then

$|a_{pq}^{(k)}|^2 \geq \text{off}(A^{(k)})/2N$ , where  $N = \frac{n(n-1)}{2}$ , and

$\text{off}(A^{(k)}) = \sum_{i \neq j}^n (a_{ij}^{(k)})^2$ , the sum of square off-diagonal elements of  $A^{(k)}$

Furthermore,

$$\begin{aligned} \text{off}(A^{(k+1)}) &= \text{off}(A^{(k)}) - 2(a_{pq}^{(k)})^2 + 2(a_{pq}^{(k+1)})^2 \\ &= \text{off}(A^{(k)}) - 2(a_{pq}^{(k)})^2, \text{ since } a_{pq}^{(k+1)} = 0 \\ &\leq \text{off}(A^{(k)}) \left(1 - \frac{1}{N}\right), \text{ since } |a_{pq}^{(k)}|^2 \geq \text{off}(A^{(k)})/2N \end{aligned}$$

Thus

$$\text{off}(A^{(k+1)}) \leq \left(1 - \frac{1}{N}\right)^{k+1} \text{off}(A^{(0)}) \rightarrow 0 \text{ as } k \rightarrow \infty$$

Example:

$$A = \begin{bmatrix} 4 & 2 & 0 \\ 2 & 3 & 1 \\ 0 & 1 & 2 \end{bmatrix}, \quad J(1, 2; \theta) = \begin{bmatrix} c & s & 0 \\ -s & c & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Then

$$A^{(1)} = J^t(1, 2; \theta)AJ(1, 2; \theta) = \begin{bmatrix} 4c^2 - 4cs + 3s^2 & 2c^2 + cs - 2s^2 & -s \\ 2c^2 + cs - 2s^2 & 3c^2 + 4cs + 4s^2 & c \\ -s & c & 1 \end{bmatrix}$$

Note that  $\text{off}(A^{(1)}) = 2 < 10 = \text{off}(A^{(0)}) = \text{off}(A)$

## Example for Convergence of Jacobi Algorithm

$$A^{(0)} = \begin{bmatrix} 1.0000 & 0.5000 & 0.2500 & 0.1250 \\ 0.5000 & 1.0000 & 0.5000 & 0.2500 \\ 0.2500 & 0.5000 & 1.0000 & 0.5000 \\ 0.1250 & 0.2500 & 0.5000 & 1.0000 \end{bmatrix}, \quad A^{(1)} = \begin{bmatrix} 1.5000 & 0.0000 & 0.5303 & 0.2652 \\ 0.0000 & 0.5000 & 0.1768 & 0.0884 \\ 0.5303 & 0.1768 & 1.0000 & 0.5000 \\ 0.2652 & 0.0884 & 0.5000 & 1.0000 \end{bmatrix}$$

$$A^{(2)} = \begin{bmatrix} 1.8363 & 0.0947 & 0.0000 & 0.4917 \\ 0.0947 & 0.5000 & 0.1493 & 0.0884 \\ 0.0000 & 0.1493 & 0.6637 & 0.2803 \\ 0.4917 & 0.0884 & 0.2803 & 1.0000 \end{bmatrix}, \quad A^{(3)} = \begin{bmatrix} 2.0636 & 0.1230 & 0.1176 & 0.0000 \\ 0.1230 & 0.5000 & 0.1493 & 0.0405 \\ 0.1176 & 0.1493 & 0.6637 & 0.2544 \\ 0.0000 & 0.0405 & 0.2544 & 0.7727 \end{bmatrix}$$

$$A^{(4)} = \begin{bmatrix} 2.0636 & 0.1230 & 0.0915 & 0.0739 \\ 0.1230 & 0.5000 & 0.0906 & 0.1254 \\ 0.0915 & 0.0906 & 0.4580 & 0.0000 \\ 0.0739 & 0.1254 & 0.0000 & 0.9783 \end{bmatrix}, \quad A^{(5)} = \begin{bmatrix} 2.0636 & 0.1018 & 0.0915 & 0.1012 \\ 0.1018 & 0.4691 & 0.0880 & 0.0000 \\ 0.0915 & 0.0880 & 0.4580 & 0.0217 \\ 0.1012 & 0.0000 & 0.0217 & 1.0092 \end{bmatrix}$$

