

# Binary numbers

## Binary number system

- ▶ Computer (electronic) systems prefer binary numbers
- ▶ Binary number: represent a number in **base-2**

a. Base ten system

3	7	5
Hundreds	Ten	One

Representation: 375  
Position's quantity

$3 \times 10^2 + 7 \times 10^1 + 5 \times 10^0$

b. Base two system

1	0	1	1
Eight	Four	Two	One

Representation: 1011  
Position's quantity

$1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1$

Bit pattern	Hexadecimal representation
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	8
1001	9
1010	A
1011	B
1100	C
1101	D
1110	E
1111	F

- ▶ Some terminology
  - ▶ **Bit:** a binary digit (0 or 1)
  - ▶ **Hexadecimal notation** (十六進位)
    - ▶ Represents each 4 bits by a single symbol
    - ▶ Example: A3 denotes 1010 0011

## Outline

- ▶ Integer: decimal-binary conversion
- ▶ Integer addition
- ▶ Negative integer (2's complement representation)
- ▶ Real numbers (floating point representation)

## Binary to decimal

- ▶ What is the decimal number of  $100101_2$ ?

Binary pattern	1	0	0	1	0	1	
							$2^0$
						1 x one = 1	$2^1$
					0 x two = 0		$2^2$
					1 x four = 4		$2^3$
					0 x eight = 0		$2^4$
					0 x sixteen = 0		$2^5$
					1 x thirty-two = 32		
							37 Total
							Value of bit      Position's quantity

## Decimal to binary

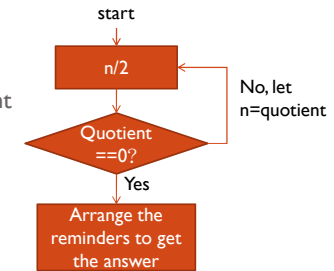
- ▶ What is the binary number of  $13_d$ ?
  - ▶ Step 1: Divide the value by 2 and record the remainder
  - ▶ Step 2: If quotient is not zero, use the quotient as the new value and repeat step 1
  - ▶ Step 3: The binary representation is the recorded remainders listed from right to left

$2 \overline{) 13}$	Quotient = 6
$2 \overline{) 6}$	Remainder = 1
$2 \overline{) 3}$	Quotient = 3
$2 \overline{) 1}$	Remainder = 0
$2 \overline{) 0}$	Quotient = 1
$2 \overline{) 0}$	Remainder = 1
$2 \overline{) 0}$	Quotient = 0
$2 \overline{) 1}$	Remainder = 0
$2 \overline{) 0}$	Quotient = 0
$2 \overline{) 1}$	Remainder = 1



## Algorithm

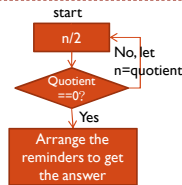
- ▶ A **finite** sequence of instructions to describe a systematical method to solve a problem
- ▶ We can represent the 'decimal-to-binary' algorithm by a flow chart
  - ▶ It has starting point [step 1]
  - ▶ Step 2 is a condition statement
  - ▶ Step 1+step 2 is a loop statement
  - ▶ The problem size is shrunk after each loop
  - ▶ The loop can be terminated [when quotient ==0]



- ▶ Homework: write an algorithm to convert binary to decimal

## 十進位轉二進位演算法 Three algorithms

- ▶ Problem: converting  $n_d$  to  $m_b$ .
- ▶ Algorithm 1: as mentioned in the last class
- ▶ Algorithm 2:
  - ▶  $m_b = 0$
  - ▶ For  $i = 1$  to  $n_d$ 
    - ▶  $m_b = m_b + 1$
- ▶ Algorithm 3:
  - ▶  $m_b = 0$
  - ▶ While  $n_d$  is not 0
    - ▶ Find  $2^k \leq n_d < 2^{k+1}$
    - ▶  $m_b = m_b + 2^k$
    - ▶  $n_d = n_d - 2^k$



What is the cost of  $m_b = m_b + 1$ ?

What is the cost of finding  $k$ ?

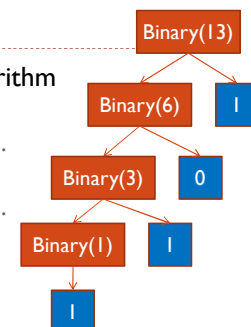
The cost of  $m_b + 2^k$  is just putting 1 at the  $k$ th position of  $m_b$

What is the cost of  $n_d = n_d - 2^k$ ?



## Recursive algorithm

- ▶ The first algorithm is a recursive algorithm
  - ▶ To answer what the binary of  $13_d$  is, you need to answer what the binary of  $6_d$  is.
  - ▶ To answer what the binary of  $6_d$  is, you need to answer what the binary of  $3_d$  is.
  - ▶ To answer what the binary of  $3_d$  is, you need to answer what the binary of  $1_d$  is.
- ▶ Recursive algorithm
  - ▶ Turn a big problem into one or several smaller subproblems
  - ▶ Each subproblem is identical to the original one except size
    - ▶ Thus, they can be solved by the same method.
  - ▶ Need a termination condition.



## Integer addition

- ▶ One bit addition

$$\begin{array}{r} 0 \\ + 0 \\ \hline 0 \end{array} \quad \begin{array}{r} 1 \\ + 0 \\ \hline 1 \end{array} \quad \begin{array}{r} 0 \\ + 1 \\ \hline 1 \end{array} \quad \begin{array}{r} 1 \\ + 1 \\ \hline 10 \end{array}$$

- ▶ What is  $5_d + 9_d$  using binary number representation?

$$\begin{array}{r} 0101 \\ + 1001 \\ \hline 1110 \end{array} \quad \begin{array}{l} \leftarrow 5 \\ \leftarrow 9 \\ \leftarrow 14 \end{array}$$



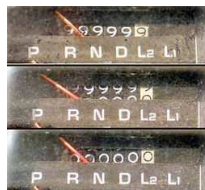
## Another example

$$\begin{array}{r} \phantom{00}1111 \\ 00111010 \\ + 00011011 \\ \hline 01010101 \end{array}$$



## Binary number in computer systems

- ▶ Mathematically, a binary integer can have arbitrary number of bits.
- ▶ In computer systems, all data has a limited number of bits.
  - ▶ For example, there are different sized (data type) integers
    - ▶ char: 8 bits: [0, 255]
    - ▶ unsigned short: 16 bits: [0, 65535]
    - ▶ unsigned int: 32 bits: [0, 4294967295]
- ▶ **Overflow:** when adding two integers, the result exceeds the numerical range of the data type
  - ▶ Ex: (char) 123 + (char) 234 = (char) 101



## How about negative integers?

- ▶ We use sign to distinguish positive and negative numbers.
- ▶ In computer system, we can use a bit to present the sign.
  - ▶ In a 4 bit integer, 0001 for 1 and 1001 for -1. (sign bit)
- ▶ Good for notation, but difficult for calculation.
  - ▶ One bit subtraction

$$\begin{array}{r} 0 \phantom{0} \phantom{0} \phantom{0} \\ - 0 \phantom{0} \phantom{0} \phantom{0} \\ \hline 0 \phantom{0} \phantom{0} \phantom{0} \end{array} \quad \begin{array}{r} 1 \phantom{0} \phantom{0} \phantom{0} \\ - 0 \phantom{0} \phantom{0} \phantom{0} \\ \hline 1 \phantom{0} \phantom{0} \phantom{0} \end{array} \quad \begin{array}{r} 1 \phantom{0} \phantom{0} \phantom{0} \\ - 1 \phantom{0} \phantom{0} \phantom{0} \\ \hline 0 \phantom{0} \phantom{0} \phantom{0} \end{array} \quad \begin{array}{r} 0 \phantom{0} \phantom{0} \phantom{0} \\ - 1 \phantom{0} \phantom{0} \phantom{0} \\ \hline -1 \phantom{0} \phantom{0} \phantom{0} \end{array}$$

- ▶ Example: 4-1

$$\begin{array}{r} \phantom{00}11 \\ 0100 \\ - 0001 \\ \hline 0011 \end{array}$$





## 2's complement addition

▶ Examples

Problem in base ten	Problem in two's complement	Answer in base ten
$\begin{array}{r} 3 \\ + 2 \\ \hline \end{array}$	$\begin{array}{r} 0011 \\ + 0010 \\ \hline 0101 \end{array}$	5
$\begin{array}{r} -3 \\ + -2 \\ \hline \end{array}$	$\begin{array}{r} 1101 \\ + 1110 \\ \hline 1011 \end{array}$	-5
$\begin{array}{r} 7 \\ + -5 \\ \hline \end{array}$	$\begin{array}{r} 0111 \\ + 1011 \\ \hline 0010 \end{array}$	2