$$A^{(6)} = \begin{bmatrix} 2.0701 & 0.0000 & 0.0969 & 0.1010 \\ 0.0000 & 0.4627 & 0.0820 & -0.0064 \\ 0.0969 & 0.0820 & 0.4580 & 0.0217 \\ 0.1010 & -0.0064 & 0.0217 & 1.0092 \end{bmatrix}, \quad A^{(15)} = \begin{bmatrix} 2.0856 & 0.0000 & 0.0000 & 0.0000 \\ 0.0000 & 0.5394 & 0.0000 & -0.0000 \\ 0.0000 & 0.0000 & 0.3750 & 0.0000 \\ 0.0000 & -0.0000 & 0.0000 & 1.0000 \end{bmatrix}$$

## Cholesky Algorithm

□ **Theorem:** Every positive definite matrix  $A$  can be decomposed as  $A = LL^t$ , where  $L$  is lower -  $\Delta$ .

□ **Algorithm:**  $A \in R^{n \times n}$ ,  $A = LL^t$ ,  $A$  is positive definite and  $L$  is lower -  $\Delta$ .

for  $j = 0, 1, \dots, n - 1$

$$L_{jj} \leftarrow [A_{jj} - \sum_{k=0}^{j-1} L_{jk}^2]^{1/2}$$

for  $i = j + 1, j + 2, \dots, n - 1$

$$L_{ij} \leftarrow [A_{ij} - \sum_{k=0}^{j-1} L_{ik}L_{jk}] / L_{jj}$$

endfor

endfor

$$C = \begin{bmatrix} 4 & -2 \\ -2 & 5 \end{bmatrix} = \begin{bmatrix} 2 & 0 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} 2 & -1 \\ 0 & 2 \end{bmatrix} = L_1 L_1^t$$

$$A = \begin{bmatrix} 9 & 3 & -3 \\ 3 & 17 & 3 \\ -3 & 3 & 27 \end{bmatrix} = \begin{bmatrix} 3 & 0 & 0 \\ 1 & 4 & 0 \\ -1 & 1 & 5 \end{bmatrix} \begin{bmatrix} 3 & 1 & -1 \\ 0 & 4 & 1 \\ 0 & 0 & 5 \end{bmatrix} = L_2 L_2^t$$

## Power of A Matrix and Its Eigenvalues

**Theorem:** Let  $\lambda_1, \lambda_2, \dots, \lambda_n$  be eigenvalues of  $A \in R^{n \times n}$ . Then  $\lambda_1^k, \lambda_2^k, \dots, \lambda_n^k$  are eigenvalues of  $A^k \in R^{n \times n}$  with the same corresponding eigenvectors of  $A$ . That is,

$$A\mathbf{v}_i = \lambda_i\mathbf{v}_i \quad \rightarrow \quad A^k\mathbf{v}_i = \lambda_i^k\mathbf{v}_i \quad \forall 1 \leq i \leq n$$

Suppose that the matrix  $A \in R^{n \times n}$  has  $n$  linearly independent eigenvectors  $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n$  corresponding to eigenvalues  $\lambda_1, \lambda_2, \dots, \lambda_n$ . Then any  $\mathbf{x} \in R^n$  can be written as

$$\mathbf{x} = c_1\mathbf{v}_1 + c_2\mathbf{v}_2 + \dots + c_n\mathbf{v}_n$$

Then

$$A^k\mathbf{x} = \lambda_1^k c_1\mathbf{v}_1 + \lambda_2^k c_2\mathbf{v}_2 + \dots + \lambda_n^k c_n\mathbf{v}_n$$

In particular, if  $|\lambda_1| > |\lambda_j|$  for  $2 \leq j \leq n$  and  $c_1 \neq 0$ , then  $A^k\mathbf{x}$  will tend to lie in the direction  $\mathbf{v}_1$  when  $k$  is *large enough*.

## Power Method for Computing the Largest Eigenvalues

Suppose that the matrix  $A \in R^{n \times n}$  is diagonalizable and that  $U^{-1}AU = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_n)$  with  $U = [\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n]$  and  $|\lambda_1| > |\lambda_2| \geq |\lambda_3| \geq \dots \geq |\lambda_n|$ . Given  $\mathbf{u}^{(0)} \in R^n$ , then power method produces a sequence of vectors  $\mathbf{u}^{(k)}$  as follows.

for  $k = 1, 2, \dots$

$$\mathbf{z}^{(k)} = A\mathbf{u}^{(k-1)}$$

$$r^{(k)} = z_m^{(k)} = \|\mathbf{z}^{(k)}\|_\infty, \text{ for some } 1 \leq m \leq n.$$

$$\mathbf{u}^{(k)} = \mathbf{z}^{(k)} / r^{(k)}$$

endfor

$\lambda_1$  must be real since the complex eigenvalues must appear in a "relatively conjugate pair".

$$A = \begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix} \Rightarrow \begin{array}{l} \lambda_1 = 3 \\ \lambda_2 = 1 \end{array}, \quad \mathbf{v}_1 = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \quad \mathbf{v}_2 = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

Let  $\mathbf{u}^{(0)} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$ , then  $\mathbf{u}^{(5)} = \begin{bmatrix} 1.0 \\ 0.9918 \end{bmatrix}$ , and  $r^{(5)} = 2.9756$ .

## QR Iterations for Computing Eigenvalues

```

%
% Script File: shiftQR.m
% Solving Eigenvalues by shift-QR factorization
%
Nrun=15;
fin=fopen('dataMatrix.txt');
fgetL(fin); % read off the header line
n=fscanf(fin,'%d',1);
A=fscanf(fin,'%f',[n n]);
A=A';
SaveA=A;
for k=1:Nrun,
    s=A(n,n);
    A=A-s*eye(n);
    [Q R]=qr(A);
    A=R*Q+s*eye(n);
end
eig(SaveA)
%
% dataMatrix.txt
%
Matrices for computing eigenvalues by QR factorization or shift-QR
5
1.0    0.5    0.25  0.125  0.0625
0.5    1.0    0.5    0.25   0.125
0.25   0.5    1.0    0.5    0.25
0.125  0.25   0.5    1.0    0.5
0.0625 0.125  0.25   0.5    1.0
4
          for shift-QR studies
2.9766  0.3945  0.4198  1.1159
0.3945  2.7328 -0.3097  0.1129
0.4198 -0.3097  2.5675  0.6079
1.1159  0.1129  0.6097  1.7231

```

## Norms of Vectors and Matrices

**Definition:** A vector norm on  $R^n$  is a function

$$\tau : R^n \rightarrow R^+ = \{x \geq 0 \mid x \in R\}$$

that satisfies

$$(1) \tau(\mathbf{x}) > 0 \quad \forall \mathbf{x} \neq \mathbf{0}, \quad \tau(\mathbf{0}) = 0$$

$$(2) \tau(c\mathbf{x}) = |c|\tau(\mathbf{x}) \quad \forall c \in R, \quad \mathbf{x} \in R^n$$

$$(3) \tau(\mathbf{x} + \mathbf{y}) \leq \tau(\mathbf{x}) + \tau(\mathbf{y}) \quad \forall \mathbf{x}, \mathbf{y} \in R^n$$

*Hölder norm (p-norm)*  $\|\mathbf{x}\|_p = (\sum_{i=1}^n |x_i|^p)^{1/p}$  for  $p \geq 1$ .

$$(\mathbf{p}=1) \quad \|\mathbf{x}\|_1 = \sum_{i=1}^n |x_i| \quad (\text{Mahattan or City-block distance})$$

$$(\mathbf{p}=2) \quad \|\mathbf{x}\|_2 = (\sum_{i=1}^n |x_i|^2)^{1/2} \quad (\text{Euclidean distance})$$

$$(\mathbf{p}=\infty) \quad \|\mathbf{x}\|_\infty = \max_{1 \leq i \leq n} \{|x_i|\} \quad (\infty\text{-norm})$$