## Another type of overflow

▶ What is 5+4 in signed 4 bit representation?

$$5_d + 4_d = 0101_b + 0100_b = 1001_b$$

▶ This is another type of **overflow**

- ▶ Adding two positive numbers results a negative number; or adding two negative numbers results a positive number.
- ▶ A 4 bits 2's complement system can only represent 7~ -8

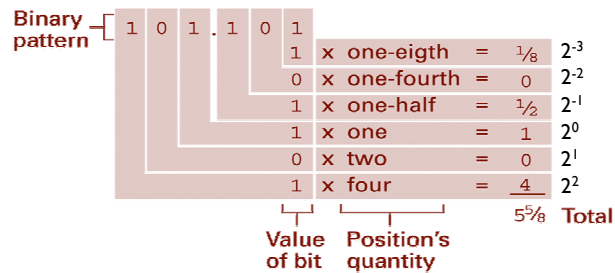
b. Using patterns of length four

Bit pattern	Value represented
0111	7
0110	6
0101	5
0100	4
0011	3
0010	2
0001	1
0000	0
1111	-1
1110	-2
1101	-3
1100	-4
1011	-5
1010	-6
1001	-7
1000	-8

## Fraction

▶ The binary number of fractions.

▶ 5.625



## Fraction point

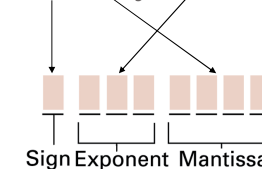
▶ To represent a wide range of numbers, we allow the decimal point to "float".

$$40.1_d = 4.01_d \times 10^1 = 401_d \times 10^{-1} = 0.401_d \times 10^2$$

- ▶ It is just like the scientific notation of numbers.

$$101.101_b = +1.01101_b \times 2^{2d} = +1.01101_b \times 2^{10_b}$$

- ▶ This is called the **floating point** representation of fractions.



### Coding the value of $2^5/8$

Binary representation:  $10.101$

Normalization:  $0.10101 \times 2^2$

Bit pattern: 0 1 1 0 1 0 1 0

Sign: 0, Exponent: 11, Mantissa: 0101 (truncated)

Exponent uses excess notation

Bit pattern	Value represented
111	3
110	2
101	1
100	0
011	-1
010	-2
001	-3
000	-4

### Signed number representations

Comparison of 4 bit signed integer representation by sign-bit notation, 2's complement, and excess notation

	Sign-bit notation	2's complement	Excess notation
8			1111
7	0111	0111	1110
6	0110	0110	1101
5	0101	0101	1100
4	0100	0100	1011
3	0011	0011	1010
2	0010	0010	1001
1	0001	0001	1000
0	0000, 1000	0000	0111
-1	1001	1111	0110
-2	1010	1110	0101
-3	1011	1101	0100
-4	1100	1100	0011
-5	1101	1011	0010
-6	1110	1010	0001
-7	1111	1001	0000
-8		1000	

### Floating-point numbers

In C (and most programming languages), there are two data types for real numbers

Data type	Size	Structure	Range	Precision
float	32 bits	Sign: 1 bit Exponent: 8 bits Mantissa: 23 bits	$\pm \sim 10^{-44.85}$ to $\sim 10^{38.53}$	$\sim 10^8$
double	64 bits	Sign: 1 bit Exponent: 11 bits Mantissa: 52 bits	$\pm \sim 10^{-323.3}$ to $\sim 10^{308.3}$	$\sim 10^{16}$

### Truncation error

- Mantissa field is not large enough
  - $2^5/8 = 2.625 \Rightarrow 2.5 + \text{round off error } (0.125)$
- Nonterminating representation
  - $0.1 = 1/16 + 1/32 + 1/256 + 1/512 + \dots$
  - Change the unit of measure
- Order of computation:
  - $2.5 + 0.125 + 0.125 \Rightarrow 2.5+0+0 = 2.5$
  - $2.5 + (0.125+0.125) \Rightarrow 2.5+0.25 = 2.75$

More in 数值分析