**Definition:** A matrix norm on  $R^{m \times n}$  is a function

$$\tau : R^{m \times n} \rightarrow R^+ = \{x \geq 0 \mid x \in R\}$$

that satisfies

- (1)  $\tau(A) > 0 \quad \forall A \neq O, \tau(O) = 0$
- (2)  $\tau(cA) = |c|\tau(A) \quad \forall c \in R, A \in R^{m \times n}$
- (3)  $\tau(A + B) \leq \tau(A) + \tau(B) \quad \forall A, B \in R^{m \times n}$

*Consistency Property:*  $\tau(AB) \leq \tau(A)\tau(B) \quad \forall A, B$

- (a)  $\tau(A) = \max\{|a_{ij}| \mid 1 \leq i \leq m, 1 \leq j \leq n\}$
- (b)  $\|A\|_F = \left[ \sum_{i=1}^m \sum_{j=1}^n a_{ij}^2 \right]^{1/2}$  (Fröbenius norm)

**Subordinate Matrix Norm:**  $\|A\| = \max_{\|\mathbf{x}\| \neq 0} \{\|A\mathbf{x}\| / \|\mathbf{x}\|\}$

- (1) If  $A \in R^{m \times n}$ , then  $\|A\|_1 = \max_{1 \leq j \leq n} (\sum_{i=1}^m |a_{ij}|)$
- (2) If  $A \in R^{m \times n}$ , then  $\|A\|_\infty = \max_{1 \leq i \leq m} (\sum_{j=1}^n |a_{ij}|)$
- (3) Let  $A \in R^{n \times n}$  be real symmetric, then  $\|A\|_2 = \max_{1 \leq i \leq n} |\lambda_i|$ , where  $\lambda_i \in \lambda(A)$

**Theorem:** Let  $\mathbf{x} \in R^n$  and let  $A = (a_{ij}) \in R^{n \times n}$ . Define  $\|A\|_1 = \text{Sup}_{\|\mathbf{u}\|_1=1} \{\|A\mathbf{u}\|_1\}$

**Proof:** For  $\|\mathbf{u}\|_1 = 1$ ,

$$\|A\|_1 = \text{Sup}\{\|A\mathbf{u}\|_1\} = \sum_{i=1}^n \left| \sum_{j=1}^n a_{ij} u_j \right| \leq \sum_{j=1}^n \sum_{i=1}^n |a_{ij}| |u_j| = \sum_{j=1}^n |u_j| \sum_{i=1}^n |a_{ij}|$$

Then

$$\|A\|_1 \leq \text{Max}_{1 \leq j \leq n} \left\{ \sum_{i=1}^n |a_{ij}| \right\} \sum_{j=1}^n |u_j| = \text{Max}_{1 \leq j \leq n} \left\{ \sum_{i=1}^n |a_{ij}| \right\}$$

On the other hand, let  $\sum_{i=1}^n |a_{ik}| = \text{Max}_{1 \leq j \leq n} \left\{ \sum_{i=1}^n |a_{ij}| \right\}$  and choose  $\mathbf{u} = \mathbf{e}_k$ , which completes the proof.

**Theorem:** Let  $A = [a_{ij}] \in R^{m \times n}$ , and define  $\|A\|_\infty = \text{Max}_{\|\mathbf{u}\|_\infty=1} \{\|A\mathbf{u}\|_\infty\}$ .

$$\text{Show that } \|A\|_\infty = \text{Max}_{1 \leq i \leq m} \left\{ \sum_{j=1}^n |a_{ij}| \right\}$$

**Proof:** Let  $\sum_{j=1}^n |a_{Kj}| = \text{Max}_{1 \leq i \leq m} \left\{ \sum_{j=1}^n |a_{ij}| \right\}$ , for any  $\mathbf{x} \in R^n$  with  $\|\mathbf{x}\|_\infty = 1$ , we have

$$\begin{aligned} \|A\mathbf{x}\|_\infty &= \text{Max}_{1 \leq i \leq m} \left\{ \left| \sum_{j=1}^n a_{ij} x_j \right| \right\} \\ &\leq \text{Max}_{1 \leq i \leq m} \left\{ \sum_{j=1}^n |a_{ij}| \cdot |x_j| \right\} \leq \text{Max}_{1 \leq i \leq m} \left\{ \sum_{j=1}^n |a_{ij}| \|\mathbf{x}\|_\infty \right\} \\ &\leq \text{Max}_{1 \leq i \leq m} \left\{ \sum_{j=1}^n |a_{ij}| \right\} = \sum_{j=1}^n |a_{Kj}| \end{aligned}$$

In particular, if we pick up  $\mathbf{y} \in R^n$  such that  $y_j = \text{sign}(a_{Kj})$ ,  $\forall 1 \leq j \leq n$ , then  $\|\mathbf{y}\|_\infty = 1$ , and  $\|A\mathbf{y}\|_\infty = \sum_{j=1}^n |a_{Kj}|$ , which completes the proof.

**Theorem:** Let  $A = [a_{ij}] \in R^{n \times n}$ , and define  $\|A\|_2 = \text{Max}_{\|\mathbf{x}\|_2=1} \{\|A\mathbf{x}\|_2\}$ . Show that

$$\|A\|_2 = \sqrt{\rho(A^t A)} = \sqrt{\text{maximum eigenvalue of } A^t A} \quad (\text{spectral radius})$$

**(Proof)** Let  $\lambda_1, \lambda_2, \dots, \lambda_n$  be eigenvalues and their corresponding unit eigenvectors  $\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_n$  of matrix  $A^t A$ , that is,

$$(A^t A)\mathbf{u}_i = \lambda_i \mathbf{u}_i \quad \text{and} \quad \|\mathbf{u}_i\|_2 = 1 \quad \forall 1 \leq i \leq n.$$

Since  $\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_n$  must be an orthonormal basis based on *spectrum decomposition*

*theorem*, for any  $\mathbf{x} \in R^n$ , we have  $\mathbf{x} = \sum_{i=1}^n c_i \mathbf{u}_i$ . Then

$$\begin{aligned} \|A\|_2 &= \text{Max}_{\|\mathbf{x}\|_2=1} \{\|A\mathbf{x}\|_2\} \\ &= \sqrt{\text{Max}_{\|\mathbf{x}\|_2=1} \{\|A\mathbf{x}\|_2^2\}} \\ &= \sqrt{\text{Max}_{\|\mathbf{x}\|_2=1} \{\mathbf{x}^t A^t A \mathbf{x}\}} \\ &= \sqrt{\text{Max}_{\|\mathbf{x}\|_2=1} \left| \sum_{i=1}^n \lambda_i c_i^2 \right|} \\ &= \sqrt{\text{Max}_{1 \leq j \leq n} \{|\lambda_j|\}} \end{aligned}$$

## A Markov Process

Suppose that 10% of the people outside Taiwan move in, and 20% of the people inside Taiwan move out in each year. Let  $y_k$  and  $z_k$  be the population at the end of the  $k$ -th year, outside Taiwan and inside Taiwan, respectively. Then we have

$$\begin{bmatrix} y_k \\ z_k \end{bmatrix} = \begin{bmatrix} 0.9 & 0.2 \\ 0.1 & 0.8 \end{bmatrix} \begin{bmatrix} y_{k-1} \\ z_{k-1} \end{bmatrix} \Rightarrow \lambda_1 = 1.0, \lambda_2 = 0.7$$

$$\begin{bmatrix} y_k \\ z_k \end{bmatrix} = \begin{bmatrix} 0.9 & 0.2 \\ 0.1 & 0.8 \end{bmatrix}^k \begin{bmatrix} y_0 \\ z_0 \end{bmatrix} = \frac{1}{3} \begin{bmatrix} 2 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1^k & 0 \\ 0 & (0.7)^k \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & -2 \end{bmatrix} \begin{bmatrix} y_0 \\ z_0 \end{bmatrix}$$

□ A *Markov* matrix  $A$  is nonnegative with each column adding to 1.

(a)  $\lambda_1 = 1$  is an eigenvalue with a nonnegative eigenvector  $\mathbf{x}_1$ .

(b) The other eigenvalues satisfy  $|\lambda_i| \leq 1$ .

(c) If any power of  $A$  has all positive entries, and the other  $|\lambda_i| < 1$ . Then  $A^k \mathbf{u}_0$  approaches the steady state of  $\mathbf{u}_\infty$  which is a multiple of  $\mathbf{x}_1$  as long as the projection of  $\mathbf{u}_0$  in  $\mathbf{x}_1$  is not zero.

◇ Check Perron-Fröbenius theorem in Strang's book.

## $e^A$ and Differential Equations

$$\clubsuit e^A = I + \frac{A}{1!} + \frac{A^2}{2!} + \cdots + \frac{A^m}{m!} + \cdots$$

$$\clubsuit \frac{du}{dt} = -\lambda u \Rightarrow u(t) = e^{-\lambda t}u(0)$$

$$\clubsuit \frac{d\mathbf{u}}{dt} = -A\mathbf{u} = \begin{bmatrix} -2 & 1 \\ 1 & -2 \end{bmatrix} \mathbf{u} \Rightarrow \mathbf{u}(t) = e^{-tA}\mathbf{u}(0)$$

$\clubsuit$   $A = U\Lambda U^t$  for an orthogonal matrix  $U$ , then

$$e^A = Ue^\Lambda U^t = U \text{diag}[e^{\lambda_1}, e^{\lambda_2}, \dots, e^{\lambda_n}] U^t$$

$\clubsuit$  Solve  $x''' - 3x'' + 2x' = 0$ .

Let  $y = x'$ ,  $z = y' = x''$ , and let  $\mathbf{u} = [x, y, z]^t$ . The problem is reduced to solving

$$\mathbf{u}' = A\mathbf{u} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & -2 & 3 \end{bmatrix} \mathbf{u}$$

Then

$$\mathbf{u}(t) = e^{tA}\mathbf{u}(0) = \begin{bmatrix} \frac{1}{\sqrt{21}} & \frac{1}{\sqrt{3}} & 1 \\ \frac{2}{\sqrt{21}} & \frac{1}{\sqrt{3}} & 0 \\ \frac{4}{\sqrt{21}} & \frac{1}{\sqrt{3}} & 0 \end{bmatrix} \begin{bmatrix} e^{2t} & 0 & 0 \\ 0 & e^t & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 & -2.2913 & 2.2913 \\ 0 & 3.4641 & -1.7321 \\ 1 & -1.5000 & 0.5000 \end{bmatrix} \mathbf{u}(0)$$

## Problems Solved by Matlab

Let  $A, B, H, \mathbf{x}, \mathbf{y}, \mathbf{u}, \mathbf{b}$  be matrices and vectors defined below, and  $H = I - 2\mathbf{u}\mathbf{u}^t$

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & 6 & 0 \\ -2 & 7 & 2 \end{bmatrix}, \quad B = \begin{bmatrix} -3 & 1 & 0 \\ 1 & -3 & 0 \\ 0 & 0 & 3 \end{bmatrix}, \quad \mathbf{u} = \begin{bmatrix} 1/2 \\ 1/2 \\ 1/2 \\ 1/2 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 6 \\ 2 \\ -5 \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}, \quad \mathbf{y} = \begin{bmatrix} 1 \\ 1 \\ -1 \\ -1 \end{bmatrix}$$

1. Let  $A=LU=QR$ , find  $L, U; Q, R$ .
2. Find determinants and inverses of matrices  $A, B$ , and  $H$ .
3. Solve  $A\mathbf{x} = \mathbf{b}$ , how to find the number of floating-point operations are required?
4. Find the ranks of matrices  $A, B$ , and  $H$ .
5. Find the characteristic polynomials of matrices  $A$  and  $B$ .
6. Find 1-norm, 2-norm, and  $\infty$ -norm of matrices  $A, B$ , and  $H$ .
7. Find the eigenvalues/eigenvectors of matrices  $A$  and  $B$ .
8. Find matrices  $U$  and  $V$  such that  $U^{-1}AU$  and  $V^{-1}BV$  are diagonal matrices.
9. Find the singular values and singular vectors of matrices  $A$  and  $B$ .
10. Randomly generate a  $4 \times 4$  matrix  $C$  with  $0 \leq C(i, j) \leq 9$ .